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hal-02868664

HAL Id: hal-02868664
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Submitted on 15 Jun 2020

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Removing inessential points in $c$-and $A$-optimal design

Luc Pronzato* and Guillaume Sagnol†

June 15, 2020

Abstract

A design point is inessential when it does not contribute to an optimal design, and can therefore be safely discarded from the design space. We derive three inequalities for the detection of such inessential points in $c$-optimal design: the first two are direct consequences of the equivalence theorem for $c$-optimality; the third one is derived from a second-order cone programming formulation of $c$-optimal design. Elimination rules for $A$-optimal design are obtained as a byproduct. When implemented within an optimization algorithm, each inequality gives a screening test that may provide a substantial acceleration by reducing the size of the problem online. Several examples are presented with a multiplicative algorithm to illustrate the effectiveness of the approach.

Keywords: $A$-optimal design; $c$-optimal design; Inessential Point; Screening Test; Support Point.
AMS subject classifications: 62K05, 62J07, 90C46

1 Introduction

Let $\mathbb{H} = \{\mathbf{H}_i, i = 1, \ldots, q\}$ denote a set of $m \times m$ symmetric non-negative definite matrices and denote by $\mathcal{H}$ its convex hull. We assume that the linear span of $\mathbb{H}$ contains a nonsingular matrix. A design corresponds to a probability measure on $\mathbb{H}$, that is, to vector of weights $\mathbf{w} = (w_1, \ldots, w_q)\top$ in the probability simplex $\mathcal{P}_q = \{\mathbf{w} \in \mathbb{R}^q, w_i \geq 0 \text{ for all } i \text{ and } \sum_{i=1}^q w_i = 1\}$. The extension to designs $\xi$ defined as probability measures on an arbitrary compact set of matrices $\mathbb{H}_t$ is straightforward, but for the sake of simplicity in the rest of the paper we shall restrict the presentation to the discrete case. A $c$-optimal design, associated with a nonzero vector $\mathbf{c} \in \mathbb{R}^m$, is defined by a vector $\mathbf{w}^\ast$ that minimizes

$$\phi(\mathbf{w}) = \Phi[\mathbf{M}(\mathbf{w})] = \mathbf{c}^\top \mathbf{M}^{-}\mathbf{(w)c}$$

with respect to $\mathbf{w}$, with $\mathbf{M}(\mathbf{w})$ the (information) matrix

$$\mathbf{M}(\mathbf{w}) = \sum_{i=1}^q w_i \mathbf{H}_i$$

and $\mathbf{M}^{-}$ a pseudo-inverse of $\mathbf{M}$. We set $\phi(\mathbf{w}) = \infty$ when $\mathbf{c}$ does not belong to the column space of $\mathbf{M}(\mathbf{w})$; the value of $\phi(\mathbf{w})$ does not depend on the choice of the pseudo-inverse $\mathbf{M}^{-}$.

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A typical example is when $H_i = a_i a_i^\top$ for all $i$, with $a_i$ the vector of regressors in the model with observations

$$Y_i = a_i^\top \theta + \varepsilon_i,$$  

(1.3)

parameters $\theta$ and i.i.d. errors $\varepsilon_i$ having zero mean and variance $\sigma^2$. Suppose that $n$ observations are collected according to the design $w \in \mathcal{P}_q$, with $w$ such that $nw_i \in \mathbb{N}$ for all $i$. Then, $\sigma^2 \phi(w)/n$ is the variance of $c^\top \hat{\theta}^n$, with $\hat{\theta}^n$ the Best Linear Unbiased Estimator (BLUE) of $\theta$.

When the observations $Y_i$ are multivariate, with $Y_i = A_i^\top \theta + \varepsilon_i$ in $\mathbb{R}^q$ where the errors $\varepsilon_i$ have the $q \times q$ covariance matrix $\Sigma_i$, $\sigma^2 \phi(w)/n$ is still equal to $\text{var}(c^\top \hat{\theta}^n)$, with now

$$H_i = A_i \Sigma_i^{-1} A_i^\top,$$  

(1.4)

in $M(w)$, and each $H_i$ may have rank larger than 1; see for instance Harman and Trnovská (2009) for the case of $D$-optimal design that maximizes $\det[M(w)]$. When setting a normal prior $\mathcal{N}(\theta^0, \Sigma)$ on $\theta$ in the regression model (1.3), with $\Sigma$ having full rank, and when the errors $\varepsilon_i$ are normal $\mathcal{N}(0, \sigma^2)$, the posterior variance of $c^\top \theta$ after $n$ observations is $\sigma^2 \phi(w)/n$, with $\phi(w)$ given by (1.1) and

$$H_i = a_i a_i^\top + \sigma^2 \Sigma^{-1}/n = [a_i \ n^{-1/2} \sigma \Sigma^{-1/2}] [a_i \ n^{-1/2} \sigma \Sigma^{-1/2}]^\top$$  

(1.5)

in (1.2); one may refer in particular to Pilz (1983) for a thorough exposition on optimal design for Bayesian estimation (Bayesian optimal design). When $H_i$ is given by (1.5), $c$-optimal design is equivalent to the minimization of $c^\top M^{-1}(w)c'$ where $c' = \Sigma^{1/2} c$, $M(w) = \sum_{i=1}^q w_i H_i'$, and $H_i' = \Sigma^{1/2} a_i a_i^\top \Sigma^{1/2} + (\sigma^2/n) I_m$ for all $i$, so that we can assume that $\Sigma = I_m$, the $m$-dimensional identity matrix, without any loss of generality. Other cases of interest correspond to $A$- and $L$-optimality, they will be considered in Section 4.

In general, not all $H_i$ contribute to an optimal design $w^*$ as some of the weights $w_i^*$ equal zero. We shall say that such an $H_i$, which does not support an optimal design, is inessential. In Section 2 we show that any design $w$ yields two simple inequalities that must be satisfied by an $H_i$ that contributes to an optimal design. Any $H_i$ that does not satisfy these screening inequalities is therefore inessential and can be safely eliminated from $H$ — we shall sometimes speak of (design-) point elimination instead of matrix elimination, which is all the more appropriate when $H_i$ has the form (1.5) and the $a_i$ are seen as points in $\mathbb{R}^m$. The idea originated from Pronzato (2003) and Harman and Pronzato (2007) for $D$-optimal design (and for the construction of the minimum-volume ellipsoid containing a set of points); it was further extended to Kiefer’s $\varphi_p$ criteria (Pronzato, 2013), to $E$-optimal design (Harman and Rosa, 2019), and to the elimination of inessential points in the smallest enclosing ball problem (Pronzato, 2019). In Section 3, we use the Second-Order Cone Programming formulation of $c$-optimal design of Sagnol (2011) to derive a third screening inequality. Since $A$-optimal design can be seen as $c$-optimal design in a higher dimensional space, see Sagnol (2011), the same screening inequalities can be used to eliminate inessential points in $A$-optimal design. This is considered in Section 4, where a comparison is made with the elimination method of Pronzato (2013). The inequalities proposed in the paper become more stringent when $w$ approaches an optimal design $w^*$, which can be used to accelerate the construction of an optimal design. The examples of Section 5 provide an illustration for the particular case of a “multiplicative algorithm”, classical in optimal design, but elimination of inessential could benefit any other algorithm.
2 Elimination of inessential points by the equivalence theorem

2.1 Optimality conditions

The function $M \rightarrow \Phi(M) = c^\top M c$ is convex on the set $M^2$ of symmetric non-negative definite matrices. When $M$ has full rank, the directional derivative of $\Phi$ at $M$ in the direction $M'$, defined by

$$F_\Phi(M, M') = \lim_{\alpha \to 0^+} \frac{\Phi((1 - \alpha)M + \alpha M') - \Phi(M)}{\alpha},$$

equals

$$F_\Phi(M, M') = -c^\top M^{-1}(M' - M)M^{-1}c.$$ Denote $\Phi_* = \Phi(M_*) = \min_{w \in \mathcal{H}} \Phi[M(w)]$, with $M_* = M(w^*)$ a $c$-optimal matrix in $\mathcal{H}$, and define

$$\delta = \delta(M) = \frac{\max_{M' \in \mathcal{H}} c^\top M^{-1}M'c - c^\top M^{-1}c}{c^\top M^{-1}c},$$

(2.1)

where the maximum is attained for $M'$ equal to some $H_i \in \mathcal{H}$. The convexity of $\Phi$ implies that $\Phi(M') \geq \Phi(M) + F_\Phi(M, M')$ for all $M, M'$ in $\mathcal{H}$, that is,

$$\Phi(M') \geq 2\Phi(M) - c^\top M^{-1}M'M^{-1}c \text{ for all } M, M' \in \mathcal{H}.$$ Similarly, using the stronger property that the function $M \rightarrow 1/\Phi(M)$ is concave, we obtain

$$\Phi(M') \geq \frac{\Phi^2(M)}{c^\top M^{-1}M'M^{-1}c} \text{ for all } M, M' \in \mathcal{H}.$$ (2.2)

A matrix $M_*$ is $c$-optimal if and only if $F_\Phi(M_*, M) \geq 0$ for all $M \in \mathcal{H}$. Therefore, when $M_*$ has full rank (which is necessarily the case when rank$(H) = m$ for all $H \in \mathcal{H}$), $M_*$ is $c$-optimal if and only if

$$c^\top M_*^{-1}H_*M_*^{-1}c \leq \Phi_* \text{ for all } H \in \mathcal{H}.$$ (2.3)

Substituting $H_i$ for $H$, since $M_* = \sum_i w_i^* H_i$ we obtain that

$$c^\top M_*^{-1}H_*M_*^{-1}c = \Phi_* \text{ for all } H_i \text{ such that } w_i^* > 0.$$ (2.4)

The properties (2.3, 2.4) correspond to what is called “equivalence theorem” in optimal design, which is central to major developments in the field since the pioneering work of Kiefer and Wolfowitz (1960) on $D$-optimal design; see for instance Silvey (1980); Pukelsheim (1993). The equivalence theorem takes a more complicated form when $M_*$ is singular, so that $\Phi$ is not differentiable at $M_*$; see Pukelsheim (1993, Chap. 2, 4). When $c$ belongs to the column space of $M$, the directional derivative at $M$ in the direction $M'$ is then

$$F_\Phi(M, M') = \max_{A:MAM=M} -c^\top A^\top (M' - M)Ac,$$

and the equivalence theorem indicates that $M_*$ is $c$-optimal if and only if $c$ belongs to the column space of $M_*$ and there exists a generalized inverse $M_*^-$ of $M_*$ (i.e., such that $M_*M_*^-M_* = M_*$) satisfying $c^\top M_*^{-1}H_*M_*^{-1}c \leq \Phi_*$ for all $H \in \mathcal{H}$, with equality for all $H_i$ with $w_i^* > 0$.  

When all $\mathbf{H}_i$ have rank one, Elfving’s theorem (1952) gives a simpler necessary and sufficient condition for c-optimality; see also (Pukelsheim, 1993, p. 50). In (Sagnol, 2011), Elfving’s theorem is extended to the case of matrices $\mathbf{H}_i$ having rank larger than one; see also Section 3. The following lemma summarizes the properties that we shall use to derive screening bounds based on the equivalence theorem.

**Lemma 2.1.** A c-optimal matrix $\mathbf{M}_* = \mathbf{M}(\mathbf{w}^*)$ satisfies

$$
\frac{1}{1 + \delta} \Phi(\mathbf{M}) \leq \Phi_* = \Phi(\mathbf{M}_*) \leq \Phi(\mathbf{M})
$$

and there exists a generalized inverse $\mathbf{M}_*^{-1}$ of $\mathbf{M}_*$ such that

$$
c^\top \mathbf{M}_*^{-1} \mathbf{H}_i \mathbf{M}_*^{-1} c = \Phi_* \text{ for all } \mathbf{H}_i \text{ such that } w_i^* > 0.
$$

**Proof.** The first inequality in (2.5) is obtained by substituting $\mathbf{M}_*$ for $\mathbf{M}$ in (2.2) and using (2.1) and $\max_{\mathbf{M}_* \in \mathcal{H}} c^\top \mathbf{M}_*^{-1} \mathbf{M}_*^{-1} c \geq c^\top \mathbf{M}_*^{-1} \mathbf{M}_*^{-1} c$; the second inequality follows from the optimality of $\mathbf{M}_*$; (2.6) is part of the equivalence theorem.

\[\square\]

### 2.2 Elimination rule

The following property allows us to eliminate inessential matrices $\mathbf{H}_i$ from $\mathcal{H}$. The proof is given in Appendix A. We denote by $\lambda_{\max}(\mathbf{A})$ and $\lambda_{\min}(\mathbf{A})$ the maximum and minimum eigenvalues, respectively, of the matrix $\mathbf{A}$.

**Theorem 2.1.** Let $\mathbf{M}$ be any full rank matrix in $\mathcal{H}$, $\delta = \delta(\mathbf{M})$ be given by (2.1), and $\mathbf{H}_i$ be a matrix in $\mathcal{H}$ such that

$$
B_1(\mathbf{M}, \mathbf{H}_i) = \left[1 - \Delta_i \left(\frac{\delta}{1 + \delta}\right)^{1/2}\right] \Phi(\mathbf{M}) - c^\top \mathbf{M}_i^{-1} \mathbf{H}_i \mathbf{M}_i^{-1} c > 0,
$$

or

$$
B_2(\mathbf{M}, \mathbf{H}_i) = \gamma(\kappa_i, \varphi) \Phi(\mathbf{M}) - c^\top \mathbf{M}_i^{-1} \mathbf{H}_i \mathbf{M}_i^{-1} c > 0,
$$

where $\Delta_i = \lambda_{\max}(\Omega_i) - \lambda_{\min}(\Omega_i)$, $\kappa_i = \lambda_{\max}(\Omega_i)/\lambda_{\min}(\Omega_i)$, with $\Omega_i = \mathbf{M}_i^{-1/2} \mathbf{H}_i \mathbf{M}_i^{-1/2}$, and

$$
\gamma(\kappa, \varphi) = \frac{\cos^2(\omega - \varphi) + \kappa \sin^2(\omega - \varphi)}{\cos^2(\omega) + \kappa \sin^2(\omega)}, \quad \text{with} \quad \omega = \omega(\kappa, \varphi) = \frac{1}{2} \left\{ \arccos \left[ \frac{\kappa - 1}{\kappa + 1} \cos(\varphi) \right] + \varphi \right\},
$$

$\varphi = \arccos((1 + \delta)^{-1/2})$ and $\mathbf{M}_i^{-1/2}$ a square-root matrix of $\mathbf{M}_i^{-1}$. Then $\mathbf{H}_i$ does not support a c-optimal design.

**Remark 2.1.** (i) In Theorem 2.1, (2.7) remains valid when the matrices $\mathbf{H}_i$ do not have full rank provided that $\mathbf{M}$ is nonsingular. In that case, on the one hand $\lambda_{\min}(\Omega_i) = 0 = \Delta_i$ can be computed more efficiently as $\Delta_i = \lambda_{\max}(\Omega_i) = \lambda_{\max}(\mathbf{A}_i^\top \mathbf{M}_i^{-1} \mathbf{A}_i)$ when we write $\mathbf{H}_i = \mathbf{A}_i \mathbf{A}_i^\top$. On the other hand, $\kappa_i$ becomes infinite and $\gamma(\kappa_i, \varphi) = 0$ in (2.8), which therefore cannot be used to eliminate singular matrices from $\mathcal{H}$. The examples in Section 5.1 will also illustrate that $B_2(\mathbf{M}, \mathbf{H}_i)$ is more sensitive than $B_1(\mathbf{M}, \mathbf{H}_i)$ to badly conditioned matrices $\mathbf{H}_i$.

(ii) The precise knowledge of $\Delta_i$ is not necessary to be able to eliminate a matrix $\mathbf{H}_i$: any $\mathbf{H}_i$ such that $c^\top \mathbf{M}_i^{-1} \mathbf{H}_i \mathbf{M}_i^{-1} c < \left\{1 - \overline{\Delta}_i [\delta/(1 + \delta)]^{1/2}\right\} \Phi(\mathbf{M})$ with $\overline{\Delta}_i > \Delta_i$ can be safely removed from $\mathcal{H}$. Also, for a fixed $\varphi$, $\gamma(\kappa, \varphi)$ is a decreasing function of $\kappa$, so that an overestimation of $\lambda_{\max}(\Omega_i)$ and an underestimation of $\lambda_{\min}(\Omega_i)$ still provide a safe elimination of $\mathbf{H}_i$ via (2.8).

(iii) The factor $\left\{1 - \Delta_i [\delta/(1 + \delta)]^{1/2}\right\}$ tends to one as $\delta$ tends to zero, which renders the sieve based on (2.7) finer as $\mathbf{M}$ gets closer to an optimal matrix $\mathbf{M}_*$. Since for a given $\kappa$, $\gamma(\kappa, \varphi)$ is a decreasing function of $\varphi$, the sieve based on (2.8) becomes also finer as $\delta$ decreases.
Theorem 3.2. Let indices $i$ can be safely eliminated when $F$ or any factor. The proof is given in Appendix B.

Theorem 2.2. For any $\delta, \epsilon > 0$ and any dimension $m \geq 2$, there exists a set $H$, a matrix $M$ in $H$ (that is, a design on $H$) and a $H_i \in H$ supporting an optimal design such that $\delta$ equals $\delta(M)$ given by (2.1) and $B_1(M, H_i) + \epsilon > 0$, $B_2(M, H_i) + \epsilon > 0$.

The artificial example used for the proof of Theorem 2.2 suggests that (2.8) can be more accurate than (2.7). It is not always the case, however, and Example 1 of Section 5 will present situations where (2.7) eliminates more points than (2.8).

3 Elimination of inessential points by the generalized Elfving theorem

3.1 A criterion based on Second Order Cone Programming

Write each $H_i$ in $H$ as $H_i = A_i A_i^\top$ with $A_i$ a $m \times \ell$ matrix; see for instance (1.4) and (1.5). Sagnol (2011) gives a generalization of Elfving theorem to multiresponse models; see also Dette and Holland-Letz (2009) for the case heteroscedastic regression models and Dette (1993) for a model-robust version of c-optimality. A consequence of this generalized Elfving theorem is that c-optimal design is equivalent to a Second Order Cone Programming problem (Sagnol, 2011); Theorem 3.2 exploits this equivalence to derive a sufficient condition for elimination of an inessential $H_i$.

Theorem 3.1 (Sagnol, 2011). Let $u^*$ and $(\mu^*, h_1^*, \ldots, h_q^*)$ be respectively primal and dual solutions of the second-order cone programs:

\[ P - \text{SOCP} : \max_{u \in \mathbb{R}^m} c^\top u \quad \text{subject to} \quad \|A_i^\top u\| \leq 1 \text{ for all } i = 1, \ldots, q, \]

\[ D - \text{SOCP} : \min_{\mu \in \mathbb{R}^q, h_1, \ldots, h_q} \sum_{i=1}^q \mu_i \quad \text{subject to} \quad c = \sum_{i=1}^q A_i h_i \quad \text{and} \quad \|h_i\| \leq \mu_i \text{ for all } i = 1, \ldots, q. \]

Then, $w^* = \mu^*/(\sum_{i=1}^q \mu_i)$ defines a c-optimal design; that is, it minimizes $\phi(w)$ given by (1.1) with respect to $w \in \mathcal{P}_q$. Moreover, $\phi(w^*) = (\sum_{i=1}^q \mu_i^*)^2 = (c^\top u^*)^2$ and strictly positive $w_i^*$ correspond to indices $i$ such that $\|A_i^\top u^*\| = 1$.

Theorem 3.2. Let $M = M(w)$ be any full-rank matrix in $H$. Any matrix $H_i \in H$ such that

\[ B_3(M, H_i) = 1 + \sup_{\beta \geq \lambda_{\max}(A_i^\top M^{-1} A_i)} \left\{ \frac{\Phi(M)}{(1 + \delta)(c^\top (\beta M - H_i)^{-1} c)} - \beta \right\} > 0 \quad (3.1) \]

does not support a c-optimal design.

The proof is given in Appendix C. It relies on the construction of an upper bound $t^*$ on $\sup_{z \in \mathcal{F}(u)} \|A_i^\top z\|$, where $\mathcal{F}(u) = \{z \in \mathbb{R}^m : c^\top z \geq c^\top u \text{ and } \|A_j^\top z\| \leq 1, j = 1, \ldots, q\}$; $H_i$ can be safely eliminated when $t^* < 1$. The value of $t^*$ is obtained by solving a one-dimensional optimization problem,

\[ t^* = \inf_{\beta \geq \lambda_{\max}(A_i^\top M^{-1} A_i)} \beta - \frac{(c^\top u)^2}{c^\top (\beta M - H_i)^{-1} c}, \quad (3.2) \]

where the function to be minimized is defined by continuity at $\beta = \lambda_{\max}(A_i^\top M^{-1} A_i)$, see Section 3.2.
Remark 3.1. (i) Theorem 3.2 is valid whatever the rank of matrices $H_i$. We have $\lambda_{\text{max}}(A_i^T M^{-1} A_i) = \lambda_{\text{max}}(\Omega_i)$, where $\Omega_i = M^{-1/2} H_i M^{-1/2}$ is used in Theorem 2.1.

(ii) The determination of $t^*$ in (C.1) need not be very precise: any upper bound smaller than one also allows us to identify eliminate $H_i$. Similarly, restricting $\beta$ to be larger that an upper bound on $\lambda_{\text{max}}(A_i^T M^{-1} A_i)$ only makes the screening test more conservative.

The following theorem extends the result of Theorem 2.2 to the bound $B_3$, showing that it cannot be improved by a constant factor. The proof is in Appendix B.

**Theorem 3.3.** For any $\delta, \epsilon > 0$ and any dimension $m \geq 2$, there exists a set $H$, a matrix $M$ in $\mathcal{M}$ (that is, a design on $H$) and a $H_i \in H$ supporting an optimal design such that $\delta$ equals $\delta(M)$ given by (2.1) and $B_3(M, H_i) + \epsilon > 0$.

### 3.2 Numerical calculation of $B_3(M, H_i)$

Consider the function $f(\cdot)$ involved in the determination of $t^*$, see (C.1),

$$ f(\beta) = \beta - \frac{\Phi(M)}{(1 + \delta)c^T (\beta M - H_i)^{-1} c}, \quad \beta > \beta_{\text{min}} = \lambda_{\text{max}}(A_i^T M^{-1} A_i), $$

for which we define by continuity $f(\beta_{\text{min}}) = \beta_{\text{min}}$. For $\beta > \beta_{\text{min}}$, its first and second derivatives $f'$ and $f''$ are

$$ f'(\beta) = 1 - \frac{\Phi(M)[u^T M u]}{(1 + \delta)[c^T u]^2}, \quad f''(\beta) = 2 \frac{\Phi(M)\{[u^T M(\beta M - H_i)^{-1} M u][c^T u] - [u^T M u]^2\}}{(1 + \delta)[c^T u]^3} $$

where we have denoted $u = (\beta M - H_i)^{-1} c$. Since $f''(\beta) > 0$ follows from Cauchy-Schwarz inequality, $f$ is convex and can easily be minimized on $[\beta_{\text{min}}, \infty)$. (The convexity of $f$ is already clear by construction, as we obtained Problem (C.1) by partial minimization over a subset of variables of a convex semi-definite program, see the proof of Theorem 3.2 in Appendix C.) In the examples of Section 5 we use the simple dichotomy line search presented hereafter. The search can be stopped as soon as a $\beta$ is found such that $f(\beta) < 1$ (and $H_i$ can be eliminated), or when a lower bound $\underline{f}$ on $f(\beta)$ for $\beta \in [\beta_{\text{min}}, \infty)$ has been determined and is such that $\underline{f} \geq 1$ (and $H_i$ cannot be eliminated). Similarly to the definition of $f(\beta_{\text{min}})$, we define by continuity

$$ f'(\beta_{\text{min}}) = 1 - \frac{\Phi(M)}{(1 + \delta)[c^T M^{-1/2} P_i D_i P_i^T M^{-1/2} c]}, $$

where $P_i D_i P_i^T$ corresponds to the eigen-decomposition of $\Omega_i = M^{-1/2} H_i M^{-1/2}$ and $D_i$ is the diagonal matrix with $(D_i)_{k,k} = 1$ if $(D_i)_{k,k} = \beta_{\text{min}}$ and $(D_i)_{k,k} = 0$ otherwise.

The fact that $\underline{f}$ in Stage 2 gives a lower bound on $f$ is a simple consequence of its convexity. The procedure converges exponentially fast since the search interval is halved at each iteration of Stage 2: when $k_1$ expansions are needed at Stage 1, $k_2 = k_1 - \lfloor 1/ \log_2(\epsilon) \rfloor$ iterations are required at Stage 2, and in general $k_1$ is very small (with often $k_1 = 1$). Also, the initial tests at Stage 0 and the tests $f(\epsilon) < 1$ and $\underline{f} \geq 1$ at Stage 2 frequently yield early termination.

### 3.3 The special case $H_i = a_i a_i^T$

Elimination with (2.8) cannot be used when $H_i = a_i a_i^T$ since $B_2(M, H_i) = 0$, see Remark 2.1-(i). However, (2.7) and (3.1) remain valid, with $\Delta_i = \lambda_{\text{max}}(A_i^T M^{-1} A_i) = \lambda_{\text{max}}(\Omega_i) = a_i^T M^{-1} a_i$. More importantly, the minimum $t^*$ in (C.1) can be expressed explicitly, so that no iterative minimization is needed, and $B_1(M, a_i a_i^T) > 0$ implies $B_3(M, a_i a_i^T) > 0$; that is, $B_3$ is more efficient than $B_1$ in c-optimal design with $H_i = a_i a_i^T$. 


Theorem 3.4. For any full rank matrix $M$ in $\mathcal{H}$ and $H_i = a_i a_i^T$ in $\mathbb{H}$, $B_3(M, a_i a_i^T)$ defined by (3.1) satisfies

$$B_3(M, a_i a_i^T) > 0 \iff a_i^T M^{-1} a_i \in [0, 1) \cup \left( \delta_i^{(3)}, \Delta_i^{(3)} \right),$$

where

$$\delta_i^{(3)} = \frac{(1 + \delta) (c_i^T M^{-1} a_i)^2}{\Phi(M)}, \quad \Delta_i^{(3)} = 1 + \frac{1}{\delta} \left[ 1 - (\delta_i^{(3)})^{1/2} \right]^2.$$

Moreover, if $H_i$ satisfies (2.7), then it also satisfies (3.1).

Proof. Using the Sherman-Morrison-Woodbury formula for matrix inversion, we obtain that the function $f(\beta)$ in (C.1) is

$$f(\beta) = \beta - \frac{\beta \Phi(M)}{(1 + \delta) \Phi(M) + \delta (c_i^T M^{-1} a_i)^2}.$$  

Direct calculation shows that $f'(a_i^T M^{-1} a_i) = 1 - (a_i^T M^{-1} a_i)/\delta_i^{(3)}$, so that $f'(a_i^T M^{-1} a_i) \geq 0$ is equivalent to $\Delta_i = a_i^T M^{-1} a_i \leq \delta_i^{(3)}$, with $\delta_i^{(3)} \leq (\delta + 1)^2$ from the definition (2.1) of $\delta$.

When $f'(a_i^T M^{-1} a_i) < 0$, the minimizer $\beta^*$ equals $\tau [\tau + (c_i^T M^{-1} a_i)^2] / [\delta \Phi(M) (c_i^T M^{-1} a_i)^2]$, where $\tau = |c_i^T M^{-1} a_i| \delta^{1/2} [\Phi(M) (a_i^T M^{-1} a_i) - (c_i^T M^{-1} a_i)^2]^{1/2}$ (with $\Phi(M) (a_i^T M^{-1} a_i) \geq (c_i^T M^{-1} a_i)^2$ from Cauchy-Schwarz inequality). We then obtain that $f(\beta^*) < 1$ is equivalent to $\Delta_i < \Delta_i^{(3)}$, with $\Delta_i^{(3)} - \delta_i^{(3)} = (1 + 1/\delta) [\Phi(M) (a_i^T M^{-1} a_i) / \Phi^{1/2}(M) - 1] \geq 0$, which completes the proof of (3.3).

Suppose that (2.7) is satisfied; that is, $B_1(M, a_i a_i^T) > 0$. It is equivalent to

$$\Delta_i < \Delta_i^{(1)} = \left( \frac{1 + \delta}{\delta} \right)^{1/2} \left[ 1 - \frac{(c_i^T M^{-1} a_i)^2}{\Phi(M)} \right].$$
If $\Delta_i < 1$, then $B_3(M, a_i a_i^\top) > 0$. Therefore, we only need to consider the case when $\Delta_i \geq 1$. We have

$$B_1(M, a_i a_i^\top) > 0 \iff (c^\top M^{-1} a_i)^2 < T_i^{(1)}(1) = \Phi(M) \left[ 1 - \Delta_i \left( \frac{\delta}{1 + \delta} \right)^{1/2} \right].$$

Denote $T_i^{(3)} = \Phi(M) \Delta_i/(1 + \delta)$. $T_i^{(3)} - T_i^{(1)}$ is an increasing function of $\Delta_i$, and for $\Delta_i = 1$ we have $T_i^{(3)} - T_i^{(1)} = \Phi(M) \{ (\delta + (1 + \delta)^{-1/2} - \delta) / (1 + \delta) > 0$, showing that $T_i^{(3)} - T_i^{(1)} > 0$ for all $\Delta_i \geq 1$. Therefore, $B_1(M, a_i a_i^\top) > 0$ and $\Delta_i \geq 1$ imply $(c^\top M^{-1} a_i)^2 < T_i^{(3)}$, which is equivalent to $\Delta_i > \delta_i^{(3)}$. It just remains to show that $\Delta_i^{(1)} \leq \Delta_i^{(3)}$. Direct calculation gives

$$\Delta_i^{(3)} - \Delta_i^{(1)} = \frac{1}{\delta (1 + \delta)} \left\{ \left( \delta_i^{(3)} \right)^{1/2} - \frac{1}{1 + \left[ \delta/(1 + \delta) \right]^{1/2}} \right\}^2 \geq 0,$$

which completes the proof.

Figure 1 shows the regions $B_1(M, a_i a_i^\top) > 0$ and $B_3(M, a_i a_i^\top) > 0$ in the plane with $x$ variable $R_i = (c^\top M^{-1} a_i)^2 / \Phi(M)$ and $y$ variable $\Delta_i$. The yellow area is a forbidden zone, due to Cauchy-Schwarz inequality $\Delta_i \geq R_i$ and the definition of $\delta$ which imposes $R_i \leq \delta + 1$. $B_1(M, a_i a_i^\top) > 0$ is equivalent to $\Delta_i$ lying below the red line; it corresponds to the region colored in red. Define $\psi = (1 + 1/\delta)^{1/2} > 1$. The line $B_1 = 0$ crosses the main diagonal $\Delta_i = R_i$ at point $A$, with both coordinates equal to $R_i(A) = \psi/(1 + \psi)$. The blue diagonal dashed-line corresponds to $\Delta_i = \delta_i^{(3)}(R_i)$ and the blue dashed curve to $\Delta_i = \Delta_i^{(3)}(R_i)$. It is tangent to the line $B_1 = 0$ at point $B$ with coordinates $(1/\delta (1 + \psi)^2, 2 \psi/(1 + \psi))$; it is also tangent to the horizontal line $\Delta_i = 1$ at point $C$ with coordinates $(1/(1 + \delta), 1)$ and tangent to the main diagonal at point $D$ with coordinates $(1 + \delta, 1 + \delta)$; $C$ is also on the line $\Delta_i = \delta_i^{(3)}(R_i)$. The region $B_3(M, a_i a_i^\top) > 0$ corresponds to the red region $B_3(M, a_i a_i^\top) > 0$ augmented by the blue region, showing the superiority of $B_3$ over $B_1$. 

![Figure 1](image.png)
4 Elimination of inessential points in $A$-optimal design

4.1 New elimination rules for $A$-optimal design

Denote by $\text{vec}(A)$ the operation that stacks the columns of the $m \times n$ matrix $A$ into a single $mn$ dimensional column vector and by $A \otimes B$ the Kronecker product of matrices $A$ and $B$. Then, for any $m \times m$ matrix $A$ and $\mathbf{c} = \text{vec}(I_m)$ we have $\mathbf{c}^\top (I_m \otimes A) \mathbf{c} = \text{trace}(A)$, and therefore, for an information matrix $M(w)$ given by (1.2),

$$
\mathbf{c}^\top (I_m \otimes M^{-1}(w)) \mathbf{c} = \text{trace}[M^{-1}(w)] = \Phi_A[M(w)] = \phi_A(w),
$$

with $\phi_A$ the $A$-optimality criterion. Denoting $\Omega(w) = I_m \otimes M(w) = \sum_{i=1}^q w_i (I_m \otimes H_i)$, we get

$$
\phi_A(w) = \mathbf{c}^\top \Omega^{-1}(w) \mathbf{c},
$$

(4.1)

showing that the screening bounds of Theorems 2.1 and 3.2 can also be used to eliminate inessential $H_i$ in $A$-optimal design. We use $\Omega^{-1}$ instead of $\Omega$ in (4.1) since the property that $\mathbf{c}$ belongs to the column space of the optimal matrix $\Omega$ implies that the corresponding $\Omega_s$ is nonsingular, with $\phi_A(w) = \infty$ for a singular $M(w)$. We thus obtain the following result, where we use the property that, for any $m \times m$ matrix $A$, $I_m \otimes A$ has the same eigenvalues as $A$ (with multiplicities $m$).

**Corollary 4.1.** Let $M$ be any full rank matrix in $\mathcal{X}$ and $H_i$ be a matrix in $\mathcal{H}$ such that

$$
B_{1,A}(M, H_i) = \left[1 - \Delta_i \left(\frac{\delta}{1+\delta}\right)^{1/2}\right] \Phi_A(M) - \text{trace}(M^{-2}H_i) > 0,
$$

(4.2)

or

$$
B_{2,A}(M, H_i) = \gamma(\kappa_i, \varphi) \Phi_A(M) - \text{trace}(M^{-2}H_i) > 0,
$$

or

$$
B_{3,A}(M, H_i) = 1 + \sup_{\beta \geq \lambda_{\max}(A_i^\top M^{-1}A_i)} \left\{ \frac{\Phi_A(M)}{(1+\delta) \text{trace}[(\beta M - H_i)^{-1}]} - \beta \right\} > 0,
$$

(4.3)

where

$$
\delta = \delta(M) = \max_{i=1,\ldots,q} \frac{\text{trace}(M^{-2}H_i)}{\Phi_A(M)} - 1,
$$

(4.4)

and $\Delta_i$ and $\gamma(\kappa_i, \varphi)$ are defined in Theorem 2.1. Then $H_i$ does not support an $A$-optimal design.

The case $H_i = a_i a_i^\top$ deserves special attention here too. We cannot directly use the results of Section 3.3 since the matrix $\Omega_i = I_m \otimes H_i = (I_m \otimes a_i)(I_m \otimes a_i^\top)$ is not of rank 1. However, similarly to Section 3.3, simplifications are obtained for $B_{1,A}$ and $B_{3,A}$ since $\Delta_i = \lambda_{\max}(A_i^\top M^{-1}A_i) = a_i^\top M^{-1}a_i$ and the test for $B_{3}(M,a_i a_i^\top) > 0$ can be made explicit. Again, $B_{1,A}(M,a_i a_i^\top) > 0$ implies $B_{3,A}(M,a_i a_i^\top) > 0$, indicating that $B_{3,A}$ is more efficient than $B_{1,A}$ for $A$-optimal design with $H_i = a_i a_i^\top$.

**Theorem 4.1.** For any full rank matrix $M$ in $\mathcal{X}$ and $H_i = a_i a_i^\top$ in $\mathcal{H}$, $B_{3,A}(M,a_i a_i^\top)$ defined by (4.3) satisfies

$$
B_{3,A}(M,a_i a_i^\top) > 0 \iff a_i^\top M^{-1}a_i \in [0,1) \cup \left(\delta_i^{(3,A)}, \Delta_i^{(3,A)}\right),
$$

(4.5)
where

\[
\delta^{(3,A)}_i = \frac{(1 + \delta) a_i^\top M^{-2} a_i}{\Phi_A(M)}, \quad \Delta^{(3,A)}_i = 1 + \frac{1}{\delta} \left[ 1 - (\delta^{(3,A)})^{1/2} \right]^2.
\]

Moreover, if \( H_i \) satisfies (4.2), then it also satisfies (4.3).

Proof. Using the Sherman-Morrison-Woodbury formula for matrix inversion, we obtain that the function \( f(\beta) \) involved in the computation of \( B_{3,A} \) is

\[
f(\beta) = \beta - \frac{\beta \Phi_A(M)}{(1 + \delta) \left[ \Phi_A(M) + \frac{a_i^\top M^{-2} a_i}{\beta a_i^\top M^{-1} a_i} \right]},
\]

that is, an expression similar to (3.4) with \( a_i^\top M^{-2} a_i \) substituted for \( (c^\top M^{-1} a_i)^2 \). The proof of (4.5) is thus similar to that of (3.3), with now \( \beta^* = \tau \left[ \tau + a_i^\top M^{-2} a_i \right] / \left[ \delta \Phi_A(M) a_i^\top M^{-2} a_i \right] \) where \( \tau = (\delta a_i^\top M^{-2} a_i)^{1/2} \left[ \Phi_A(M) (a_i^\top M^{-1} a_i) - a_i^\top M^{-2} a_i \right]^{1/2} \) and \( a_i^\top M^{-2} a_i \leq (a_i^\top M^{-1} a_i) \lambda_{\text{max}}(M^{-1}) < (a_i^\top M^{-1} a_i) \Phi_A(M) \). The condition \( B_1(M, a_i a_i^\top) > 0 \) is equivalent to

\[
\Delta_i < (1 + 1/\delta)^{1/2} \left[ 1 - (a_i^\top M^{-2} a_i) / \Phi_A(M) \right].
\]

The rest of the proof is identical to that of Theorem 3.4, with \( a_i^\top M^{-2} a_i \) substituted for \( (c^\top M^{-1} a_i)^2 \).

\[\tag*{\square}\]

Remark 4.1. More generally, following the same line as in Sagnol (2011, Section 4.1), similar results can be derived for \( L_r \) (linear) optimality — sometimes also called \( A_K \)-optimality by some authors — where the criterion to be minimized is \( \Phi_L[M(w)] = \text{trace}(C M^{-r}(w)) \), with \( C \) symmetric nonnegative definite. Indeed, for \( C = K K^\top \), with \( K \) an \( m \times r \) matrix, \( r \leq m \), we can write \( \text{trace}(C M^{-r}) = \text{trace}(K^\top M^{-r} K) = \bar{c}^\top M^{-1} \bar{c} \), with \( \bar{c} = \text{vec}(K) \) and \( M = I_r \otimes M \), so that \( L \)-optimal design corresponds to a particular case of \( c \)-optimal design in a higher dimensional space. In particular, when \( H_i = a_i a_i^\top \), the function \( f(\beta) \) in the computation of \( B_3 \) is

\[
f(\beta) = \beta - \frac{\beta \Phi_L(M)}{(1 + \delta) \left[ \Phi_L(M) + \frac{a_i^\top M^{-1} C M^{-1} a_i}{\beta a_i^\top M^{-1} a_i} \right]},
\]

and the result of Theorem 4.1 remains valid when \( \Phi_A(M) \) is replaced by \( \Phi_L(M) \) and \( a_i^\top M^{-2} a_i \) by \( a_i^\top M^{-1} C M^{-1} a_i \) in \( \delta^{(3,A)} \).

4.2 Comparison between screening bounds

In (Pronzato, 2013) a screening rule is proposed for all criteria of the form \( \Phi_p(M) = \text{trace}(M^{-p}) \), for \( p \in (-1, \infty) \). The specialization of this result to \( A \)-optimal design, which corresponds to \( p = 1 \), writes as follows.

Theorem 4.2 (Pronzato, 2013). Let \( M \) be any full rank matrix in \( \mathcal{K} \) and \( H_i \) be a matrix in \( \mathbb{H} \) such that

\[
B_{4,A}(M, H_i) = \omega^2 \frac{\Phi_A(M)}{1 + \delta} - \text{trace}(M^{-2} H_i) > 0,
\]

where \( \omega \) is the unique root in the interval \((\alpha^{1/2}, 1)\) of the fourth degree polynomial \( P(x) = (\alpha - x^2)(1 + \delta - \alpha x)^2 + (1 - \alpha)^3 x^2 \), with \( \delta \) given by (4.4) and \( \alpha = \lambda_{\text{min}}(M^{-1}) / \Phi_A(M) < 1 \). Then \( H_i \) does not support an \( A \)-optimal design.
Besides the resolution of a fourth-order polynomial equation, the calculation of $B_{4,A}(M,H_i)$ requires the computation of the minimum eigenvalue of $M^{-1}$ (or equivalently of the maximum eigenvalue of $M$).

The comparison of $B_{4,A}$ with $B_{1,A}$ and $B_{3,A}$ is complicated due to the absence of an explicit expression for $\omega$ and its dependence on $\lambda_{\min}(M^{-1})$.

Consider the particular situation where $M = \alpha M_\epsilon$ for some $\alpha \in (0,1)$, with $M_\epsilon$ the $A$-optimal matrix in $\mathcal{M}$. Denote $\gamma_i = \text{trace}(M^{-2}_\epsilon H_i)/\phi^*_A$, where $\phi^*_A = \Phi_A(M_\epsilon)$ and $\gamma_i \in (0,1]$ for all $i$ from the optimality of $M_\epsilon$. Direct calculation gives $\text{trace}(M^{-2} H_i) = (\gamma_i/\alpha^2) \phi^*_A$, $\delta = 1/\alpha - 1$ and $\text{trace}(M^{-1} H_i) = \text{trace}(M^{-1} H_i)/\alpha > \text{trace}(M^{-2} H_i)/\Phi_A(M) = \gamma_i/\alpha$, with $\alpha = \phi^*_A/\Phi_A(M)$ the $A$-efficiency of $M$.

For $H_i = a_i a_i^\top$, denote $t_i^* = \text{trace}(M^{-1}_\epsilon H_i) = a_i^\top M^{-1}_\epsilon a_i$. The definition (4.2) of $B_{1,A}$ gives

$$B_{1,A}(M,a_i a_i^\top) = \left[1 - t_i^* \frac{(1-\alpha)^{1/2}}{\alpha} - \frac{\gamma_i}{\alpha}\right] \Phi_A(M) \leq \left\{1 - \frac{\gamma_i}{\alpha} \left[1 + (1-\alpha)^{1/2}\right]\right\} \Phi_A(M).$$

Therefore, $B_{1,A}(M,a_i a_i^\top) > 0$ is equivalent to $t_i^* < t_1(\alpha, \gamma_i) = (1 - \gamma_i/\alpha)/(1 - \alpha)^{1/2}$ and implies $\gamma_i < \gamma_1(\alpha) = \alpha [1 + (1-\alpha)^{1/2}]^{-1}$. From $\omega < 1$ we obtain that $B_{4,A}(M,a_i a_i^\top) > 0$ implies $\gamma_i < \alpha^2$. If we compare with the necessary condition $\gamma_i < \gamma_1(\alpha)$ for $B_{1,A}(M,a_i a_i^\top) > 0$, we see that elimination with $B_{1,A}$ is more demanding than with $B_{1,A}$ for small $\alpha$, but less demanding for $\alpha > \phi$, the Golden-Section number $(\sqrt{5} - 1)/2 \approx 0.6180$, that is, when $M$ is close to being optimal.

To make the comparison more precise, we specialize the example to the case $M_\epsilon = I_m$, with therefore $a_i^\top M^{-2}_\epsilon a_i = \|a_i\|^2 \leq \phi^*_A = m$ and $M = \alpha I_m$, $\alpha \in (0,1)$. It may correspond for instance to $A$-optimal design in the hypercube $[-1,1]^m$ but also covers other situations. We obtain that $B_{1,A}(M,a_i a_i^\top) > 0$ is equivalent to $\|a_i\|^2 < C_1(\alpha, m) = m \alpha/[1 + (m - 1)/(1 - \alpha)^{1/2}]$, and $B_{3,A}(M,a_i a_i^\top) > 0$ is equivalent to $\|a_i\|^2 < C_3(\alpha, m) = m/[1 + (m - 1)/(1 - \alpha)^{1/2}]^{1/2}$ for $\alpha \in (1/m, 1)$ (and $C_3(\alpha, m) \geq \alpha$, with $C_3(1/m, m) = 1/m$). Moreover, $B_{4,A}(M,a_i a_i^\top) = m \omega^2 - \|a_i\|^2/\alpha$, with $\omega = \omega(m, \alpha)$ defined in Theorem 4.2, and $B_{4,A}(M,a_i a_i^\top) > 0$ is equivalent to $\|a_i\|^2 < C_4(\alpha, m) = m \alpha^2 \omega^2(m, \alpha)$.

The left panel of Fig. 2 shows the bounds $C_1(\alpha, m)$, $C_3(\alpha, m)$ and $C_4(\alpha, m)$ as functions of $\alpha$ for $m = 3$, $C_1(\alpha, m) > C_3(\alpha, m)$ for $\alpha \in (0,1)$ and any $m \geq 2$ since $B_{3,A}$ is always more efficient than $B_{1,A}$, see Theorem 4.1. On the one hand, $C_4(\alpha, m) < C_1(\alpha, m)$ on the left (and major) part of the interval $[0,1]$, which makes $B_{4,A}$ the less powerful for those $\alpha$. On the other hand, for any $m \geq 2$ there exist $\alpha_4^*(m) < \alpha_3^*(m)$ in $(0,1)$ such that $C_4(\alpha, m) > C_1(\alpha, m)$ for $\alpha > \alpha_4^*(m)$ and $C_4(\alpha, m) > C_3(\alpha, m)$ for $\alpha > \alpha_3^*(m)$, which makes $B_{4,A}$ the most powerful close to optimality. The right panel of Fig. 2 shows the $A$-efficiency levels $\alpha_4^*(m)$ and $\alpha_3^*(m)$ above which $B_{4,A}$ is more powerful than $B_{1,A}$ and $B_{3,A}$, respectively, for $m = 2$ to 25. The numerical comparison in Section 5.3 indicates, however, that this behavior is not generic, $B_{4,A}$ being there much less powerful than $B_{1,A}$ and $B_{3,A}$.

5 Examples

We use rank-one matrices $H_i = a_i a_i^\top$ in Section 5.3 on $A$-optimal design, but full rank matrices $H_i = a_i a_i^\top + \lambda I_m$, $\lambda > 0$, in the examples of Sections 5.1 and 5.2 which concern $c$-optimal design. When $H_i = a_i a_i^\top$ for all $i$, $c$-optimal design corresponds to a linear programming problem (Harman and Jurik, 2008), and elimination of inessential $H_i$ is less a problem than in the more difficult situation where rank($H_i$) > 1.
Figure 2: Left: $C_1(\alpha, m)$ (red ★), $C_3(\alpha, m)$ (magenta +) and $C_4(\alpha, m)$ (black ♦), such that $H_i$ is eliminated by $B_{3,A}$ when $||a_i||^2 < C_j$, as functions of the $A$-efficiency $\alpha \in (0, 1)$ for $m = 3$. Right: $A$-efficiencies $\alpha_1^*(m)$ (red ★) and $\alpha_3^*(m)$ (magenta +) above which $B_{4,A}$ is more powerful than $B_{1,A}$ and $B_{3,A}$ for $m = 2, \ldots, 25$.

To illustrate the potential interest of the elimination of inessential $H_i$ in the construction of an optimal design, we consider the multiplicative algorithm of Fellman (1974), defined by

$$w_i^{k+1} = \frac{w_i^k \left[c^\top M^{-1}(w)H_iM^{-1}(w)c\right]^{1/2}}{\sum_{j=1}^q w_i^k \left[c^\top M^{-1}(w)H_jM^{-1}(w)c\right]^{1/2}}$$

(5.1)

for $c$-optimality, and

$$w_i^{k+1} = \frac{w_i^k \text{trace}^{1/2} [M^{-2}(w)H_i]}{\sum_{j=1}^q w_i^k \text{trace}^{1/2} [M^{-2}(w)H_j]}$$

(5.2)

for $A$-optimality; see also Yu (2010).

5.1 Example 1: quadratic regression

In this example $H_i = a_i a_i^\top + \lambda I_2$ with $a_i = a(t_i) = (t_i, t_i^2)^\top$ ($m = 2$) and $t_i \in [0, 1]$ for all $i$. We take $c = (1, \cdot)^\top$. When $\lambda = 0$, the $c$-optimal design over the compact set of matrices $H_i$ with $t \in [0, 1]$ is the probability measure

$$\xi_0^* = \begin{cases} \alpha^* \delta_{\sqrt{2} - 1} + (1 - \alpha^*) \delta_1 & \text{if } c \in [0, \sqrt{2} - 1] \cup [1, \infty), \\ \delta_c & \text{otherwise}, \end{cases}$$

where $\alpha_0^* = (\sqrt{2}/2)(1 - c)/[2(\sqrt{2} - 1) - c]$ and $\delta_t$ denotes the Dirac delta measure at $t$. In the rest of the example, we take $c = (\sqrt{2} - 1)/2$, which gives a 2-point optimal design with $\alpha_0^* = (\sqrt{2}/2)(1 - \sqrt{2}/3)/(\sqrt{2} - 1) \approx 0.902369$.

Direct calculation based on the optimality condition (2.3) shows that for $\lambda \geq \lambda_0 = (1/2)(15 + 16 \sqrt{2})/(18 + 11 \sqrt{2})$, the optimal design is $\xi_0^* = \delta_1$. For $\lambda_0 \geq \lambda \geq \lambda_1 = (\sqrt{2} - 1)^2/(4 \sqrt{2})$, $\xi_0^* = \delta_t$, with $t = 1$ for $\lambda = \lambda_0$, $t = \sqrt{2} - 1$ for $\lambda = \lambda_1$, and $t$ and $\lambda$ related by $\lambda = t^3(\sqrt{2} + 1)(2t + 1 - \sqrt{2})/[2(t + 1 + \sqrt{2})]$ in between. For $\lambda < \lambda_1$, $\xi_0^*$ is supported on the two points $\sqrt{2} - 1$ and 1, like
\[ \xi^*, \text{ and } \alpha^* \text{ can be computed numerically, with } \alpha^* \simeq 0.980081, \alpha^* \simeq 0.910140 \text{ and } \alpha^* \simeq 0.902377 \text{ for } \lambda = 10^{-2}, \lambda = 10^{-3} \text{ and } \lambda = 10^{-6}, \text{ respectively.} \]

First, we test and compare the capability of the inequalities \( B_j(M(w), H_i) > 0, \) for \( j = 1, 2, 3, \) to eliminate inessential \( H_i \) on a set of \( N \) random designs. These designs have a fixed support \( H \) of size \( q, \) corresponding to \( t_k = \sqrt{2} - 1 \) and \( t_k = (i - 2)/(q - 2) \) for \( i = 2, \ldots, q; \) the associated weights are \( w_i = (\gamma_i + \xi_i) / [\sum_{i=1}^{q}(\gamma_i + \xi_i)], \) where \( \gamma_1 = \alpha^*, \gamma_q = (1 - \alpha^*), \gamma_i = 0 \) for \( i \neq 1, q, \) and where the \( \zeta \) are i.i.d. in \([0, s] \). For \( \lambda \) small enough, the designs are close to optimality when \( s \) is close to 0. For each design generated, characterized by a vector of weights \( w, \) we compute how many \( H_i \) are eliminated by \( B_j(M(w), H_i) > 0, j = 1, 2, 3. \)

Figure 3 present box plots, constructed from the \( N \) random design, of the proportions of inessential \( H_i \) detected when using each test \( B_j(M, H_i) > 0 \) separately \((T_j)\), when using pairs of tests jointly \((T_j \cup T_k)\), and when using the three tests altogether \((T_1 \cup T_2 \cup T_3)\). We take \( q = 500, s = 10^{-3}, N = 1000 \) and \( \lambda = 10^{-3} \) (left), \( \lambda = 10^{-6} \) (right); \( \epsilon = 0.01 \) in the dichotomy procedure to compute \( B_3. \) As anticipated in Remark 2.1-(i), \( T_2 \) is rather inefficient when \( \lambda \) is small, a typical situation when \( \lambda \) is just a regularization term introduced to avoid singularity of the \( H_i. \) T1 is more powerful than T2 for the two values of \( \lambda \) considered, but not uniformly better for \( \lambda = 10^{-3} \) since the combination of T1 and T2 eliminates more points than T1 alone. Although the \( H_i \) have full rank, so that Theorem 3.4 does not apply, T3 always performs best on this example and combination with another bound is not helpful.

![Figure 3: Box plots for 1000 random design \((q = 500, s = 10^{-3})\) of proportions of inessential \( H_i \) when using each test separately \((T_j)\), by pairs \((T_j \cup T_k)\), and altogether \((T_1 \cup T_2 \cup T_3)\), for \( \lambda = 10^{-3} \) (left) and \( \lambda = 10^{-6} \) (right).](image)

The left panel of Fig. 4 shows \( B_1(M, H(t)) \), \( B_2(M, H(t)) \) and \( 10 B_3(M, H(t)) \) as functions of \( t \in [0, 1], \) for \( H(t) = a(t)a^+(t) + 10^{-3}I_2 \) and \( M = M(w^{1000}) \) the matrix obtained after 1000 iterations of the multiplicative algorithm (5.1), initialized at the uniform measure on \( H \) that gives weight \( 1/q \) at each of the \( q \) matrices \( H_i = H(t_i) \) \((q = 500)\). All \( H_i \) such that one of the \( B_j(M, H_i) \) is positive can be safely removed from \( H. \) The right panel plots the proportions \( \rho_j(k) \) of \( H_i \) such that \( B_j(M(w^k), H_i) > 0 \) as a function of the iteration number \( k, \) for \( k = 1, \ldots, 1000). \)

In the rest of the example, we investigate the potential interest of eliminating inessential points in terms of computational cost \((\text{time})\). We take \( q = 500, \lambda = 10^{-3} \) and \( \epsilon = 0.01 \) for \( B_3. \)

The left panel of Fig. 5 shows the evolution of \( \delta \) given by (2.1) as a function of computational time \( T \) for the multiplicative algorithm (5.1), initialized as above, when no \( H_i \) is removed \((\text{black squares})\) and when inessential \( H_i \) are eliminated at iteration \( k \) if \( B_j(M(w^k), H_i) > 0. \) \( T \) is measured
in seconds (calculations with Matlab, on a PC with a clock speed of 2.50 GHz and 32 Go RAM), but only the comparison between curves in meaningful. It indicates that, for the multiplicative algorithm, it is more efficient not to perform any screening test than to try to eliminate points at every iteration with one of the screening test proposed in the paper. In particular, $B_3(M,H_i)$ has a non negligible computational cost due to dichotomy search. If we consider that the computational cost of one iteration (5.1) is roughly proportional to the size of $w$, $q(1 - \rho_j(i))$ at iteration $i$ for elimination by $T_j$, we may rescale the iteration counter $k$ into $\sum_{i=1}^k (1 - \rho_j(i))$ which counts pseudo iterations accounting for the decreasing size of $w^k$. The right panel of Fig. 5 provides such a comparison of computational costs that neglects the costs of the screening tests, and only shows performance in terms of proportions of points eliminated; this is in agreement with Fig. 4-right.

To reduce the total computational cost, we now perform the screening tests more rarely, every 500 iterations only. Figure 6 presents the same information as Fig. 5 in this periodic screening situation. The left panel shows that the test based on $B_1(M,H_i)$ is then the most powerful on this example and provides a significant acceleration of the algorithm. The curves on the right panel are confounded up to iteration $k = 500$ when the first test is performed, compare with Fig. 5.

We have stopped the algorithm after 10 000 iterations. Figure 6 indicates that removing inessential points can accelerate the algorithm by lightening the iterations, but does not reduce the number of iterations required to reach a given accuracy, as the value of $\delta$ reached after the 10 000 iterations is approximately always the same.

5.2 Example 2: uniform points in the ball $\mathcal{B}_m(0,1)$

We take $H_i = a_i a_i^\top + \lambda I_m$ with the $a_i$ in the $m$-dimensional unit ball $\mathcal{B}_m(0,1)$. When the design space is the whole ball, the delta measure at $a^* = c/\|c\|$ is $c$-optimal, with $\Phi(M_*) = \|c\|^2/(\lambda + 1)$ and, for any $a_i \in \mathcal{B}_m(0,1)$, $\sum_{i=1}^k (1 - \rho_j(i))$ which counts pseudo iterations accounting for the decreasing size of $w^k$. The right panel of Fig. 5 provides such a comparison of computational costs that neglects the costs of the screening tests, and only shows performance in terms of proportions of points eliminated; this is in agreement with Fig. 4-right.

To reduce the total computational cost, we now perform the screening tests more rarely, every 500 iterations only. Figure 6 presents the same information as Fig. 5 in this periodic screening situation. The left panel shows that the test based on $B_1(M,H_i)$ is then the most powerful on this example and provides a significant acceleration of the algorithm. The curves on the right panel are confounded up to iteration $k = 500$ when the first test is performed, compare with Fig. 5.

We have stopped the algorithm after 10 000 iterations. Figure 6 indicates that removing inessential points can accelerate the algorithm by lightening the iterations, but does not reduce the number of iterations required to reach a given accuracy, as the value of $\delta$ reached after the 10 000 iterations is approximately always the same.

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We first visualize, for the case $m = 2$, regions of the disk where design points are eliminated.
Figure 5: Left: Evolution of $\delta$ given by (2.1) as a function of computational time $T$ for 10,000 iterations of algorithm (5.1) with and without elimination of inessential $H_i$. Right: $\delta$ as a function of pseudo iteration number. No elimination: black $\square$, elimination by $T_1$ (red $\bigstar$), $T_2$ (blue $\lhd$), or $T_3$ (magenta $+$), with screening at every iteration.

Figure 6: Same as Fig. 5 but with a screening test every 500 iterations only.

Figure 7 shows the outside contours of those regions, defined by the curves $A_j = \{a_i \in B_2(0,1) : B_j(M(w), H_i) = 0\}$, for $\lambda = 1$ (left) and $\lambda = 0.1$ (right). The matrix $M(w)$ corresponds to a $w$ close to optimality, obtained with a similar approach to previous example: we set $a_1 = c$ and generate $q-1$ other random $a_i$ independently and uniformly distributed in $B_m(0,1)$; the $q$ associated weights are $w_i = (\gamma_i + \zeta_i)/[\sum_{i=1}^q (\gamma_i + \zeta_i)]$, where $\gamma_1 = 1$ and $\gamma_i = 0$ for $i > 1$, and where the $\zeta_i$ are i.i.d. in $[0, s]$. We have taken $q = 10^5$ and $s = 10^{-5}$, and the designs we use are such that $\Phi[M(w)]/\Phi(M_s) \simeq 1.14$ and 1.29 for $\lambda = 1$ and $\lambda = 0.1$, respectively. On the left panel, the regions for $T_2$ and $T_3$ are hardly distinguishable and the two tests dominate $T_1$; on the right panel, $T_3$ dominates $T_1$ and $T_2$, but none of them dominates the other. The complicated shapes of the contours, even in such a simple example, indicates the difficulty of determining which of the three tests is best in a given situation.

We take $\lambda = 0.1$ through the rest of the example. Figure 8 presents box plots of the proportions
of inessential points detected when using each screening test separately ($T_j$), when using pairs of tests jointly ($T_j U T_k$), and when using the three tests altogether ($T_1 U T_2 U T_3$). The plots are obtained from $N$ random designs generated as above, with $m = 5$, $q = 1000$, $s = 10^{-4}$ and $N = 1000$. Here $T_1$ is the less powerful, while $T_2$ and $T_3$ perform almost equally. There is no clear benefit in using several tests jointly.

As in Example 1, we investigate the effect of point elimination on computational time: we perform 2000 iterations of the multiplicative algorithm, initialized at the uniform design on a set of $q$ points independently and uniformly distributed in $B_m(0,1)$, with $m = 50$ and $q = 1000$. Figure 9 presents the same information as Fig. 6, but when screening tests are performed more frequently, every 100 iterations. The staircase shape of the curves on the left panel is due to the long computational time required by the evaluation of $B_j(M(w^k), H_i)$ for all $H_i$ not eliminated yet:
there are 20 steps along each of the two curves, one every 100 iterations, shorter and shorter as more $H_i$ are discarded. The left panel shows that $T_3$ is hardly competitive due to its high computational cost; it becomes the most powerful if we omit this cost, see the right panel. $T_2$ is more powerful than $T_1$ and quicker to compute than $T_3$; it is the most efficient in this example.

Figure 9: Left: Evolution of $\delta$ given by (2.1) as function of computational time $T$ for algorithm (5.1) with and without elimination of inessential $H_i$. Right: $\delta$ as a function of pseudo iteration number. No elimination: black □, elimination by $T_1$ (red ★), $T_2$ (blue ▽), or $T_3$ (magenta +), with screening every 100 iterations.

5.3 Example 3: $A$-optimal design

We consider $A$-optimal design for a product-type regression model where $a_i = a(x_i), \ x_i \in \mathcal{X} = [-1,1]^2$ and, for $x = (a,b), a(x) = (1, a, a^2, b, b^2, ab, a^2 b^2, a b^2)^T \in \mathbb{R}^9$. The $A$-optimal design is the cross product of two $A$-optimal designs for quadratic regression on $[-1,1]$, given by $(1/4) \delta_{-1} + (1/2) \delta_0 + (1/4) \delta_1$, with $\text{trace}(M^{-1}) = 64$. To compute optimal designs, we discretize $\mathcal{X}$ into a regular grid with $201 \times 201$ points, which yields a set $H$ of $q = 40401$ rank-one matrices $H_i = a_i a_i^T$.

The left panel of Fig. 10 shows the evolution of $\delta$ given by (2.1) as a function of computational time $T$ for the multiplicative algorithm (5.2), initialized with the uniform distribution on $H$, when no $H_i$ is removed (black squares) and when inessential $H_i$ are eliminated at iteration $k$ if $B_{j,A}(M(w^k), H_i) > 0$, for $j = 1, 3, 4$. Although the screening tests are performed at every iteration, we see that elimination of inessential points with $T_1$ and $T_3$ yields a significant acceleration of the algorithm. The higher computational cost of $B_{4,A}$ (it requires the computation of the minimum eigenvalue of a $m \times m$ matrix, with here $m = 9$) is not the only reason for the poorer performance of $T_4$: the evolution of $\delta$ as a function of the pseudo iteration counter $\sum_{i=1}^k (1 - \rho_j(i))$ on the right panel indicates that $T_4$ is much less powerful than $T_1$ and $T_3$ on this example, a situation opposite to that in Section 4.2. The superiority of $T_3$ over $T_1$ is no due to a smaller computational cost (see the right panel), but to its more efficient elimination; see Theorem 4.1. When $H_i = a_i a_i^T$, the easy calculation of $B_{3,A}$ and its good performance in various situations make it the recommended method for support point elimination in $A$-optimal design.
Figure 10: Left: Evolution of $\delta$ given by (2.1) as a function of computational time $T$ for 1,000 iterations of algorithm (5.2) with and without elimination of inessential $H_i$. Right: $\delta$ as a function of pseudo iteration number. No elimination: black □, elimination by $T_1$ (red ⭐), $T_3$ (magenta +) or $T_4$ (blue ◊), with screening at every iteration.

6 Conclusion and further developments

We have derived three inequalities that must be satisfied by any elementary information matrix $H_i$ that contributes to a c-optimal design. The first two, $B_1$ and $B_2$, are direct consequences of the equivalence theorem for c-optimality; the third one $B_3$ is derived from the dual problem. Each of those inequalities can be used to safely eliminate, from a set $H$ of $q$ candidates, inessential $H_i$ that cannot be support points of an optimal design. When implemented within an optimization algorithm, these inequalities can thereby provide a substantial acceleration by reducing the size of the problem online. Screening tests do not need to be done every iteration, and examples have been presented showing the value of periodic screening.

In the general situation where $\text{rank}(H_i) > 1$, elimination by $B_3$ necessitates an iterative optimization, which induces a non-negligible computational cost and makes it less useful than $B_1$ and $B_2$. When $\text{rank}(H_i) = m$, both $B_1$ and $B_2$ can be used, and the major contribution to their computational cost comes from the calculation of the maximum and minimum eigenvalues of a $m \times m$ matrix. $B_2$ is less efficient than $B_1$ when $H_i$ is badly conditioned, but may perform better otherwise; since there is a priori no strict superiority of one bound compared to the other, we recommend to use both in parallel. When $1 < \text{rank}(H_i) < m$, $B_2$ cannot be used and we recommend to use $B_1$ due to its much easier calculation than $B_3$. The situation is somewhat reverse when $\text{rank}(H_i) = 1$: $B_3$ can be calculated explicitly and Theorem 3.4 shows that it is more efficient than $B_1$.

Since A-optimal design can be seen as a particular case of c-optimal design in a higher dimensional space, these results extend directly to A-optimal design, and the same recommendations as above can be made. In particular, when $\text{rank}(H_i) = 1$, $B_3$ can still be calculated explicitly and Theorem 4.1 indicates that it is more efficient than $B_1$. Numerical examples show that in general both $B_1$ and $B_3$ outperform a screening bound already proposed for Kiefer’s $\varphi_p$ class of criteria, although an example has been provided indicating that the situation is reverse in particular situations.

The fact that concurrent inequalities have been derived suggests that there may still be room for improvement. The equivalence between c-optimal design and a lasso problem (Sagnol and Pauwels, 2019) is also a strong motivation for trying to derive other screening tests with low computational
cost, especially for situations with large $m$ and $q$.

Appendix

A Proof of Theorem 2.1

The proof relies on the following property.

**Lemma A.1.** Let $\Omega$ be a $m \times m$ symmetric non-negative definite matrix with minimum and maximum eigenvalues $\lambda_a$ and $\lambda_b$, respectively. Then for any pair of unit vectors $u$, $v$ in $\mathbb{R}^m$ making an angle $\varphi \in (-\pi/2, \pi/2)$, we have

$$\min_{u,v} u^T \Omega u - v^T \Omega v = - (\lambda_b - \lambda_a) |\sin(\varphi)|, \quad (A.1)$$

and

$$\min_{u,v} \frac{u^T \Omega u}{v^T \Omega v} = \frac{\lambda_a \cos^2(\omega - \varphi) + \lambda_b \sin^2(\omega - \varphi)}{\lambda_a \cos^2(\omega) + \lambda_b \sin^2(\omega)}, \quad (A.2)$$

where

$$\omega = \frac{1}{2} \left\{ \arccos \left[ \frac{\lambda_b - \lambda_a}{\lambda_b + \lambda_a} \cos(\varphi) \right] + \varphi \right\}. \quad (A.3)$$

The minimum is attained for $u$ and $v$ lying in the plane spanned by eigenvectors $e_a$, $e_b$ respectively associated with $\lambda_a$ and $\lambda_b$. In that plane, there exist four configurations that yield the minimum in (A.1), respectively in (A.2), obtained by changing the sign of one coordinate, or of both coordinates, of $u$ and $v$ (simultaneously) in the basis formed by $e_a$ and $e_b$. For (A.1), one of these configurations corresponds to $u$ (respectively $v$) making an angle of $\pi/4 - |\varphi|/2$ (respectively, $\pi/4 + |\varphi|/2$) with $e_a$. For (A.2), one configuration corresponds to $u$ (respectively $v$) making an angle of $\omega - \varphi$ (respectively, $\omega$) with $e_a$. When $\lambda_a = 0$, $\omega = \varphi$ and $\min_{u,v}(u^T \Omega u)/(v^T \Omega v) = 0$.

**Proof.** (i) Denote $f(u, v) = u^T \Omega u - v^T \Omega v$. We assume that $\lambda_b > \lambda_a$ to avoid the trivial solution $f(u, v) = 0$ for all unit vectors $u$ and $v$. The minimum of $f$ cannot be achieved for $u$ and $v$ belonging to the null space of $\Omega$, since otherwise $u^T \Omega u = v^T \Omega v = 0$. Therefore, suppose that $\text{span}(u, v)$ intersects the column space of $\Omega$, and denote by $0 \leq \lambda_1 < \lambda_2$ the eigenvalues of $\Pi_{u,v} \Omega \Pi_{u,v}$, with $\Pi_{u,v}$ the orthogonal projector onto the plane $(u, v)$, and by $e_1$ and $e_2$ the canonical orthonormal basis defined from the associated eigenvectors. In this coordinate system, we can write $u = [\cos(\alpha), \sin(\alpha)]^T$ and $v = [\cos(\beta), \sin(\beta)]^T$, which gives $f(u, v) = \lambda_1 [\cos^2(\alpha) - \cos^2(\beta)] + \lambda_2 [\sin^2(\alpha) - \sin^2(\beta)] = (\lambda_2 - \lambda_1) \sin(\alpha + \beta) \sin(\alpha - \beta)$ and $u^T v = \cos(\alpha - \beta) = \cos(\varphi)$. The minimum of $f$ for $u, v$ in this plane is thus equal to $-(\lambda_2 - \lambda_1) |\sin(\varphi)|$ and is attained when $u$ and $v$ are in one of the configurations indicated in the lemma. The overall minimum of $f$ is achieved when $u$ and $v$ lie in the plane defined by eigenvectors $e_a$ and $e_b$ associated with $\lambda_a$ and $\lambda_b$.

(ii) Denote now $g(u, v) = (u^T \Omega u)/(v^T \Omega v)$. When $\Omega$ is singular ($\lambda_a = 0$), $g$ reaches its minimum value zero when $u$ is in the null space of $\Omega$. We thus suppose that $\lambda_a > 0$ and proceed as above by writing $u = [\cos(\alpha), \sin(\alpha)]^T$ and $v = [\cos(\beta), \sin(\beta)]^T$ in the coordinate system defined by $e_1$ and $e_2$. Since $u^T v = \cos(\alpha - \beta) = \cos(\varphi)$, we obtain

$$g(u, v) = \frac{\lambda_a \cos^2(\beta - \varphi) + \lambda_b \sin^2(\beta - \varphi)}{\lambda_a \cos^2(\beta) + \lambda_b \sin^2(\beta)}.$$
Direct calculation shows that the derivative with respect to $\beta$ is proportional to $(\lambda_1 + \lambda_2) \cos(2\beta - \varphi) - (\lambda_2 - \lambda_1) \cos(\varphi)$, and the minimum $g_*$ of $g(u, v)$ is achieved at $\beta_*$ obtained by substituting $\lambda_1$ and $\lambda_2$ for $\lambda_0$ and $\lambda_0$ in (A.3). Finally, $\beta_*$ and $g_*$ only depend on $\lambda_1$ and $\lambda_2$ through their ratio $\kappa = \lambda_2/\lambda_1$, with $g_*$ being a decreasing function of $\kappa$. The overall minimum of $g$ is thus achieved when $u$ and $v$ lie in the plane defined by $e_a$ and $e_b$.

Proof. (of Theorem 2.1) We show that any $H_i$ such that $w_i^* > 0$ for some c-optimal design $w^*$ satisfies

(i) $c^\top M^{-1}H_iM^{-1}c \geq \{1 - \Delta_i[\delta/(1 + \delta)]^{1/2}\} \Phi(M)$

and

(ii) $c^\top M^{-1}H_iM^{-1}c \geq \gamma(\kappa_i, \varphi) \Phi(M)$,

where $\Delta_i$, $\kappa_i$, $\varphi$ and $\gamma(\kappa, \varphi)$ are defined in the theorem.

Let $M_* = M(w^*)$ be a c-optimal matrix in $\mathcal{H}$, and denote by $M_*^{-}$ its generalized inverse in the equivalence theorem, see Section 2.1. Denote $u = M^{-1/2}c$, $v = M^{1/2}M_*^{-}c$, $\hat{u} = u/\|u\|$, $\hat{v} = v/\|v\|$ and $\Omega_i = M^{-1/2}H_iM^{-1/2}$ (the choice of the square root matrix unimportant). We have

$$u^\top \hat{u} = \Phi(M),$$

$$u^\top \hat{v} = \Phi_*,$$

$$v^\top \hat{v} = c^\top M_*^{-}\hat{u} \leq \Phi_*,$$

where the last equality follows from the equivalence theorem for c-optimality. Also, $c^\top M^{-1}H_iM^{-1}c = u^\top \Omega_i u$ and the equivalence theorem gives $c^\top M_*^{-}\hat{H}_iM_*^{-}c = v^\top \Omega_i v = \Phi_* = u^\top \hat{v}$. Moreover, denoting by $\psi$ the angle between $u$ and $v$, we have

$$\cos(\psi) = \frac{u^\top \hat{v}}{\|u\|\|v\|} \geq \left[\frac{\Phi_*}{\Phi(M)}\right]^{1/2} \geq \frac{1}{(1 + \delta)^{1/2}},$$

see (2.5).

We first prove (i). Using (A.1), we have

$$c^\top M^{-1}H_iM^{-1}c = u^\top \Omega_i u = \|u\|^2 \hat{u}^\top \Omega_i \hat{u} \geq \|u\|^2 \left(\hat{v}^\top \Omega_i \hat{v} - \Delta_i |\sin(\psi)|\right)$$

$$= \Phi(M) \left(\frac{\Phi_*}{c^\top M_*^{-}\hat{u} c} - \Delta_i |\sin(\psi)|\right)$$

$$\geq \Phi(M) \left(1 - \Delta_i |\sin(\psi)|\right) \geq \Phi(M) \left\{1 - \Delta_i [\delta/(1 + \delta)]^{1/2}\right\}.$$

This proves (i) and therefore (2.7).

We now prove (ii). Denote $R_i = (c^\top M^{-1}H_iM^{-1}c)/\Phi(M) = (u^\top \Omega_i u)/(u^\top \hat{u})$. We can write

$$R_i = \frac{\Phi_*}{\|v\|^2} \hat{v}^\top \Omega_i \hat{v} \geq \frac{\hat{u}^\top \Omega_i \hat{u}}{\hat{v}^\top \Omega_i \hat{v}}.$$

Applying (A.2) of Lemma A.1 to the right-hand side and using the property that for a fixed $\kappa$ $\gamma(\kappa, \psi)$ is a decreasing function of the angle $\psi$, with $\gamma(\kappa, 0) = 1$ and $\gamma(\kappa, \pi/2) = 1/\kappa$, we obtain that $R_i \geq \gamma(\kappa_i, \varphi)$ with $\varphi$ such that $\cos(\varphi) = 1/(1 + \delta)^{1/2}$. This proves (ii) and thus (2.8).
B Proof of Theorems 2.2 and 3.3

Proof. (of Theorem 2.2) The dimension is irrelevant, since for $m > 2$ we can consider block matrices $A_m$ of the form

$$A_m = \left( \begin{array}{cc} A_2 & O_{2,m-2} \\ O_{m-2,2} & I_{m-2} \end{array} \right),$$

with $A_2$ a $2 \times 2$ matrix, $O_{p,q}$ the $p \times q$ null matrix, and a vector $c = (c_1, c_2, 0, \ldots, 0)^\top$. We can thus restrict our attention to the case $m = 2$. Take $H = \{H_1, H_2, M\}$ with

$$H_1 = \begin{pmatrix} 1 & 0 \\ 0 & 1 + a + t \end{pmatrix}, \quad H_2 = \begin{pmatrix} 1 + a + t & 0 \\ 0 & 1 \end{pmatrix} \quad \text{and} \quad M = \begin{pmatrix} 1 + t & 0 \\ 0 & 1 \end{pmatrix},$$

with $a, t > 0$.

For $c = (\sqrt{2}/2, 1, 1)^\top$, the $c$-optimal matrix is $M_* = (H_1 + H_2)/2$, with $\Phi_* = \Phi(M_*) = 2/(2 + a + t) < \Phi(M) = (2 + t)/(2(1 + t))$. We have

$$c^\top M^{-1}H_1M^{-1}c - c^\top M^{-1}H_2M^{-1}c = \frac{t(2 + t)(a + t)}{2(1 + t)^2},$$

which is positive for $a, t > 0$, and therefore $\delta = (c^\top M^{-1}H_1M^{-1}c)/\Phi(M) - 1 = [t^2 + at + a/(2 + t)]/(1 + t)$. The matrix $\Omega_2 = M^{-1/2}H_2M^{-1/2}$ is

$$\Omega_2 = \begin{pmatrix} 1 + a + t & 0 \\ 0 & 1 \end{pmatrix}.$$

Consider first the bound (2.7) and denote $b_1(a, t) = B_1(M, H_2) = \{1 - \Delta_2 [\delta/(1 + \delta)]^{1/2}\} \Phi(M) - c^\top M^{-1}H_2M^{-1}c$, with $\Delta_2 = \lambda_{\max}(\Omega_2) - \lambda_{\min}(\Omega_2) = a/(1 + t)$. We obtain

$$b_1(a, t) = -a \{ (2 + t) [\delta/(1 + \delta)]^{1/2} + 1 \} / 2(1 + t)^2.$$

The expression of $\delta$ gives

$$a = a_\delta(t) = \frac{(2 + t)(\delta + \delta t - t^2)}{(1 + t)^2}, \quad (B.1)$$

which is positive for $0 < t < t_\delta = \{\delta + [\delta(4 + \delta)]^{1/2}\}/2$. We then obtain $b_1(a_\delta(t), t) < 0$ for $t \in (0, t_\delta)$, with $b_1(a_\delta(t), t) = 0$ and $\text{db}_1(a_\delta(t), t)/\text{dt}\big|_{t = t_\delta} > 0$ for all $\delta > 0$, with $\text{db}_1(a_\delta(t), t)/\text{dt}\big|_{t = t_\delta} = 0$ for $\delta = 0$, $\text{db}_1(a_\delta(t), t)/\text{dt}\big|_{t = t_\delta} = 2 \delta^{1/2} - 3 \delta + (3/4) \delta^{3/2} + O(\delta^2)$ for small $\delta$, and $\text{db}_1(a_\delta(t), t)/\text{dt}\big|_{t = t_\delta} > 0$ for all $\delta > 0$; reaching its maximum $\simeq 0.409910$ at $\delta \simeq 0.257688$.

We proceed similarly for the bound (2.8). Denote $b_2(a, t) = B_2(M, H_2) = \gamma(\kappa, \varphi) \Phi(M) - c^\top M^{-1}H_2M^{-1}c$, with $\kappa = \lambda_{\max}(\Omega_2)/\lambda_{\min}(\Omega_2) = 1 + a/(1 + t)$ and $\varphi = \arccos\left([1 + \delta]^{-1/2}\right)$. We get $b_2(a_\delta(t), t) < 0$ for $t \in (0, t_\delta)$, $b_2(a_\delta(t_\delta), t_\delta) = 0$ and $\text{db}_2(a_\delta(t), t)/\text{dt}\big|_{t = t_\delta} > 0$ for all $\delta > 0$; $\text{db}_2(a_\delta(t), t)/\text{dt}\big|_{t = t_\delta} = 0$ at $\delta = 0$, equals $2 \delta^{1/2} - 11 \delta + (99/4) \delta^{3/2} + O(\delta^2)$ for small $\delta$, with a maximum $\simeq 0.121304$ at $\delta \simeq 0.022087$. ($b_2(a_\delta(t), t) > b_1(a_\delta(t), t)$ for all $t \in (0, t_\delta)$, $\delta > 0$, which suggests that here (2.8) is more accurate than (2.7).)

For any $\epsilon > 0$, we can thus choose a $t_* \in (0, t_\delta)$ such that $b_2(a_\delta(t_*), t_*) + \epsilon > 0$ and $b_1(a_\delta(t_*), t_*) + \epsilon > 0$. For $t = t_*$ and $a = a_\delta(t_*)$, the inequalities $B_1(M, H_2) + \epsilon > 0$ and $B_2(M, H_2) + \epsilon > 0$ are thus satisfied although $H_2$ supports the optimal design.

\[\square\]
Proof. (of Theorem 3.3) We use the same construction as in the proof of Theorem 2.2, and denote $b_3(a,t) = B_3(M,H_2)$. We get $\lambda_{\max}(A_2^T M^{-1} A_2) = 1 + a/(1 + t)$ in (C.1), the optimal $\beta$ is

$$
\beta^* = \frac{\delta^2 (t + 1) + \delta (2 t + 1) + [t/(1 + t)]^{1/2} [\delta (t + 1) - t^2]}{\delta (1 + t)^2}.
$$

It gives $b_3(a_3(t), t) < 0$ for $0 < t < t_3 = \{\delta + |\delta(4 + \delta)|^{1/2}\}/2$, with $a_3(t)$ given by (B.1), and $b_3(a_3(t_3), t_3) = 0$. Also, $db_3(a_3(t), t)/dt|_{t=t_3} > 0$ for all $\delta > 0$, with $db_3(a_3(t), t)/dt|_{t=t_3} = 0$ for $\delta = 0$, $db_3(a_3(t), t)/dt|_{t=t_3} = 2 \delta^{1/2} - 2 \delta - (3/4) \delta^{3/2} + O(\delta^2)$ for small $\delta$, and $db_3(a_3(t), t)/dt|_{t=t_3}$ reaching its maximum $\approx 0.510064$ at $\delta \approx 0.346719$. For any $\epsilon > 0$, we can thus choose a $t_* \in (0, t_3)$ such that $b_3(a_3(t_*), t_*) + \epsilon > 0$, and for $t = t_*$ and $a = a_3(t_*)$ we have $B_3(M,H_2) + \epsilon > 0$ although $H_2$ supports the optimal design.

\[\square\]

C Proof of Theorem 3.2

Proof. Define the set

$$
\mathcal{F}(u) = \{ z \in \mathbb{R}^m : c^T z \geq c^T u \text{ and } \|A_j^T z\| \leq 1, \ j = 1, \ldots, q \},
$$

where $u = M^{-1} c/(\max_j \|A_j^T M^{-1} c\|)$ is feasible for $P - \text{SOCP}$. Then, $\sup_{z \in \mathcal{F}(u)} \|A_j^T z\| < 1$ implies that $\|A_j^T u\| < 1$ and, from Theorem 3.1, $H_j$ cannot support an optimal design. Checking the condition $\sup_{z \in \mathcal{F}(u)} \|A_j^T z\| < 1$ requires the solution of a non-convex quadratic optimization problem, but a much simpler condition can be derived using the so-called S-procedure; see, e.g., Pölk and Térakly (2007). Let $t$ be an upper bound on $\sup_{z \in \mathcal{F}(u)} \|A_j^T z\|$, so that $\|A_j^T z\| \leq t$ for any $z \in \mathcal{F}(u)$. This is equivalent to

$$
\begin{pmatrix}
1 \\
(1)\top
\end{pmatrix}
\begin{pmatrix}
1 \\
(1)\top
\end{pmatrix}
\begin{pmatrix}
-\langle c, u \rangle^2 & 0_m^T \\
0_m & cc^\top
\end{pmatrix}
\begin{pmatrix}
1 \\
(1)\top
\end{pmatrix}
\geq 0 \text{ and } \begin{pmatrix}
1 \\
(1)\top
\end{pmatrix}
\begin{pmatrix}
1 & 0_m \\
0_m & -H_j
\end{pmatrix}
\begin{pmatrix}
1 \\
(1)\top
\end{pmatrix}
\geq 0 \text{ for all } j = 1, \ldots, q,
$$

\[\implies \begin{pmatrix}
1 \\
(1)\top
\end{pmatrix}
\begin{pmatrix}
t \\
0_m
\end{pmatrix}
\begin{pmatrix}
0_m^T \\
0_m & -H_j
\end{pmatrix}
\begin{pmatrix}
1 \\
(1)\top
\end{pmatrix}
\geq 0,
\]

with $0_m$ the null vector of length $m$. A sufficient condition for this implication is the existence of nonnegative $\mu$ and $\tau_j$ ($j = 1, \ldots, q$) such that

$$
\begin{pmatrix}
t \\
0_m
\end{pmatrix}
\begin{pmatrix}
0_m^T \\
0_m & -H_j
\end{pmatrix}
\leq \mu
\begin{pmatrix}
-\langle c, u \rangle^2 & 0_m^T \\
0_m & cc^\top
\end{pmatrix}
+ \sum_{j=1}^q \tau_j
\begin{pmatrix}
1 & 0_m \\
0_m & -H_j
\end{pmatrix},
$$

or equivalently, $t \geq \sum_{j=1}^q \tau_j - \mu \langle c, u \rangle^2$ and $\sum_{j=1}^q \tau_j H_j - H_j \succeq \mu cc^\top$, where, for $A$ and $B$ two square matrices of the same size, $A \succeq B$ means that $A - B$ is nonnegative definite.

Finding the smallest $t \geq 0$ such that the above linear matrix inequality holds for some nonnegative $\mu$ and $\tau_j$ is a Semi-Definite Program (SDP), the solution of which would allow the elimination of $H_j$ when $t < 1$. However, solving this SDP is probably as hard (if not harder) than solving the original problem, and we restrict the search for solutions to $\tau = (\tau_1, \ldots, \tau_q)^\top$ being proportional to $w$, that is, to $\tau = \beta w$ with $\beta > 0$. This shows that $H_j$ can be safely eliminated if the optimal value $t^*$ of the following simplified SDP is strictly less than 1:

$$
t^* = \inf_{\beta, \mu \geq 0} \beta - \mu \langle c, u \rangle^2 \text{ subject to } \beta M - H_j \succeq \mu cc^\top.
$$

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Any feasible $\beta$ must satisfy $\beta \mathbf{M} \succeq \mathbf{H}_i$, that is, $\beta \geq \lambda_{\max}(\mathbf{A}_i^T \mathbf{M}^{-1} \mathbf{A}_i)$, and a simple application of the Schur-complement lemma shows that the largest feasible $\mu$ is $\mu^* = [c^\top (\beta \mathbf{M} - \mathbf{H}_i)]^{-1}$. Hence the computation of $t^*$ reduces to solving a one-dimensional optimization problem,

$$t^* = \inf_{\beta \geq \lambda_{\max}(\mathbf{A}_i^T \mathbf{M}^{-1} \mathbf{A}_i)} \beta - \frac{(c^\top \mathbf{u})^2}{c^\top (\beta \mathbf{M} - \mathbf{H}_i)}^{-1} c,$$

(C.1)

where the function $f(\cdot)$ to be minimized is defined by continuity at $\beta_{\min} = \lambda_{\max}(\mathbf{A}_i^T \mathbf{M}^{-1} \mathbf{A}_i)$ by $f(\beta_{\min}) = \beta_{\min}$, and $\mathbf{H}_i$ can be eliminated when $t^* < 1$. From the definition (2.1) of $\delta$, we have $(c^\top \mathbf{u})^2 = \Phi(\mathbf{M})/(1 + \delta)$, so that the condition $t^* < 1$ is equivalent to $B_3(\mathbf{M}, \mathbf{H}_i) > 0$, which concludes the proof.

### Acknowledgement

The work of Luc Pronzato was partly supported by project INDEX (INcremental Design of EXperiments) ANR-18-CE91-0007 of the French National Research Agency (ANR). The work of Guillaume Sagnol was supported by the Deutsche Forschungsgemeinschaft (DFG, German Research Foundation) under Germany’s Excellence Strategy – The Berlin Mathematics Research Center MATH+ (EXC-20461, project ID: 390685689).

### References


