Degrees bounding principles and universal instances in reverse mathematics
Ludovic Patey

To cite this version:

HAL Id: hal-01888599
https://hal.archives-ouvertes.fr/hal-01888599
Submitted on 5 Oct 2018

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L’archive ouverte pluridisciplinaire HAL, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d’enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.
DEGREES BOUNDING PRINCIPLES
AND UNIVERSAL INSTANCES IN REVERSE MATHEMATICS

LUDOVIC PATEY

ABSTRACT. A Turing degree $d$ bounds a principle $P$ of reverse mathematics if every computable instance of $P$ has a $d$-computable solution. $P$ admits a universal instance if there exists a computable instance such that every solution bounds $P$. We prove that the stable version of the ascending descending sequence principle ($SADS$) as well as the stable version of the thin set theorem for pairs ($STS_2^2$) do not admit a bound of low$_2$ degree. Therefore no principle between Ramsey’s theorem for pairs ($RT^2_2$) and SADS or STS$_2^2$ admit a universal instance. We construct a low$_2$ degree bounding the Erdős-Mosser theorem ($EM$), thereby showing that the previous argument does not hold for $EM$. Finally, we prove that the only $\Delta^0_2$ degree bounding a stable version of the rainbow Ramsey theorem for pairs ($SRRT^2_2$) is $0'$. Hence no principle between the stable Ramsey theorem for pairs ($SRT^2_2$) and $SRRT^2_2$ admit a universal instance. In particular the stable version of the Erdős-Mosser theorem does not admit one. It remains unknown whether $EM$ admits a universal instance.

1. INTRODUCTION

Reverse mathematics is a program whose goal is to classify theorems according to their computational strength, within the framework of subsystems of second-order arithmetic. Proofs are done relatively to a very weak system ($RCA_0$) meant to capture computational mathematics. $RCA_0$ is composed of basic Peano axioms, $\Delta^0_1$ comprehension and $\Sigma^0_1$ induction schemes. See [12] for a good introductory book. Most of statements in reverse mathematics are of the form

$$\forall X(\Phi(X) \rightarrow \exists Y \Psi(X,Y))$$

where $\Phi$ and $\Psi$ are arithmetic formulas.

A set $X$ such that $\Phi(X)$ holds is called an instance of $P$ and a set $Y$ such that $\Psi(X,Y)$ holds is a solution to $X$. We can see relations between two instances $X_1,X_2$ of a statement $P$ as a mass problem consisting of computing a solution to $X_1$ given any solution to $X_2$.

Definition 1.1 Given a statement $P$, a degree $d$ is $P$-bounding ($d \gg_P \emptyset$) if every computable instance $X$ of $P$ has a $d$-computable solution. A statement $P$ admits a universal instance if it has a computable instance $X$ such that every solution to $X$ bounds $P$.

The notation $d \gg \emptyset$ historically means that the degree $d$ is PA and therefore is equivalent to $d \gg_{WKL_0} \emptyset$ where $WKL_0$ is weak König’s lemma principle, i.e., König’s lemma restricted to subtrees of $2^{<\omega}$. It is well-known that WKL$_0$ admits a universal instance – e.g. take the $\Pi^0_1$ class of completions of Peano arithmetics. A few principles have been proven to admit universal instances – $WKL_0$ [22], König’s lemma ($KL$) [12], the

Date: June 23, 2015.
Ramsey-type weak weak König’s lemma (RWWKL) \cite{1}, the finite intersection property (FIP) \cite{9}, the omitting partial type theorem (OPT) \cite{15}, or even the rainbow Ramsey theorem for pairs (RRT\(^2\)) \cite{21} – but most of principles do not admit one. An important notion for proving such a result is computable reducibility.

**Definition 1.2** Fix two statements \(P\) and \(Q\). We say that \(P\) is **computably reducible** to \(Q\) (written \(P \leq_c Q\)) if for every instance \(X \leq P\) there is an \(X\)-computable instance \(Y\) of \(Q\) such that each solution to \(Y\) computes relative to \(X\) a solution to \(X\). \(P\) and \(Q\) are **computably equivalent** if \(P \leq_c Q\) and \(Q \leq_c P\).

Mileti proved in \cite{20} that the stable Ramsey theorem for pairs (SRT\(^2\)) admits no bound of low\(_2\) degree. Therefore every statement \(P\) having an \(\omega\)-model with only low\(_2\) sets, and such that SRT\(^2\) \(\leq_c\) \(P\), admits no universal instance. In particular none of Ramsey’s theorem for pairs (RT\(^2\)), SRT\(^2\) and the Ramsey-type weak König’s lemma relative to \(\emptyset’\) (RWKL[\(\emptyset’\)]) admit a universal instance. Independently, Hirschfeldt & Shore proved in \cite{14} that the stable ascending descending sequence principle (SADS) admits no bound of low degree. Hence neither SADS nor the stable chain antichain principle (SCAC) admit a universal instance.

We generalize both results by proving that SADS does not admit a bound of low\(_2\) degree, proving therefore that if a statement \(P\) has an \(\omega\)-model with only low\(_2\) sets and SADS \(\leq_c\) \(P\) then \(P\) admits no universal instance. We also extend the result to statements to which the stable thin set theorem for pairs (STS(2)) computably reduces. Hence we deduce that none of the ascending descending sequence principle (ADS), the chain antichain principle (CAC), the thin set theorem for pairs (TS(2)), the free set theorem for pairs (FS(2)) and their stable versions admit a universal instance.

We generalize the result to arbitrary tuples and prove that none of RT\(^n\), FS\((n)\), TS\((n)\) and their stable versions admit a universal instance for \(n \geq 2\). The question remains open for the rainbow Ramsey theorem for \(n\)-tuples (RRT\(^n\)) with \(n \geq 3\). We construct a low\(_2\) degree bounding the Erdős Moser theorem (EM), thereby showing that the previous argument does not hold for EM.

Mileti proved in \cite{20} that the only \(\Delta^0_3\) degree bounding SRT\(^2\) is \(\emptyset’\). Using the fact that every \(\Delta^0_3\) set has an infinite incomplete \(\Delta^0_2\) subset in either it or its complement \cite{13}, we obtain another proof that SRT\(^2\) admits no universal instance. We extend this result by proving that the only \(\Delta^0_3\) degree bounding a stable version of the rainbow Ramsey theorem for pairs (SRRT\(^2\)) is \(\emptyset’\). Hence none of the statements \(P\) satisfying SRRT\(^2\) \(\leq_c\) \(P\) \(\leq_c\) SRT\(^2\) admit a universal instance. In particular we deduce that neither SRRT\(^2\) nor the stable version of the Erdős-Moser theorem (SEM) admits a universal instance.

### Notation

**Formulas.** The notation \((\forall^∞ s)\varphi(s)\) means that \(\varphi(s)\) holds for all but finitely many \(s\), i.e., is translated to \((\exists \delta)(\forall s \geq \delta)\varphi(s)\). Given two sets \(X\) and \(Y\), we denote by \(X \preceq^\ast Y\) the statement \((\forall^∞ s \in X)[s \in Y]\). Accordingly, \(X =^\ast Y\) means that both \(X \subseteq^\ast Y\) and \(Y \subseteq^\ast X\) hold, i.e., \(X\) and \(Y\) differ by finitely many elements.

**Turing functional and lowness.** We fix an effective enumeration of all Turing functionals \(\Phi_0, \Phi_1, \ldots\). We denote by \(\Phi_\varepsilon\) the partial approximation of the Turing functional \(\Phi_\varepsilon\) at stage \(s\). Given a set \(X\), we denote by \(X’\) the jump of \(X\) and by \(X^{(n)}\) the \(n\)th jump of \(X\). A set \(X\) is low\(_n\) over \(Y\) if \((X \oplus Y)^{(n)} \preceq Y^{(n)}\). A set is low\(_n\) if it is low\(_n\) over \(\emptyset\). A low\(_n\)-ness index of a set \(X\) low\(_n\) over \(Y\) is a Turing index \(e\) such that \(\Phi_\varepsilon^{(n)} = (X \oplus Y)^{(n)}\).
Mathias forcing. Given two sets $E$ and $F$, we denote by $E \prec F$ the formula $(\forall x \in E)(\forall y \in F)x < y$. A Mathias condition is a pair $(F, X)$ where $F$ is a finite set, $X$ is an infinite set and $F \prec X$. A condition $(\tilde{F}, \tilde{X})$ extends $(F, X)$ (written $(\tilde{F}, \tilde{X}) \preceq (F, X)$) if $F \subseteq \tilde{F}$, $\tilde{X} \subseteq X$ and $\tilde{F} \setminus F \subseteq X$. A set $G$ satisfies a Mathias condition $(F, X)$ if $F \subseteq G$ and $G \setminus F \subseteq X$.

2. Degrees bounding cohesiveness

A standard proof of Ramsey’s theorem for pairs consists of reducing an arbitrary coloring of pairs into a stable one using the cohesiveness principle. The understanding of the links between cohesiveness and stability is a very active subject of research in reverse mathematics [4, 13, 5].

Definition 2.1 (Cohesiveness) An infinite set $C$ is $\bar{R}$-cohesive for a sequence of sets $R_0, R_1, \ldots$ if for each $i \in \omega$, $C \subseteq^* R_i$ or $C \subseteq^* \overline{R_i}$. A set $C$ is cohesive (resp. $r$-cohesive) if it is $\bar{R}$-cohesive where $\bar{R}$ is an enumeration of all c.e. (resp. computable) sets. COH is the statement “Every uniform sequence of sets $\bar{R}$ has an $\bar{R}$-cohesive set.”

Jockusch et al. proved in [16] the existence of a low$_2$ cohesive set. Degrees bounding COH are quite well understood and admit a simple characterization:

Theorem 2.2 (Jockusch & Stephan [16]) Fix an $n \in \omega$.

1. For every set $C$ such that $C \not\geq T C$, $C \not\geq_{COH} \emptyset$.
2. There exists a uniformly $\emptyset^{(n)}$-computable sequence of sets $\bar{R}$ such that for every $\bar{R}$-cohesive set $C$, $(C \oplus \emptyset^{(n)})' \not\geq \emptyset^{(n+1)}$.

In particular, taking a set $P \not\geq T \emptyset$ low over $\emptyset'$ and a set $C$ such that $C' = T P$ whose existence is ensured by Friedberg’s jump inversion theorem, we obtain a low$_2$ degree bounding COH. The canonical $\emptyset^{(n)}$-computable sequence of sets $\bar{R}$ whose existence is claimed in clause 2 of Theorem 2.2 is

$$\bar{R}_e = \{s : \Phi_{e, \mathcal{E}}^\emptyset(e) \downarrow = 1\}$$

Every $\bar{R}$-cohesive set $C$ computes a function $f(\cdot, \cdot)$ such that $\lim_{e \in \omega} f(e, s)$ exists for each $e \in \omega$ and $\lim_{e \in \omega} f(e, s) = \Phi_{e, \mathcal{E}}^{\emptyset^{(n+1)}}(e)$ for each Turing index $e$ such that $\Phi_{e, \mathcal{E}}^{\emptyset^{(n+1)}}(e) \downarrow$. By a relativized version of Schoenfield’s limit lemma, $(C \oplus \emptyset^{(n)})'$ computes the function $\check{f}(x) = \lim_{e \in \omega} f(x, s)$ and is therefore of PA degree relative to $\emptyset^{(n+1)}$.

Corollary 2.3 COH admits a universal instance.

Proof. The uniformly computable sequence of sets $\bar{R}$ such that the jump of every $\bar{R}$-cohesive set is of PA degree relative to $\emptyset'$ is a universal instance by the previous theorem.

Wang proved in [26] that for every set $P \not\geq T \emptyset''$ and every uniformly $\emptyset'$-computable sequence of sets $\bar{R}$, there exists an $\bar{R}$-cohesive set $C$ such that $C'' \leq_T C \oplus \emptyset'' \leq_T P$. Cholak et al. used in [4] the existence of a low subuniform degree to deduce the existence, for every set $P \not\geq T \emptyset'$, of an $r$-cohesive set $C$ such that $C' \leq_T P$. We can apply a similar reasoning for $\emptyset'$-computable sets, using the fact that degrees bounding COH are somehow subuniform degrees for $\Delta^0_2$ approximations.
Theorem 2.4 For every set \( P \gg \emptyset'' \), there exists an \( \vec{R} \)-cohesive set \( C \) such that \( C'' \leq_T C \oplus \emptyset'' \leq_T P \), where \( \vec{R} \) is the (non-uniformly computable) sequence of all \( \emptyset' \)-computable sets.

Proof. Let \( \vec{U} \) be the uniformly computable sequence of sets defined by

\[
U_{e,x} = \{ s : \Phi_{e,s}(x) = 1 \}
\]

Fix a low\(_2\) \( \vec{U} \)-cohesive set \( C_0 \) and its \( C_0' \)-computable bijection \( f : \omega \to C_0 \). Every set \( P \gg \emptyset'' \), \( P \gg C_0'' \). Consider the uniformly \( C_0' \)-computable sequence of sets

\[
V_e = \{ x : \lim s \Phi_{e,s}(x) = 1 \}
\]

The sequence \( \vec{V} \) contains every \( \emptyset' \)-computable set. In particular, every \( \vec{V} \)-cohesive set is \( \vec{R} \)-cohesive. By a relativization of Wang’s result, there exists an \( \vec{V} \)-cohesive set \( C \) such that \( (C \oplus C_0')'' \leq_T C \oplus \emptyset'' \leq_T P \). \( \square \)

The proof of the previous theorem shows that an application of \text{COH} followed by an application of \text{COH}\([\emptyset']\) are enough to obtain a set of degree bounding \text{COH}\([\emptyset']\). The following question remains open:

Question 2.5 Does \text{COH}\([\emptyset']\) admit a universal instance?

3. Degrees bounding the atomic model theorem

The atomic model theorem is a statement of model theory admitting a simple, purely computability-theoretic characterization over \( \omega \)-models. This statement happens to have a weak computational content and is therefore a consequence of many other principles in reverse mathematics. For those reasons, the atomic model theorem is a good candidate for factorizing proofs of properties which are closed upward by the consequence relation.

Definition 3.1 (Atomic model theorem) A formula \( \varphi(x_1, \ldots, x_n) \) of \( T \) is an atom of a theory \( T \) if for each formula \( \psi(x_1, \ldots, x_n) \), one of \( T \vdash \varphi \rightarrow \psi \) and \( T \vdash \varphi \rightarrow \neg \psi \) holds, but not both. A theory \( T \) is atomic if for every formula \( \psi(x_1, \ldots, x_n) \) consistent with \( T \), there exists an atom \( \varphi(x_1, \ldots, x_n) \) of \( T \) extending it, i.e., one such that \( T \vdash \varphi \rightarrow \psi \). A model \( \mathcal{A} \) of \( T \) is atomic if every \( n \)-tuple from \( \mathcal{A} \) satisfies an atom of \( T \). AMT is the statement “Every complete atomic theory has an atomic model”.

AMT has been introduced as a principle by Hirschfeldt et al. in [15]. They proved that \( \text{WKL}_0 \) and AMT are incomparable on \( \omega \)-models, proved over \( \text{RCA}_0 \) that AMT is strictly weaker than SADS. The author proved in [23] that \( \text{STS}(2) \) implies AMT over \( \text{RCA}_0 \). In this section we use the fact that AMT is not bounded by any \( \Delta^0_2 \) low\(_2\) degree to deduce that none of AMT, SADS and SCAC admits a universal instance. The principle AMT has been proven in [15, 6] to be computably equivalent to the following principle:

Definition 3.2 (Escape property) For every \( \Delta^0_2 \) function \( f \), there exists a function \( g \) such that \( f(x) \leq g(x) \) for infinitely many \( x \).
This equivalence does not hold over RCA₀ as, unlike AMT, the escape property implies Σ²₀ over Σ²₀ [15]. Using this characterization, we can easily deduce the two following theorems:

**Theorem 3.3** (Hirschfeldt et al. [15]) There is no low₂ Δ²₀ degree bounding AMT.

**Theorem 3.4** No principle $P$ having an ω-model with only low sets and such that AMT $\leq_c P$ admits a universal instance.

Theorem 3.3 and Theorem 3.4 can be easily proven using the following characterization of Δ²₀ low₂ sets in terms of domination:

**Lemma 3.5** (Martin, [19]) A set $A \leq_T \emptyset'$ is low₂ iff there exists an $f \leq_T \emptyset'$ dominating every $A$-computable function.

**Proof.** A set $A$ is low₂ iff $\emptyset'$ is high relative to $A$. We conclude the lemma from the observation that a set $X$ is high relative to a set $A \leq_T \emptyset'$ iff it computes a function dominating every $A$-computable function. □

**Remark.** As explained Conidis in [6], Theorem 3.3 cannot be extended to every low₂ sets: Soare [6] constructed a low₂ set bounding the escape property using a forcing argument. So there exists a low₂ degree bounding AMT.

**Proof of Theorem 3.4.** Suppose for the sake of contradiction that $P$ has a universal instance $U$ and an ω-model $\mathcal{M}$ with only low sets. As $U$ is computable, $U \in \mathcal{M}$. Let $X \in \mathcal{M}$ be a (low) solution to $U$. In particular, $X$ is low₂ and Δ²₀, so by Lemma 3.5 and the computable equivalence of AMT and the escape property, there exists a computable instance $Y$ of AMT such that $X$ does not compute a solution to $Y$. As AMT $\leq_c P$, there exists a $Y$-computable (hence computable) instance $Z$ of $P$ such that every solution to $Z$ computes a solution to $Y$. Thus $X$ does not compute a solution to $Z$, contradicting universality of $U$. □

Hirschfeldt et al. proved in [14] the existence of an ω-model of SADS and SCAC with only low sets. Therefore we obtain another proof that neither SADS nor SCAC admits a universal instance. The result was first proven in [14] using an ad-hoc notion of reducibility.

**Corollary 3.6** None of AMT, SADS and SCAC admit a universal instance.

The previous argument cannot directly be applied to SRT²₂, SEM or STS(2) as none of those principles admit an ω-model with only low sets [10, 17, 23]. However Lemma 3.4 can be extended to principles such that every computable instance has a Δ²₀ low₂ solution. It is currently unknown whether every Δ²₀ set admits a Δ²₀ low₂ infinite subset in either it or its complement. A positive answer would lead to a proof that SRT²₂, SEM and STS(2) have no universal instance, and more importantly, would provide an ω-model of SRT²₂ that is not a model of DNR[\emptyset'] as explained in [13]. We shall see later that none of SRT²₂, SEM and STS(2) admits a universal instance.

4. DEGREES BOUNDING STS(2) AND SADS

Mileti originally proved in [20] that no principle $P$ having an ω-model with only low₂ sets and satisfying SRT²₂ $\leq_c P$ admits a universal instance, and deduced that none of
SRT\(_2^2\) and RT\(_2^2\) admit one. In this section, we reapply his argument to much weaker statements and derive non-universality results to a large range of principles in reverse mathematics. Thin set theorem and ascending descending sequence are example of statements weak enough to be a consequence of many others, and surprisingly strong enough to diagonalize against low\(_2\) sets.

**Definition 4.1** (Thin set) Let \(k \in \omega\) and \(f : [\omega]^k \to \omega\). A set \(A\) is thin for \(f\) if \(f([A]^k) \neq \omega\), that is, if the set \(A\) “avoids” at least one color. TS\((k)\) is the statement “every function \(f : [\omega]^k \to \omega\) has an infinite set thin for \(f\)”. A function \(f : [\omega]^k \to \omega\) is stable if \(\forall \sigma \in [\omega]^{k-1}, \lim f(\sigma, s)\) exists. STS\((k)\) is the restriction of TS\((k)\) to stable functions.

Cholak et al. studied extensively thin set principle in [3]. Some of the results where already stated by Friedman without giving a proof, notably there exists an \(\omega\)-model of WKL\(_0\) which is not a model of TS\((2)\), and the arithmetical comprehension axiom (ACA\(_0\)) does not imply (\(\forall k\))TS\((k)\) over RCA\(_0\). Wang showed in [28] that (\(\forall k\))TS\((k)\) does not imply ACA\(_0\) on \(\omega\)-models. Rice [24] proved that STS\((2)\) implies DNR over RCA\(_0\). The author proved in [23] that RCA\(_0\) \(\vdash\) TS\((2)\) \(\rightarrow\) RT\(_2^2\).

**Definition 4.2** (Ascending descending sequence) ADS is the statement “Every infinite linear order admits an infinite ascending or descending sequence”. SADS is the restriction of ADS to order types \(\omega + \omega^*\).

Tennenbaum [25] constructed a computable linear order of order type \(\omega + \omega^*\) with no computable ascending or descending sequence. Therefore SADS does not hold over RCA\(_0\). Hirschfeldt & Shore [14] studied ADS within the framework of reverse mathematics, proving that ADS implies both COH and BΣ\(^0_2\) over RCA\(_0\) and that SADS implies AMT over RCA\(_0\). They constructed an \(\omega\)-model of ADS that is not a model of DNR, and an \(\omega\)-model of COH + WKL\(_0\) that is not a model of SADS.

The study of degrees bounding a statement and the existence of a universal instance are closely related. As does Miletin in [20], we deduce two kind of theorems by the application of his proof technique.

**Theorem 4.3** There exists no low\(_2\) degree bounding any of STS\((2)\) or SADS.

**Theorem 4.4** No principle \(P\) having an \(\omega\)-model with only low\(_2\) sets and such that any of STS\((2)\), SADS is computably reducible to \(P\) admits a universal instance.

The proof of the two theorems is split into three lemmas. Lemma 4.7 provides a general way of obtaining bounding and universality results, assuming the ability of a principle to diagonalize against a particular set. Lemma 4.8 and Lemma 4.9 state the desired diagonalization for respectively STS\((2)\) and SADS.

**Corollary 4.5** None of the following principles admits a universal instance: RT\(_2^2\), RWKL[\(\emptyset^0\)], FS\((2)\), TS\((2)\), CAC, ADS and their stable versions.

**Proof.** Each of the above mentioned principles is a consequence of RT\(_2^2\) over RCA\(_0\) and computably implies either SADS or STS\((2)\). See [11] for RWKL[\(\emptyset^0\)], [3] for FS\((2)\) and TS\((2)\), and [14] for CAC and ADS. By Theorem 3.1 of [4], there exists an \(\omega\)-model of RT\(_2^2\) having only low\(_2\) sets. The result now follows from Theorem 4.4. \(\square\)
In order to prove Theorem 4.3 and Theorem 4.4, we need the following theorem proven by Mileti. It simply consists of applying a relativized version of the low basis theorem to a $\Pi^0_1$ class of completions of the enumeration of all partial computable sets.

**Theorem 4.6** (Mileti, Corollary 5.4.5 of [20]) For every set $X$, there exists $f : \omega^2 \to \{0, 1\}$ low over $X$ such that for every $X$-computable set $Z$, there exists an $e \in \omega$ with $Z = \{a \in \omega : f(e, a) = 1\}$.

**Lemma 4.7** Fix an $n \in \omega$ and two principles $P$ and $Q$ such that $P \leq_e Q$. Suppose that for any $f : \omega^2 \to \{0, 1\}$ satisfying $f'' \leq_T \emptyset^{(n+2)}$, there exists a computable instance $I$ of $P$ such that for each $e \in \omega$, if $\{a \in \omega : f(e, a) = 1\}$ is infinite then it is not a solution to $I$. Then the following holds:

(i) For any degree $d$ low$_2$ over $\emptyset^{(n)}$ there is a computable instance $U$ of $P$ such that $d$ does not bound a solution to $U$.
(ii) There is no degree low$_2$ over $\emptyset^{(n)}$ bounding $P$.
(iii) If every computable instance $I$ of $Q$ has a solution low$_2$ over $\emptyset^{(n)}$, then $Q$ has no universal instance.

**Proof.**

(i) Consider any set $X$ of degree low$_2$ over $\emptyset^{(n)}$. By Theorem 4.6, there exists a function $f : \omega^2 \to \{0, 1\}$ low over $X$, hence low$_2$ over $\emptyset^{(n)}$, such that any $X$-computable set $Z$ is of the form $\{a \in \omega : f(e, a) = 1\}$ for some $e \in \omega$. Take a computable instance $I$ of $P$ having no solution of the form $\{a \in \omega : f(e, a) = 1\}$ for any $e \in \omega$. Then $X$ does not compute a solution to $I$.

(ii) Immediate from (i).

(iii) Take any computable instance $U$ of $Q$. By assumption, $U$ has a solution $X$ low$_2$ over $\emptyset^{(n)}$. By (i), there exists an instance $I$ of $P$ such that $X$ does not compute a solution to $I$. As $P \leq_e Q$, there exists an $I$-computable (hence computable) instance $J$ of $Q$ such that any solution to $J$ computes a solution to $I$. Then $X$ does not compute a solution to $J$, hence $U$ is not a universal instance.

We will prove the following lemmas which, together with Lemma 4.7, are sufficient to deduce Theorem 4.3 and Theorem 4.4.

**Lemma 4.8** Fix a set $X$. Suppose $f : \omega^2 \to \{0, 1\}$ satisfies $f'' \leq_T X''$. There exists an $X$-computable stable coloring $g : [\omega] \to \omega$ such that for all $e \in \omega$, if $\{a \in \omega : f(e, a) = 1\}$ is infinite then it is not thin for $g$.

**Lemma 4.9** Fix a set $X$. Suppose $f : \omega^2 \to \{0, 1\}$ satisfies $f'' \leq_T X''$. There exists a stable $X$-computable linear order $L$ such that for all $e \in \omega$, if $\{a \in \omega : f(e, a) = 1\}$ is infinite then it is neither an ascending nor a descending sequence in $L$.

Before proving the two remaining lemmas, we relativize the results to colorings over arbitrary tuples.

**Theorem 4.10** For any $n$, there exists no degree low$_2$ over $\emptyset^{(n)}$ bounding $\text{STS}(n + 2)$. 


Proof. Apply Lemma 4.8 relativized to $X = \emptyset^{(n)}$ together with Lemma 4.7. Simply notice that if $f : [\omega]^n \to \omega$ is a $\emptyset'$-computable coloring, the computable coloring $g : [\omega]^{n+1} \to \omega$ obtained by an application of Schoenfield’s limit lemma is such that every infinite set thin for $g$ is thin for $f$.

**Theorem 4.11** For any $n$, no principle $P$ having an $\omega$-model with only low$_2$ over $\emptyset^{(n)}$ sets and such that $\text{STS}(n+2) \leq_e P$ admits a universal instance.

**Proof.** Same reasoning as Theorem 4.4 using the notice in the proof of Theorem 4.10.

**Theorem 4.12** For any $n$, none of $\text{RT}_2^{n+2}$, $\text{RWKL}[\emptyset^{(n+1)}]$, $\text{FS}(n+2)$, $\text{TS}(n+2)$ and their stable versions admits a universal instance.

**Proof.** Fix an $n \in \omega$. Each of the above cited principles $P$ satisfies $\text{STS}(n+2) \leq_e P$ and is a consequence of $\text{RT}_2^{n+2}$ over $\omega$-models. Cholak et al. [4] proved the existence of an $\omega$-model of $\text{RT}_2^{n+2}$ having only low$_2$ over $\emptyset^{(n)}$ sets. Apply Theorem 4.11.

We now turn to the proofs of Lemma 4.8, and Lemma 4.9.

**Proof of Lemma 4.8.** We prove it in the case when $X = \emptyset$. The general case follows by a straightforward relativization. For each $e \in \omega$, let $Z_e = \{a \in \omega : f(e, a) = 1\}$. The proof is very similar to [20, Theorem 5.4.2]. We build a $\emptyset'$-computable function $c : \omega \to \omega$ such that for all $e \in \omega$, if $Z_e$ is infinite then it is not thin for $c$. Given such a function $c$, we can then apply Schoenfield’s limit lemma to obtain a stable computable function $h : [\omega]^{2} \to \omega$ such that for each $x \in \omega$, $\lim h(x, s) = c(x)$. Every set thin for $h$ is thin for $c$, and therefore for all $e \in \omega$, if $Z_e$ is infinite then it is not thin for $h$.

Suppose by Kleene’s fixpoint theorem that we are given a Turing index $d$ of the function $c$ as computed relative to $\emptyset$. The construction is done by a finite injury priority argument satisfying the following requirements for each $e, i \in \omega$:

$$\mathcal{R}_{e,i} : Z_e \text{ is finite or } (\exists a)(f(e, a) = 1 \text{ and } \Phi^\emptyset_d)(a) = i$$

The requirements are ordered in a standard way, that is, following the pairing of the indexes. Notice that each of these requirement is $\Sigma^I_2$, and furthermore we can effectively find an index for each as such. Therefore, for each $e$ and $i \in \omega$, we can effectively find an integer $m_{e,i}$ such that $R_{e,i}$ is satisfied if and only if $m_{e,i} \in f''$. By Schoenfield’s limit Lemma relativized to $\emptyset'$ and low$_2$-nes of $f$, there exists a $\emptyset'$-computable function $g : \omega^2 \to 2$ such that for all $m$, we have $m \in f'' \iff \lim g(m, s) = 1$ and $m \not\in f'' \iff \lim g(m, s) = 0$. Notice that for all $e$ and $i \in \omega$, $R_{e,i}$ is satisfied if and only if $\lim g(m_{e,i}, s) = 1$.

At stage $s$, assume we have defined $c(u)$ for every $u < s$. If there exists a least strategy $\mathcal{R}_{e,i}$ (in priority order) with $\langle e, i \rangle < \langle e', i' \rangle$ such that $g(m_{e,i}, s) = 0$, set $c(s) = i$. Otherwise set $c(s) = 0$. This ends the construction. We now turn to the verification.

**Claim.** Every requirement $\mathcal{R}_{e,i}$ is satisfied.

**Proof.** By induction over ordered pairs $\langle e, i \rangle$ in lexicographic order. Suppose that $R_{e',i'}$ is satisfied for all $\langle e', i' \rangle < \langle e, i \rangle$, but $\mathcal{R}_{e,i}$ is not satisfied. Then there exists a threshold $t \geq \langle e, i \rangle$ such that $g(m_{e',i'}, s) = 1$ for all $\langle e', i' \rangle < \langle e, i \rangle$ and $g(m_{e,i}, s) = 0$ whenever $s \geq t$. By construction, $c(s) = i$ for every $s \geq t$. As $Z_e$ is infinite, there exists an element
s ∈ Z such that c(s) = i, so Z is not thin for c with witness i and therefore R_e,i is satisfied. Contradiction.

Proof of Lemma 4.9. Again, we prove it in the case when X = ∅. For each e ∈ ω, let Z_e = {a ∈ ω : f(e, a) = 1}. The proof is very similar to [20, Theorem 5.4.2.]. We build a ∆^0_2 set U together with a stable computable linear order L such that U is the ω part of L, that is, U is the collection of elements L-below cofinitely many other elements. We furthermore ensure that for each e ∈ ω, if Z_e is infinite, then it intersects both U and U. Therefore, if Z_e is infinite, it is neither an ascending, nor a descending sequence in L as otherwise it would be included in either U or U.

Assume by Kleene’s fixpoint theorem that we are given the Turing index d of U as computed relative to ϕ. The set U is built by a finite injury priority construction with the following requirements for each e ∈ ω:

• R_{2e} : Z_e is finite or (∃a)(f(e, a) = 1 and Φ^ϕ_m(a) = 1)
• R_{2e+1} : Z_e is finite or (∃a)(f(e, a) = 1 and Φ^ϕ_m(a) = 0)

Notice again that each of these requirement is Σ^0_2, and furthermore we can effectively find an index for each as such. Therefore, for each i ∈ ω, we can effectively find an m_i such that R_i is satisfied if and only if m_i ∈ f”. By two applications of Schoenfield’s limit Lemma and low_2-ness of f, there exists a computable function g : ω^3 → 2 such that for all m ∈ ω, we have m ∈ f” ⇐ lim, lim, g(m, s, t) = 1 and m ∉ f” ⇐ lim, lim, g(m, s, t) = 0. Notice that for all i ∈ ω,

R_i is satisfied ⇐ lim lim g(m_i, s, t) = 1

At stage 0, U_0 = ∅ and every integer is a leader and follows itself. We say that R_i requires attention for u at stage s if i ≤ u ≤ s, u is leader and g(m_i, s, u) = 0. At stage s + 1, assume we have decided u <_L v or u >_L v for every u, v < s. Set u <_L s if u ∈ U_s and u >_L s if u ∉ U_s. Initially set U_{s+1} = U_s. For each leader u ≤ s which has not been claimed at stage s + 1 and for which some requirement R_i, i < u requires attention, say that the least such R_i claims u and act as follows.

(a) If i = 2e and u ∉ U_s, then add [u, s] to U_{s+1}, where the interval [u, s] is taken in the usual order on ω and not in <_L. Elements of [u+1, s] follow u and are no more considered as leaders from now on and at any further stage.
(b) If i = 2e + 1 and u ∈ U_s, then remove [u, s] from U_{s+1}. Similarly, elements of [u+1, s] are no more leaders and follow u.

Then go to the next leader u ≤ s. This ends the construction. An immediate verification shows that at every stage,

• if u stops being a leader it never becomes again a leader
• if u follows v then v ≤ u, u is a leader, every w between v and u follows v and thus u will never follow any w > v.

So the leader that u follows eventually stabilizes. Moreover, because g is limit-computable, each leader eventually stops increasing the number of followers and therefore there are infinitely many leaders.

Claim. L is a linear order.

Proof. As L is a tournament, it suffices to check there is no 3-cycle. By symmetry, we check only the case where u <_L s <_L v <_L u forms a 3-cycle with s the maximal element
in $<_{\omega}$ order. By construction, this means that $u \in U_s$, $v \not\in U_s$. If $u <_{\omega} v$, then $u \not\in U_v$ and so there exists a leader $w \leq_{\omega} u$ and an even number $i \leq w$ such that $R_i$ requires attention for $w$ at a stage $t \geq v$. Case (a) of the construction applies and the interval $[w+1, t]$ is included $U$ at least until stage $s$. As $v \in [w+1, t]$, $v \in U_s$ contradicting our hypothesis. Case $u >_{\omega} v$ is symmetric.

\textbf{Claim.} $U$ is $\Delta^0_2$.

\textbf{Proof.} Suppose for the sake of contradiction that there exists a least element $u$ entering $U$ and leaving it infinitely many times. Such a $u$ must be a leader, otherwise it would not be the least one. Let $R_i$ be the least requirement claiming $u$ infinitely many times. As $\lim_i g(m_i, s, u)$ exists, it will claim $u$ cofinitely many times and therefore $u$ will be in $U$ or in $\overline{U}$ cofinitely many times. Contradiction.

It immediately follows that $L$ is stable.

\textbf{Claim.} Every requirement $R_i$ is satisfied.

\textbf{Proof.} By induction over $R_i$, in priority order. Suppose that $R_j$ is satisfied for all $j < i$, but $R_i$ is not satisfied. Then there exists a threshold $t_0 \geq i$ such that $\lim_i g(m_i, s, t) = 1$ for all $j < i$ and $\lim_i g(m_i, s, t) = 0$ whenever $t \geq t_0$.

Then for every leader $u \geq t_0$, $R_i$ will claim $u$ cofinitely many times, and therefore $u$ will be in $U$ if $i$ is even and in $\overline{U}$ if $i$ is odd. As every element follows the least leader below itself, every $v$ above the least leader greater than $t_0$ will be in $U$ if $i$ is even and in $\overline{U}$ if $i$ is odd. So if $Z_s$ is infinite, there will be such a $v \in Z_s$ satisfying $R_i$. Contradiction.

\section{5. Degrees bounding the Erdős Moser theorem}

Another approach to the strength analysis of Ramsey’s theorem for pairs consists in seeing a coloring $f : [\omega]^2 \to 2$ as an infinite tournament $T$ such that $T(x, y)$ holds for $x < y$ if and only if $f(x, y) = 1$. The Erdős Moser theorem states the existence of an infinite transitive subtournament, that is, an infinite subset on which the tournament behaves like a linear order. Therefore the Erdős Moser theorem can be seen as a principle reducing instances of $\text{RT}^2_2$ into instances of $\text{ADS}$.

\textbf{Definition 5.1} (Erdős Moser theorem) A tournament $T$ on a domain $D \subseteq N$ is an irreflexive binary relation on $D$ such that for all $x, y \in D$ with $x \neq y$, exactly one of $T(x, y)$ or $T(y, x)$ holds. A tournament $T$ is transitive if the corresponding relation $T$ is transitive in the usual sense. A tournament $T$ is stable if $(\forall x \in D)[(\forall^{\infty} s)T(x, s) \vee (\forall^{\infty} s)T(s, x)]$. $\text{EM}$ is the statement “Every infinite tournament $T$ has an infinite transitive subtournament.” $\text{SEM}$ is the restriction of $\text{EM}$ to stable tournaments.

Bovykin and Weiermann proved in [2] that $\text{EM} + \text{ADS}$ is equivalent to $\text{RT}^2_2$ over $\text{RCA}_0$, equivalence still holding between their stable versions. Lerman et al. [18] proved over $\text{RCA}_0 + \Sigma^0_2$ that $\text{EM}$ implies $\text{OPT}$ and constructed an $\omega$-model of $\text{EM}$ that is not a model of $\text{SRT}^2_2$. Kreuzer proved in [17] that $\text{SEM}$ implies $\Sigma^0_2$ over $\text{RCA}_0$. Bienvenu et al. proved in [1] that $\text{RCA}_0 + \text{SEM} \rightarrow \text{RWKL}$, hence there exists an $\omega$-model of $\text{RRT}^2_2$ that is not a model of $\text{SEM}$. Wang constructed in [27] an $\omega$-model
of EM + COH that is not a model of STS(2). Finally, the author proved in [23] that $\text{RCA}_0 \vdash \text{EM} \rightarrow [\text{STS}(2) \lor \text{COH}]$.

The following notion of minimal interval plays a fundamental role in the analysis of EM. See [18] for a background analysis of EM.

**Definition 5.2 (Minimal interval)** Let $T$ be an infinite tournament and $a, b \in T$ be such that $T(a, b)$ holds. The interval $(a, b)$ is the set of all $x \in T$ such that $T(a, x)$ and $T(x, b)$ hold. Let $F \subseteq T$ be a finite transitive subtournament of $T$. For $a, b \in F$ such that $T(a, b)$ holds, we say that $(a, b)$ is a minimal interval of $F$ if there is no $c \in F \cap (a, b)$, i.e., no $c \in F$ such that $T(a, c)$ and $T(c, b)$ both hold.

We provide in the next subsections two different proofs of the existence of a low$_2$ degree bounding EM. More precisely, we construct a low$_2$ set $G$ which is, up to finite changes, transitive for every infinite computable tournament.

The author proved in [23] that $[\text{STS}(2) \lor \text{COH}] \leq_c \text{EM}$. Therefore every low$_2$ degree bounding EM bounds also COH. The proof does not seem adaptable to prove that COH is a consequence of EM even in $\omega$-models. However we can prove a weaker statement:

**Lemma 5.3** For every set $X$, there exists an infinite $X$-computable tournament $T$ such that for every infinite $T$-transitive subtournament $U$, $U \subseteq^* X$ or $U \subseteq^* \overline{X}$.

**Proof.** Fix a set $X$. We define a tournament $T$ as follows: For each $a < b$, set $T(a, b)$ to hold iff $a \in X$ and $b \in X$ or $a \notin X$ and $b \notin X$. Suppose for the sake of absurd that $U$ is an infinite transitive subtournament of $T$ which intersects infinitely often $X$ and $\overline{X}$. Take any $a, c \in U \cap X$ and $b, d \in U \cap \overline{X}$ such that $a < b < c < d$. Then $T(a, c)$, $T(c, b)$, $T(b, d)$ and $T(d, a)$ hold contradicting transitivity of $U$. □

Using the previous lemma, the constructed set $G$ must be cohesive and therefore provides another proof of the existence of a low$_2$ cohesive set. Finally, we can deduce a statement slightly weaker than Theorem 4.10 simply by the existence of a low$_2$ degree bounding EM.

**Lemma 5.4** There exists a set $C$ such that there is no low$_2$ over $C$ degree $d \gg_{\text{SADS}} C$.

**Proof.** Fix a low$_2$ set $C \gg_{\text{EM}} \emptyset$ and a set $X$ low$_2$ over $C$. By low$_2$-ness of $C$, $X$ is low$_2$. Consider the stable coloring $f : [\omega]^2 \rightarrow 2$ constructed by Mileti in [20, Corollary 5.4.5], such that $X$ computes no infinite $f$-homogeneous set. We can see $f$ as a stable tournament $T$ such that for each $x < y$, $T(x, y)$ holds iff $f(x, y) = 1$. As $C \gg_{\text{EM}} \emptyset$, there exists an infinite $C$-computable transitive subtournament $U$ of $T$. $U$ is a stable linear order such that every infinite ascending or descending sequence is $f$-homogeneous. Therefore $X$ computes no infinite ascending or descending sequence in $U$. □

The following question remains open:

**Question 5.5** Does EM admit a universal instance?

5.1. A low$_2$ degree bounding EM using first jump control. The following theorem uses the proof techniques introduced in [4] for producing low$_2$ sets by controlling the first jump. It is done in the same spirit as Theorem 3.6 in [4].
Theorem 5.6 For every set $P \gg \emptyset$, there exists a set $G \gg_{EM} \emptyset$ such that $G' \leq_{T} P$.

Before proving Theorem 5.6, we introduce the notion of Erdős Moser condition.

Definition 5.7 An Erdős Moser condition (EM condition) for an infinite tournament $T$ is a Mathias condition $(F, X)$ where

(a) $F \cup \{x\}$ is $T$-transitive for each $x \in X$
(b) $X$ is included in a minimal $T$-interval of $F$.

Extension is usual Mathias extension. EM conditions have good properties for tournaments as stated by the following lemmas. Given a tournament $T$ and two sets $E$ and $F$, we denote by $E \rightarrow_{T} F$ the formula $(\forall x \in E)(\forall y \in F)T(x, y)$ holds.

Lemma 5.8 Fix an EM condition $(F, X)$ for a tournament $T$. For every $x \in F$, $\{x\} \rightarrow_{T} X$ or $X \rightarrow_{T} \{x\}$.

Proof. Fix an $x \in F$. Let $(u, v)$ be the minimal $T$-interval containing $X$, where $u, v$ may be respectively $-\infty$ and $+\infty$. By definition of interval, $\{u\} \rightarrow_{T} X \rightarrow_{T} \{v\}$. By definition of minimal interval, $T(x, u)$ or $T(v, x)$ holds. Suppose the former holds. By transitivity of $F \cup \{y\}$ for every $y \in X$, $T(x, y)$ holds, therefore $\{x\} \rightarrow_{T} X$. In the latter case, by symmetry, $X \rightarrow_{T} \{x\}$.

Lemma 5.9 Fix an EM condition $c = (F, X)$ for a tournament $T$, an infinite subset $Y \subseteq X$ and a finite $T$-transitive set $F_{1} \subseteq X$ such that $F_{1} < Y$ and $[F_{1} \rightarrow_{T} Y \lor Y \rightarrow_{T} F_{1}]$. Then $d = (F \cup F_{1}, Y)$ is a valid extension of $c$.

Proof. Properties of a Mathias condition for $d$ are immediate. We prove property (a). Fix an $x \in Y$. To prove that $F \cup F_{1} \cup \{x\}$ is $T$-transitive, it suffices to check that there exists no 3-cycle in $F \cup F_{1} \cup \{x\}$. Fix three elements $u < v < w \in F \cup F_{1} \cup \{x\}$.

- Case 1: $\{u, v, w\} \cap F \neq \emptyset$. Then $u \in F$ as $F < F_{1} < \{x\}$ and $u < v < w$. If $v \in F$ then using the fact that $F_{1} \cup \{x\} \subseteq X$ and property (a) of condition $c$, $\{u, v, w\}$ is $T$-transitive. If $v \notin F$, then by Lemma 5.8, $\{u\} \rightarrow_{T} X(\supseteq F \cup \{x\})$ or $X \rightarrow_{T} \{u\}$ hence $\{u\} \rightarrow_{T} \{v, w\}$ or $\{v, w\} \rightarrow_{T} \{u\}$ so $\{u, v, w\}$ is $T$-transitive.
- Case 2: $\{u, v, w\} \cap F = \emptyset$. Then at least $u, v \in F_{1}$ because $F_{1} < \{x\}$. If $w \in F_{1}$, then $\{u, v, w\}$ is $T$-transitive by $T$-transitivity of $F_{1}$. Otherwise, as $F_{1} \rightarrow_{T} Y$ or $Y \rightarrow_{T} F_{1}$, $\{u, v\} \rightarrow_{T} \{w\}$ or $\{w\} \rightarrow_{T} \{u, v\}$ and $\{u, v, w\}$ is $T$-transitive.

We now prove property (b). Let $(u, v)$ be the minimal $T$-interval of $F$ in which $X$ (hence $Y$) is included by property (b) of condition $c$. $u$ and $v$ may be respectively $-\infty$ and $+\infty$. By assumption, either $F_{1} \rightarrow_{T} Y$ or $Y \rightarrow_{T} F_{1}$. As $F_{1}$ is a finite $T$-transitive set, it has a minimal and a maximal element, say $x$ and $y$. If $F_{1} \rightarrow_{T} Y$ then $Y$ is included in the $T$-interval $(y, v)$. Symmetrically, if $y \rightarrow_{T} F_{1}$ then $Y$ is included in the $T$-interval $(u, x)$. To prove minimality for the first case, assume that some $w$ is in the interval $(y, v)$. Then $w \notin F$ by minimality of the interval $(u, v)$ w.r.t. $F$, and $w \notin F_{1}$ by maximality of $y$. Minimality for the second case holds by symmetry.

Proof of Theorem 5.6. Let $C$ be a low set such that there exists a uniformly $C$-computable enumeration $\mathcal{F}$ of infinite tournaments containing every computable tournament. Note that $P \gg C'$. Our forcing conditions are tuples $(\sigma, F, X)$ where $\sigma \in \omega^{<\omega}$ and the following holds:
(a) \((F,X)\) forms a Mathias condition and \(X\) is a set low over \(C\).
(b) \((F \setminus [0, \sigma(v)], X)\) is an EM condition for \(T_v\) for each \(v < |\sigma|\).

A condition \((\tilde{\sigma}, \tilde{F}, \tilde{X})\) extends a condition \((\sigma, F, X)\) if \(\sigma \leq \tilde{\sigma}\) and \((\tilde{F}, \tilde{X})\) Mathias extends \((F,X)\). A set \(G\) satisfies the condition \((\sigma, F, X)\) if \(G \setminus [0, \sigma(v)]\) is \(T_v\)-transitive for each \(v < |\sigma|\) and \(G\) satisfies the Mathias condition \((F,X)\). An index of a condition \((\sigma, F, X)\) is a code of the tuple \(\langle \sigma, F, e \rangle\) where \(e\) is a lowness index of \(X\).

The first lemma simply states that we can ensure that \(G\) will be infinite and eventually transitive for each tournament in \(\tilde{T}\).

**Lemma 5.10** For every condition \(c = (\sigma, F, X)\) and every \(i, j \in \omega\), one can \(P\)-compute an extension \((\tilde{\sigma}, \tilde{F}, \tilde{X})\) such that \(|\tilde{\sigma}| \geq i\) and \(|\tilde{F}| \geq j\) uniformly from \(i, j\) and an index of \(c\).

**Proof.** Let \(x\) be the first element of \(X\). As \(X\) is low over \(C\), \(x\) can be found \(C\')-computably from a lowness index of \(X\). The condition \((\tilde{\sigma}, \tilde{F}, \tilde{X})\) is a valid extension of \(c\) where \(\tilde{\sigma} = \sigma \downarrow x \ldots x\) so that \(|\tilde{\sigma}| \geq i\). It suffices to prove that we can \(C\')\)-compute an extension \((\tilde{\sigma}, \tilde{F}, \tilde{X})\) with \(|\tilde{F}| > |F|\) and iterate the process. Define the computable coloring \(g : X \rightarrow 2^{[\sigma]}\) by \(g(s) = \rho\) where \(\rho \in \dot{2}^{[\sigma]}\) such that \(\rho(v) = 1\) iff \(T_v(x, s)\) holds. One can find uniformly in \(P\) a \(\rho \in 2^{[\sigma]}\) such that the following \(C\)-computable set is infinite:

\[Y = \{s \in X \setminus \{x\} : g(s) = \rho\}\]

By Lemma 5.9, \(\{(F \cup \{x\}) \setminus [0, \sigma(v)], Y\}\) is a valid EM extension for \(T_v\). As \(Y\) is low over \(C\), \((\tilde{\sigma}, F \cup \{x\}, Y)\) is a valid extension for \(c\).

It remains to be able to decide \(e \in (G \oplus C)'\) uniformly in \(e\). We first need to define a forcing relation.

**Definition 5.11** Fix a condition \(c = (\sigma, F, X)\) and two integers \(e\) and \(x\).

1. \(c \Vdash \Phi_e^{G \oplus C}(x) \uparrow\) if \(\Phi_e^{(F \cup \{x\}) \oplus C}(x) \uparrow\) for all finite subsets \(F_1 \subseteq X\) such that \(F_1\) is \(T_v\)-transitive simultaneously for each \(v < |\sigma|\).
2. \(c \Vdash \Phi_e^{G \oplus C}(x) \downarrow\) if \(\Phi_e^{F \oplus C}(x) \downarrow\).

Note that the way we defined our forcing relation \(c \Vdash \Psi_e^{G \oplus C}(x) \uparrow\) differs slightly from the “true” forcing notion \(\Vdash^*\) inherited by the notion of satisfaction of \(G\). The true forcing definition of this statement is the following:

\(c \Vdash^* \Phi_e^{G \oplus C}(x) \uparrow\) if \(\Phi_e^{(F \cup \{x\}) \oplus C}(x) \uparrow\) for all finite extensible subsets \(F_1 \subseteq X\) such that \(F_1\) is \(T_v\)-transitive simultaneously for each \(v < |\sigma|\), i.e., for all finite subsets \(F_1 \subseteq X\) such that there exists an extension \(d = (\tilde{\sigma}, F \cup F_1, \tilde{X})\).

However \(c \Vdash^* \Phi_e^{G \oplus C}(x) \uparrow\) is not a \(\Pi^0_1\) statement whereas \(c \Vdash \Phi_e^{G \oplus C}(x) \uparrow\) is. In particular the fact that \(c \not\Vdash \Phi_e^{G \oplus C}(x) \uparrow\) does not mean that \(c\) has an extension forcing its negation. This subtlety is particularly important in Lemma 5.13. The following lemma gives a sufficient constraint, namely being included in a part of a particular partition, on finite transitive sets to ensure that they are extensible.

**Lemma 5.12** Let \(c = (\sigma, F, X)\) be a condition and \(E \subseteq X\) be a finite set. There exists a \(2^{[\sigma]}\) partition \((E_\rho : \rho \in 2^{[\sigma]}\) of \(E\) and an infinite set \(Y \subseteq X\) low over \(C\) such that \(E < Y\) and for all \(\rho \in 2^{[\sigma]}\) and \(v < |\sigma|\), if \(\rho(v) = 0\) then \(E_\rho \rightarrow_{T_v} Y\) and if \(\rho(v) = 1\) then \(Y \rightarrow_{T_v} E_\rho\).
Moreover this partition and a lowness index of $Y$ can be uniformly $P$-computed from an index of $c$ and the set $E$.

**Proof.** Given a set $E$, define $P_e$ to be the finite set of ordered $2^{[\sigma]}$-partitions of $E$, that is,

$$P_e = \{ (E_{\rho} : \rho \in 2^{[\sigma]} ) : \bigcup_{\rho \in 2^{[\sigma]}} E_{\rho} = E \text{ and } \rho \neq \xi \rightarrow E_{\rho} \cap E_{\xi} = \emptyset \}$$

Define the $C$-computable coloring $g : X \rightarrow P_e$ by $g(x) = (E^x_{\rho} : \rho \in 2^{[\sigma]})$ where

$$E^x_{\rho} = \{ a \in E : (\forall v < |\sigma|)[T_e(a,x) \text{ holds iff } \rho(v) = 0] \}.$$ 

On can find uniformly in $P$ a partition $(E_{\rho} : \rho \in 2^{[\sigma]})$ such that the following $C$-computable set is infinite:

$$Y = \{ x \in X \setminus E : g(x) = (E_{\rho} : \rho \in 2^{[\sigma]}) \}$$

By definition of $g$, for all $\rho \in 2^{[\sigma]}$ and $v < |\sigma|$, if $\rho(v) = 0$ then $E_{\rho} \rightarrow_T Y$ and if $\rho(v) = 1$ then $Y \rightarrow_T E_{\rho}$.

We are now ready to prove the key lemma of this forcing, stating that we can $P$-decide whether or not $e \in G'$ for any $e \in \omega$.

**Lemma 5.13** For every condition $(\sigma, F, X)$ and every $e \in \omega$, there exists an extension $d = (\mathcal{Q}, F, X)$ such that one of the following holds:

1. $d \models \Phi^{G \& C}(e) \downarrow$
2. $d \models \Phi^{G \& C}(e) \uparrow$

This extension can be $P$-computed uniformly from an index of $c$ and $e$. Moreover there is a $C'$-computable procedure to decide which case holds from an index of $d$.

**Proof.** Let $k = |\sigma|$. Using a $C'$-computable procedure, we can decide from an index of $c$ and $e$ whether there exists a finite set $E \subseteq X$ such that for every $2^k$-partition $(E_i : i < 2^k)$ of $E$, there exists an $i < 2^k$ and a subset $F_1 \subseteq E_i$ $T_e$-transitive simultaneously for each $v < k$ and satisfying $\Phi^{(\mathcal{Q}, F, X)}_{\epsilon}(e) \downarrow$.

1. If such a set $E$ exists, it can be $C'$-computably found. By Lemma 5.12, one can $P$-computably find a $2^k$-partition $(E_{\rho} : \rho \in 2^k)$ of $E$ and a set $Y \subseteq X$ low over $C$ such that for all $\rho \in 2^k$ and $v < k$, if $\rho(v) = 0$ then $E_{\rho} \rightarrow_T Y$ and if $\rho(v) = 1$ then $Y \rightarrow_T E_{\rho}$. We can $C'$-computably find a $\rho \in 2^k$ and a set $F_1 \subseteq E_\rho$ which is $T_e$-transitive simultaneously for each $v < k$ and satisfying $\Phi^{(\mathcal{Q}, F, X)}_{\epsilon}(e) \downarrow$. By Lemma 5.9, $(F \setminus \{0, \sigma(v)\}) \cup F_1, Y)$ is a valid EM extension of $(\mathcal{Q}, F, X)$ for each $v < k$. As $Y$ is low over $C$, $(\sigma, F \cup F_1, Y)$ is a valid extension of $c$ forcing $\Phi^{G \& C}(e) \downarrow$.

2. If no such set exists, then by compactness, the $\Pi^{0}_{1,C}$ class of all $2^k$-partitions $(X_i : i < 2^k)$ of $X$ such that for every $i < 2^k$ and every finite set $F_1 \subseteq X_i$ which is $T_e$-transitive simultaneously for each $v < k$, $\Phi^{(\mathcal{Q}, F, X)}_{\epsilon}(e) \uparrow$ is non-empty. In other words, the $\Pi^{0}_{1,C}$ class of all $2^k$-partitions $(X_i : i < 2^k)$ of $X$ such that for every $i < 2^k$, $(\sigma, F, X_i) \models \Phi^{G \& C}(e) \uparrow$ is non-empty. By the relativized low basis theorem, there exists a $2^k$-partition $(X_i : i < 2^k)$ of $X$ low over $C$. Furthermore, a lowness index for this partition can be uniformly $C'$-computably found. Using $P$, one can find an $i < 2^k$ such that $X_i$ is infinite. $(\sigma, F, X_i)$ is a valid extension of $c$ forcing $\Phi^{G \& C}(e) \uparrow$.
Using Lemma 5.10 and Lemma 5.13, one can $P$-compute an infinite decreasing sequence of conditions $c_0 = (e, 0, \omega) \geq c_1 \geq \ldots$ such that for each $s > 0$

1. $|\sigma_i| \geq s$, $|F_i| \geq s$
2. $c_i \models \Phi^G(c) \downarrow$ or $c_i \models \Phi^G(c) \uparrow$

where $c_i = (\sigma_i, F_i, X_i)$. The resulting set $G = \bigcup_s F_s$ is $T_v$-transitive up to finite changes for each $v \in \omega$ and $G \leq_T P$.

\[ \Box \]

5.2. A low, degree bounding EM using second jump control. We now use the second proof technique used in [4] for producing a low$_2$ set. It consists of directly controlling the second jump of the produced set.

Theorem 5.14 There exists a low$_2$ degree bounding EM.

Proof. Similar to Theorem 5.6, we fix a low set $C$ such that there exists a uniformly $\hat{C}$-computable enumeration $\check{T}$ of infinite tournaments containing every computable tournament. In particular $P \gg C'$.

Our forcing conditions are the same as in Theorem 5.6. We can release the constraints of infinity and lowness over $C$ for $X$ in a condition $(\sigma, F, X)$. This gives the notion of precondition. The forcing relations extend naturally to preconditions.

Definition 5.15 Fix a finite set of Turing indexes $\check{\varepsilon}$. A condition $(\sigma, F, X)$ is $\check{\varepsilon}$-small if there exists a number $x$ and a sequence $(\sigma_i, F_i, X_i : i < n)$ such that for each $i < n$

1. $(\sigma_i, F_i, X_i)$ is a precondition extending $c$
2. $(X_i : i < n)$ is a partition of $X \cap (x, +\infty)$
3. $\max(X_i) < x$ or $(\sigma_i, F \cup F_i, X_i) \models (\exists e \in \check{\varepsilon})(\exists y < x)\Phi^G(c)(y) \uparrow$

A condition is $\check{\varepsilon}$-large if it is not $\check{\varepsilon}$-small.

A condition $(\check{\sigma}, \check{F}, \check{X})$ is a finite extension of $(\sigma, F, X)$ if $\check{X} =^* X$. Finite extensions do not play the same fundamental role as in the original forcing in [4] as adding elements to the set $F$ may require to remove infinitely many elements of the promise set $X$ to obtain a valid extension. We nevertheless prove the following traditional lemma.

Lemma 5.16 Fix an $\check{\varepsilon}$-large condition $c = (\sigma, F, X)$.

1. If $\check{\varepsilon}' \subseteq \check{\varepsilon}$ then $c$ is $\check{\varepsilon}'$-large.
2. If $d$ is a finite extension of $c$ then $d$ is $\check{\varepsilon}$-large.

Proof. Clause 1 is trivial as $\check{\varepsilon}$ appears only in a universal quantification in the definition of $\check{\varepsilon}$-largeness. We prove clause 2. Let $d = (\check{\sigma}, \check{F}, \check{X})$ be an $\check{\varepsilon}$-small finite extension of $c$. We will prove that $c$ is $\check{\varepsilon}$-small. Let $x \in \omega$ and $(\sigma_i, F_i, X_i : i < n)$ witness $\check{\varepsilon}$-smallness of $d$. Let $y = \max(x, X \setminus \check{X})$. For each $i < n$, set $\check{X}_i = X_i \cap (y, +\infty)$. Then $y$ and $(\sigma_i, F_i, \check{X}_i : i < n)$ witness $\check{\varepsilon}$-smallness of $c$.

\[ \Box \]

Lemma 5.17 There exists a $C'$-effective procedure to decide, given an index of a condition $c$ and a finite set of Turing indexes $\check{\varepsilon}$, whether $c$ is $\check{\varepsilon}$-large. Furthermore, if $c$ is $\check{\varepsilon}$-small, there exists sets $(X_i : i < n)$ low over $C$ witnessing this, and one may $C'$-compute a value of $n$, $x$, lowness indexes for $(X_i : i < n)$ and the corresponding sequences $(\sigma_i, F_i, X_i : i < n)$ which witness that $c$ is $\check{\varepsilon}$-small.
Proof. Fix a condition \( c = (\sigma, F, X) \) The predicate “\((\sigma, F, X)\) is \( \vec{c} \)-small” can be expressed as a \( \Sigma^0_2 \) statement
\[
(3z)(3Z)P(z, Z, F, X, \vec{c}, \vec{\varepsilon})
\]
where \( P \) is a \( \Pi^0_1 \) predicate. Here \( z \) codes \( n \) and \( x \), and \( Z \) codes \( (X_i : i < n) \). The predicate \( (3Z)P(z, Z, F, X, \sigma, \vec{\varepsilon}) \) is \( \Pi^0_1 \) by compactness. As \( X \) is low over \( C \) and \( F \) and \( \sigma \) are finite, one can compute a \( \Delta^0_2 \) index for the same predicate \( P \) with parameter \( z \), an index of \( c \) and \( \vec{\varepsilon} \), from a lowness index for \( X, F \) and \( \sigma \). Therefore there exists a \( \Sigma^0_2 \) statement with parameters an index of \( c \) and \( \vec{\varepsilon} \) which holds iff \( c \) is \( \vec{c} \)-small.

If \( c \) is \( \vec{c} \)-small, there exists sets \( (X_i : i < n) \) low over \( X \) (hence low over \( C \)) witnessing it by the low basis theorem relativized to \( C \). By the uniformity of the proof of the low basis theorem, one can compute lowness indexes of \( (X_i : i < n) \) uniformly from a lowness index of \( X \).

As the extension produced in Lemma 5.10 is not a finite extension, we need to refine it to ensure largeness preservation.

Lemma 5.18 For every \( \vec{c} \)-large condition \( c = (\sigma, F, X) \) and every \( i, j \in \omega \), one can \( P \)-compute an \( \vec{c} \)-large extension \( (\tilde{\sigma}, \tilde{F}, \tilde{X}) \) such that \( \tilde{\sigma} \geq i \) and \( |\tilde{F}| \geq j \) uniformly from an index of \( c \), \( i \), \( j \) and \( \vec{c} \).

Proof. Let \( x \) be the first element of \( X \). As \( X \) is low over \( C \), \( x \) can be found \( C^\prime \)-computably from a lowness index of \( X \). The condition \( d = (\tilde{\sigma}, F, X) \) is a valid extension of \( c \) where \( \tilde{\sigma} = \sigma \times \ldots \times x \) so that \( |\tilde{\sigma}| \geq i \). As \( d \) is a finite extension of \( c \), it is \( \vec{c} \)-large by Lemma 5.16. It suffices to prove that we can \( C^\prime \)-compute an \( \vec{c} \)-large extension \( (\tilde{\sigma}, \tilde{F}, \tilde{X}) \) with \( |\tilde{F}| > |F| \) and iterate the process. Define the \( C \)-computable coloring \( g : X \to 2^{[\sigma]} \) as in Lemma 5.10. For each \( \rho \in 2^{[\sigma]} \), define the following set:
\[
Y_\rho = \{ s \in X \setminus \{ x \} : g(s) = \rho \}
\]
There must be a \( \rho \in 2^{[\sigma]} \) such that \( Y_\rho \) is infinite and \( (\tilde{\sigma}, F \cup \{ x \}, Y_\rho) \) is \( \vec{c} \)-large, otherwise the witnesses of \( \vec{c} \)-smallness for each \( \rho \in 2^{[\sigma]} \) would witness \( \vec{c} \)-smallness of \( c \). By Lemma 5.17, one can \( C^\prime \)-find a \( \rho \in 2^{[\sigma]} \) such that \( (\tilde{\sigma}, F \cup \{ x \}, Y_\rho) \) is \( \vec{c} \)-large. As seen in Lemma 5.10, \( (\tilde{\sigma}, F, \{ x \}, Y_\rho) \) is a valid extension.

The following lemma is a refinement of Lemma 5.12 controlling largeness preservation.

Lemma 5.19 Let \( c = (\sigma, F, X) \) be an \( \vec{c} \)-large condition and \( E \subseteq X \) be a finite set. There is a \( 2^{[\sigma]} \) partition \( (E_\rho : \rho \in 2^{[\sigma]}) \) of \( E \) and an infinite set \( Y \subseteq X \) low over \( C \) such that \( E \subseteq Y \) and

1. for all \( \rho \in 2^{[\sigma]} \), if \( \rho(v) = 0 \) then \( E_\rho \to_{T_v} Y \) and if \( \rho(v) = 1 \) then \( Y \to_{T_v} E_\rho \).
2. \( (\sigma, F \cup F_1, Y) \) is an \( \vec{c} \)-large condition extending \( c \) for every \( \rho \in 2^{[\sigma]} \) and every finite set \( F_1 \subseteq E_\rho \), which is \( T_v \)-transitive for each \( v < |\sigma| \).

Moreover this partition and a lowness index of \( Y \) can be uniformly \( C^\prime \)-computed from an index of \( c \) and the set \( E \).

Proof. Given a set \( E \), recall from Lemma 5.12 that \( P_E \) is the finite set of ordered \( 2^k \)-partitions of \( E \). Define again the computable coloring \( g : X \to P_E \) by \( g(x) = (E^x_\rho : \rho \in 2^{[\sigma]}) \) where \( E^x_\rho = \{ a \in E : (\forall v < |\sigma|)[T_v(a, x) \text{ holds iff } \rho(v) = 0] \} \). If for each
Lemma 5.21 Suppose that \((T, \rho) \in 2^{[\omega]}\) and a \(F_1 \subseteq E_\rho\) which is \(T,\)-transitive simultaneously for each \(v < |\sigma|\) and such that \((\sigma, F \cup F_1, Y)\) is \(\bar{\epsilon}\)-small where
\[
Y = \{x \in X \setminus E : g(x) = (E_\rho : \rho \in 2^{[\omega]})\}
\]
Then we could construct a witness of \(\bar{\epsilon}\)-smallness of \(c\) using smallness witnesses of \((\sigma, F \cup F_1, Y)\) for each partition \((E_\rho : \rho \in 2^{[\omega]})\). Therefore there must exist a partition \((E_\rho : \rho \in 2^{[\omega]})\) such that \(Y\) is infinite and \(d = (\sigma, F \cup F_1, Y)\) is \(\bar{\epsilon}\)-large for every \(\rho \in 2^{[\omega]}\) and every \(F_1 \subseteq E_\rho\) which is \(T,\)-transitive for each \(v < |\sigma|\).

By Lemma 5.17, such a partition can be found \(\bar{C}^{\prime}\)-computably. By definition of \(\epsilon\), for all \(\rho \in 2^{[\omega]}\) and \(v < k\), if \(\rho(v) = 0\) then \(E_\rho \rightarrow_{T, v} Y\) and if \(\rho(v) = 1\) then \(Y \rightarrow_{T, v} E_\rho\). Therefore, by Lemma 5.9, \((F \setminus [0, \sigma(v)]) \cup F_1, Y)\) is a valid EM extension of \((F \setminus [0, \sigma(v)], X)\) for \(T,\) for each \(v < |\sigma|\), so \(d\) is a valid condition. 

Lemma 5.20 Suppose that \(c = (\sigma, F, X)\) is \(\bar{\epsilon}\)-large. For every \(y \in \omega\) and \(e \in \bar{\epsilon}\), there exists an \(\bar{\epsilon}\)-large extension \(d\) such that \(d \models \Phi^{G\bar{\epsilon}C}(y)\). Furthermore, an index for \(d\) can be computed from an oracle for \(C\) from an index of \(c\), \(e\) and \(y\).

**Proof.** Let \(k = |\sigma|\). As \(c\) is \(\bar{\epsilon}\)-large, then by a compactness argument, there exists a finite set \(E \subseteq X\) such that for every \(2^k\)-partition \((E_i : i < 2^k)\) of \(E\), there exists an \(i < k\) and a finite subset \(F_i \subseteq E_i\) which is \(T,\)-transitive simultaneously for each \(v < k\), and \(\Phi^{(F \cup F_i)\bar{\epsilon}C}(y)\). Moreover this set \(E\) can be \(C\)-computably found. By Lemma 5.19, on can uniformly \(\bar{C}^{\prime}\)-find a partition \((E_\rho : \rho \in 2^k)\) of \(E\) and a lowness index for an infinite set \(Y \subseteq X\) low over \(C\) such that

1. for all \(\rho \in 2^k\) and \(v < k\), if \(\rho(v) = 0\) then \(E_\rho \rightarrow_{T, v} Y\) and if \(\rho(v) = 1\) then \(Y \rightarrow_{T, v} E_\rho\).
2. \((\sigma, F \cup F_1, Y)\) is an \(\bar{\epsilon}\)-large condition extending \(c\) for every \(\rho \in 2^k\) and every finite set finite set \(F_1 \subseteq E_\rho\) which is \(T,\)-transitive for each \(v < k\).

We can then produce by a \(C\)-computable search a \(\rho \in 2^k\) and a finite set \(F_i \subseteq E_\rho\) which is \(T,\)-transitive for each \(v < k\) and such that \(\Phi^{(F \cup F_i)\bar{\epsilon}C}(y)\). By Lemma 5.9, \((F \setminus [0, \sigma(v)]) \cup F_1, Y)\) is a valid EM extension of \((F \setminus [0, \sigma(v)], X)\) for \(T,\) for each \(v < k\). As \(Y\) is low over \(C\), \((\sigma, F \cup F_1, Y)\) is a valid \(\bar{\epsilon}\)-large extension. 

Lemma 5.21 Suppose that \(c = (\sigma, F, X)\) is \(\bar{\epsilon}\)-large and \((\bar{\epsilon} \cup \{u\})\)-small. There exists a \(\bar{\epsilon}\)-large extension \(d\) such that \(d \models \Phi^{G\bar{\epsilon}C}(y)\) for some \(y \in \omega\). Furthermore one can find an index for \(d\) by applying a \(C\)-computable function to an index of \(c\), \(\bar{\epsilon}\) and \(u\).

**Proof.** By Lemma 5.17, we may choose the sets \((X_i : i < n)\) witnessing that \(c \in (\bar{\epsilon} \cup \{u\})\)-small to be low over \(C\). Fix the corresponding \(x\) and \((\sigma_i, F_i : i < n)\). Consider the \(i\)'s such that \((\sigma_i, F_i, X_i) \models \Phi^{G\bar{\epsilon}C}(y)\) for some \(y < x\). As \(c\) is \(\bar{\epsilon}\)-large, there must be such an \(i < n\) such that \((\sigma_i, F_i, X_i)\) is an \(\bar{\epsilon}\)-large condition. By Lemma 5.17 we can find \(\bar{C}^{\prime}\)-computably such an \(i < n\). \((\sigma_i, F_i, X_i)\) is the desired extension.

Using the previous lemmas, we can \(\bar{C}^{\prime}\)-compute an infinite descending sequence of conditions \(c_0 = (\epsilon, \emptyset, \omega) \geq c_1 \geq \ldots\) together with an infinite increasing sequence of Turing indices \(\bar{c}_0 = \emptyset \leq \bar{c}_1 \leq \ldots\) such that for each \(s > 0\)

1. \(|\sigma_i| \geq s, |F_i| \geq s, c_i\) is \(\bar{\epsilon}\)-large
2. Either \(s \in \bar{c}_i\) or \(c_i \models \Phi^{G\bar{\epsilon}C}(y)\) for some \(y \in \omega\)
3. \(c_i \models \Phi^{G\bar{\epsilon}C}(x)\) if \(s = (\epsilon, x)\) and \(e \in \bar{\epsilon}_i\)
where \( c = (\sigma, F, X) \). The resulting set \( G = \bigcup F \) is \( T \)-transitive up to finite changes simultaneously for each \( v \in \omega \) and \( G'' \leq_T C'' \leq_T \emptyset'' \).

\[ \square \]

6. Degree bounding the rainbow Ramsey theorem

The rainbow Ramsey theorem intuitively states that when a coloring over tuples uses each color a bounded number of times then it has an infinite subset on which each color is used at most once. This statement has been extensively studied over the past few years [8, 7, 26, 23]. Remarkably, the restriction of the rainbow Ramsey theorem to coloring over pairs of integers coincides with a well-known notion of algorithmic randomness.

**Definition 6.1** (Rainbow Ramsey theorem) Let \( n, k \in \omega \). A coloring function \( f : [\omega]^n \to \omega \) is \( k \)-bounded if for every \( y \in \omega \), \( |f^{-1}(y)| \leq k \). A set \( R \) is a rainbow for \( f \) if \( f(R) \) is injective. \( \text{RRT}_n^k \) is the statement “Every \( k \)-bounded function \( f : [\omega]^n \to \omega \) has an infinite rainbow”.

A proof of the rainbow Ramsey theorem is due to Galvin who noticed that it follows easily from Ramsey’s theorem. Hence every computable 2-colored bounding function \( f \) over \( n \)-tuples has an infinite \( \Pi^0_2 \)-rainbow. Csima and Mileti proved in [8] that every 2-random is \( \text{RRT}_2^2 \)-bounding and deduced that \( \text{RRT}_2^2 \) implies neither SADS nor \( \text{WKL}_0 \) over \( \omega \)-models. Conidis & Slaman adapted in [7] the argument from Csima and Mileti to obtain \( \text{RCA}_0 \vdash \text{2-RAN} \to \text{RRT}_2^2 \).

**Definition 6.2** A function \( f : \omega \to \omega \) is diagonally non-computable (DNC) relative to \( X \) if \( f(e) \neq \Phi^X_e(e) \) for each \( e \in \omega \). \( \text{DNR}[\emptyset'] \) is the statement “For every set \( X \), there exists a function DNC relative to the jump of \( X \)”.

**Theorem 6.3** (J.S. Miller [21]) \( \text{RRT}_2^2 \) and \( \text{DNR}[\emptyset'] \) are computably equivalent.

**Corollary 6.4** \( \text{RRT}_2^2 \) admits a universal instance.

**Proof.** If \( P \) and \( Q \) are two principles computably equivalent and \( Q \) admits a universal instance, then so does \( P \). As \( \text{DNR}[\emptyset'] \) admits a universal instance (any function DNC relative to \( \emptyset' \)), so does \( \text{RRT}_2^2 \). \( \square \)

**Corollary 6.5** For every \( X \gg \emptyset' \), there exists a \( Y \gg_{\text{RRT}_2^2} \emptyset \) such that \( Y' \leq_T X \).

**Proof.** Let \( f : [\omega]^2 \to \omega \) be a universal instance of \( \text{RRT}_2^2 \). By Csima & Mileti [8], \( \text{RRT}_2^2 \leq_{\text{c}} \text{RT}_2^2 \), so there exists a computable coloring \( g : [\omega]^2 \to 2 \) such that every infinite \( g \)-homogeneous set computes an infinite \( f \)-rainbow, hence bounds \( \text{RRT}_2^2 \). By Cholak et al. [4], for every \( X \gg \emptyset' \) there exists an infinite \( g \)-homogeneous set \( H \) such that \( H' \leq_T X \). In particular \( H \gg_{\text{RRT}_2^2} \emptyset \). \( \square \)

**Corollary 6.6** There exists a low_2 degree bounding \( \text{RRT}_2^2 \).

**Proof.** By the relativized low basis theorem, there exists a set \( X \gg \emptyset' \) low over \( \emptyset' \). By Corollary 6.5, there exists a set \( Y \gg_{\text{RRT}_2^2} \emptyset \) such that \( Y' \leq_T X \), hence \( Y'' \leq_T X' \leq_T \emptyset'' \). So \( Y \) is low_2. \( \square \)
We can generalize Corollary 6.6 to colorings over arbitrary tuples. For this, we need to restrict ourselves to the study of a particular class of colorings.

**Definition 6.7** A coloring \( f : [\omega]^{n+1} \to \omega \) is normal if \( f(\sigma, a) \neq f(\tau, b) \) for each \( \sigma, \tau \in [\omega]^n \), whenever \( a \neq b \).

Wang proved in [26] that for every 2-bounded coloring \( f : [\omega]^n \to \omega \), every \( f \)-random computes an infinite set \( X \) on which \( f \) is normal. The author refined in [23] this result by proving that every function d.n.c. relative to \( f \) computes such a set.

**Theorem 6.8** For each \( n \geq 0 \), there exists a set \( X \gg \text{RRT}^2_2 \emptyset \) low over \( \emptyset(n) \).

**Proof.** We prove by induction over \( n \) that for every set \( A \) there exists a set \( X \) low \( 2 \) over \( A(n) \) such that \( X \gg \text{RRT}^2_2 A \). Case \( n = 0 \) is a relativization of Corollary 6.6. Suppose for each set \( A \), there exists a set \( X \) low \( 2 \) over \( A(n) \) such that \( X \gg \text{RRT}^2_2 A \). Fix a set \( A \), an \( A \)-random set \( R \) low over \( A \) and a set \( C \) low \( 2 \) over \( A \oplus R \) such that \( C' \gg (A \oplus R)' \). In particular \( R \oplus C \) is low \( 2 \) over \( A \). By induction hypothesis, there exists a set \( X \) low \( 2 \) over \( (A \oplus R \oplus C)' \) such that \( X \gg \text{RRT}^2_2 (A \oplus R \oplus C)' \). In particular \( X \) is low \( 2 \) over \( A^{n+1} \). We can assume without loss of generality that \( X \) computes \( A \), since \( X \oplus A \) is low \( 2 \) over \( A^{n+1} \) and \( X \oplus A \gg \text{RRT}^2_2 (A \oplus R \oplus C)' \). We claim that \( X \gg \text{RRT}^2_2 A \).

Fix an \( A \)-computable 2-bounded coloring \( f : [\omega]^{n+3} \to \omega \). By relativizing Lemma 4.3 in [26], every \( A \)-random computes an infinite set \( Y \) such that \( f \) restricted to \( Y \) is normal. So \( A \oplus R \) computes such a set \( Y \). For each \( \sigma, \tau \in [Y]^{n+2} \), let

\[
U_{\sigma, \tau} = \{ s \in Y : f(\sigma, s) = f(\tau, s) \}
\]

By Jockusch & Frank [16], as \( C' \gg (A \oplus R)' \), \( A \oplus R \oplus C \) computes an infinite \( U \)-cohesive set \( Z \subseteq Y \). In particular the following limit exists

\[
\hat{f}(\sigma) = \lim_{s \in Z} \min \{ \tau \leq_{\text{lex}} \sigma : f(\sigma, s) = f(\tau, s) \}
\]

\( \hat{f} \) is a 2-bounded \( (A \oplus R \oplus C)' \)-computable coloring of \( (n+2) \)-tuples, so \( X \) bounds an infinite \( \hat{f} \)-rainbow \( H \). \( A \oplus H \) computes an infinite \( f \)-rainbow, so \( X \) bounds an infinite \( f \)-rainbow. \( \square \)

### 6.1. A stable rainbow Ramsey theorem.

A common process in the strength analysis of a principle consists of splitting the statement into a stable and a cohesive version. The standard notion of stability does not apply for the rainbow Ramsey theorem as no stable coloring is \( k \)-bounded for some \( k \in \omega \). Nevertheless one can define certain notions of stability for the rainbow Ramsey theorem [23]. Mileti proved in [20] that the only \( \Delta^0_5 \) degree bounding \( \text{SRT}^2_2 \) is \( \mathcal{O}' \). In fact, his priority argument can be adapted to prove the same result on a much weaker principle coinciding with a stable version of the rainbow Ramsey theorem for pairs.

**Definition 6.9** A coloring \( f : [\omega]^2 \to \omega \) is rainbow-stable if for every \( x \in \omega \), one of the following holds:

(a) There exists a \( y \neq x \) such that \( (\forall \infty)s f(x, s) = f(y, s) \)

(b) \( (\forall \infty s) \{ y \neq x : f(x, s) = f(y, s) \} \)

\( \text{SRRT}^2_2 \) is the statement "every rainbow-stable 2-bounded coloring \( f : [\omega]^2 \to \omega \) has a rainbow."
Introduced by the author in [23], he proved that SRRT\$^2_2\$ is computably reducible to SEM and STS(2). This principle admits many computably equivalent formulations. We are particularly interested in a characterization which can be seen as a stable notion of DNR[\$\emptyset\$].

**Definition 6.10**
Given a function \( f : \omega \to \omega \), a function \( g \) is \( f \)-diagonalizing if \( (\forall x)[f(x) \neq g(x)] \). SDNR[\$\emptyset\$] is the statement “Every \( \Delta^0_2 \) function \( f : \omega \to \omega \) has an \( f \)-diagonalizing function”.

**Theorem 6.11** (Patey [23]) SRRT\$^2_2\$ and SDNR[\$\emptyset\$] are computably equivalent.

The following theorem extends Mileti’s result to SDNR[\$\emptyset\$]. As SDNR[\$\emptyset\$] is computably below many stable principles, we shall deduce a few more non-universality results.

**Theorem 6.12** For every \( \Delta^0_2 \) incomplete set \( X \), there exists a \( \Delta^0_2 \) function \( f : \omega \to \omega \) such that \( X \) computes no \( f \)-diagonalizing function.

**Corollary 6.13** A \( \Delta^0_2 \) degree \( d \) bounds SRRT\$^2_2\$ iff \( d = \emptyset' \).

**Proof.** As SRRT\$^2_2\$ \( \leq_c \) SRT\$^2_2\$, any computable instance of SRRT\$^2_2\$ has a \( \Delta^0_2 \) solution. So \( \emptyset' \) bounds SRRT\$^2_2\$. If \( d \) is incomplete, then by Theorem 6.12 and by SRRT\$^2_2\$ \( \leq_c \) SDNR[\$\emptyset'\$], there is a computable instance of SRRT\$^2_2\$ such that \( d \) bounds no solution.

**Corollary 6.14** No statement \( P \) such that SRRT\$^2_2\$ \( \leq_c \) \( P \leq_c \) SRRT\$^2_2\$ admits a universal instance.

**Proof.** By [13, Corollary 4.6] every \( \Delta^0_2 \) set or its complement has an incomplete \( \Delta^0_2 \) infinite subset. As \( P \leq_c \) SRRT\$^2_2\$ \( \leq_c \) D\$^2_2\$, every computable instance \( U \) of \( \emptyset \) has a \( \Delta^0_2 \) incomplete solution \( X \). By Theorem 6.12, there exists a computable coloring \( f : [\omega]^2 \to \omega \) such that \( X \) computes no infinite \( f \)-rainbow. As SRRT\$^2_2\$ \( \leq_c \) P, there exists a computable instance of \( P \) such that \( X \) does not compute a solution to it. Hence \( U \) is not a universal instance of \( P \).

**Corollary 6.15** None of SRRT\$^2_2\$, SEM, STS(2) and SFS(2) admits a universal instance.

**Proof of Theorem 6.12.** The proof is an adaptation of [20, Theorem 5.3.7]. Suppose that \( D \) is a \( \Delta^0_2 \) incomplete set. We will construct a \( \Delta^0_2 \) function \( f : \omega \to \omega \) such that \( D \) does not compute any \( f \)-diagonalizing function. We want to satisfy the following requirements for each \( e \in \omega \):

\[ R_e : \text{If } \Phi^D_e \text{ is total, then there is an } a \text{ such that } \Phi^D_e(a) = f(a). \]

For each \( e \in \omega \), define the partial function \( u_e \) by letting \( u_e(a) \) be the use of \( \Phi^D_e \) on input \( a \) if \( \Phi^D_e(a) \downarrow \) and letting \( u_e(a) \uparrow \) otherwise. We can assume w.l.o.g. that whenever \( u_e(a) \downarrow \) then \( u_e(a) \geq a \). Also define a computable partial function \( \theta \) by letting \( \theta(a) = (\mu t)[a \in \emptyset'] \) if \( a \in \emptyset' \) and \( \theta(a) \uparrow \) otherwise.

The local strategy for satisfying a single requirement \( R_e \) works as follows. If \( R_e \) receives attention at stage \( s \), this strategy does the following. First it identifies a number
a ≥ e that is not restrained by strategies of higher priority such that the following conditions holds:

(i) \( \Phi_{e,a}^0(a) \downarrow \)
(ii) \( u_{e,s}(a) < \max(0, \theta_i(a)) \)

If no such number \( a \) exists, the strategy does nothing. Otherwise it puts a restraint on \( a \) and commits to assigning \( f_s(a) = \Phi_{e,a}^0(a) \). For any such \( a \), this commitment will remain active as long as the strategy has a restraint on this element. Having done all this, the local strategy is declared to be satisfied and will not act again unless either a higher priority puts restraints on \( a \), or the value of \( u_{e,s}(a) \) or \( \theta_i(a) \) changes. In both cases the strategy gets injured and has to reset, releasing all its restraints.

To finish stage \( s \), the global strategy assigns values \( f_s(y) \) for all \( y ≤ s \) as follows: if \( y \) is committed to some value assignment of \( f_s(y) \) due to a local strategy, then define \( f_s(y) \) to be this value. If not, let \( f_s(y) = 0 \). This finishes the construction and we now turn to the verification.

For each \( e, a ∈ ω \), let \( Z_{e,a} = \{ s ∈ ω : R_e \text{ restraints } a \text{ at stage } s \} \).

Claim. For each \( e, a ∈ ω \),

(a) if \( \Phi_{e}^0(a) \uparrow \) then \( Z_{e,a} \) is finite;
(b) if \( \Phi_{e}^0(a) \downarrow \) then \( Z_{e,a} \) is either finite or cofinite.

Proof. By induction on the priority order. We consider \( Z_{e,a} \), assuming that for all \( R e \) of higher priority, the set \( Z_{e,a} \) is either finite or cofinite. First notice that \( Z_{e,a} = ς \) if \( a < e \) or \( a ∈ T_f \), so we may assume that \( a ≥ e \) and \( a ∈ \Tf \). If there exists \( e' < e \) such that \( Z_{e',a} \) is cofinite, then \( Z_{e,a} \) is finite because at most one requirement may claim \( a \) at a given stage. Suppose that \( Z_{e',a} \) is finite for all \( e' < e \). Fix \( t_0 \) such that for all \( e' < e \) and \( s ≥ t_0 \), \( R_e \) does not restrain \( a \) at stage \( s \) and \( \theta_i(a) = \theta(a) \).

Suppose that \( \Phi_{e}^0(a) \uparrow \). Fix \( t_1 ≥ t_0 \) such that \( D(c) = D_i(c) \) for all \( b ≤ \theta(a) \) and all \( s ≥ t_1 \). Then for all \( s ≥ t_1 \), if \( \Phi_{e}^0(a) \downarrow \) then we must have \( u_{e,s}(a) > \theta(a) \) because otherwise \( \Phi_{e}^0(a) \downarrow \). It follows that for all \( s ≥ t_1 \), requirement \( R_e \) does not restrain \( a \) at stage \( s \). Hence \( Z_{e,a} \) is finite.

Suppose now that \( \Phi_{e}^0(a) \downarrow \). Fix \( t_1 ≥ t_0 \) such that for all \( s ≥ t_1 \) we have \( \Phi_{e}^0(a) \downarrow \) and \( D_i(c) = D(c) \) for every \( c ≤ u_{e}(a) \). For every \( s ≥ t_1 \), \( u_{e,s}(a) = u_{e,t_1}(a) \) and \( \theta_i(a) = \theta_i(a) \) for each \( i ≤ a \). So properties (i) and (ii) will either hold at each stage \( s ≥ t_1 \), or not hold at each stage \( s ≥ t_1 \). Therefore \( Z_{e,a} \) is either finite or cofinite.

Claim. Each requirement \( R_e \) is satisfied.

Proof. Suppose that \( \Phi_{e}^0 \) is total for some \( e ∈ ω \). We will prove that \( \Phi_{e}^0 \) is not an f-diagonalizing function. Let \( A = \{ a ≥ e : (\forall e' < e)Z_{e',a} \text{ is finite} \} \). Notice that \( A \) is cofinite since for each \( e' < e \), there is at most one \( a \) such that \( Z_{e',a} \) is cofinite.

If for all but finitely many \( k ∈ ω \), we have \( k ∈ \Tf \rightarrow k ∈ \Tf'_{u_{e}(k)} \), then \( \Tf' ≤_T u_{e} ≤_T D \), contrary to hypothesis. Thus we may let \( a \) be the least element of \( \{ k ∈ A : k ∈ \Tf' \text{ or } \Tf'_{u_{e}(k)} \} \).

We then have

(1) \( a ≥ e, \Phi_{e}^0(a) \downarrow, \theta(a) > u_{e}(a) \)
(2) For all \( e' < e \), there exists \( t \) such that \( R_e \) does not claim \( a \) at any stage \( s ≥ t \).

Therefore we may fix \( t ≥ a \) such that for all \( s ≥ t \), we have \( \Phi_{e,s}^0(a) \downarrow, \theta_i(a) = \theta(a) \), \( u_{e,s}(a) = u_{e}(a) \), and for each \( e' < e \), \( R_{e'} \) does not claim \( a \) at stage \( s \). Thus, for every \( s ≥ t \), requirement \( R_e \) claims \( a' < a \) at stage \( s \). Since \( Z_{e,t} \) is either finite or cofinite
for each \( i \leq a \), it follows that \( Z_{e,a} \) is cofinite. By the above argument, we must have \( \Phi_e^D(a) \downarrow \), and by construction, \( f(a) = \Phi_e^D(a) \). Therefore \( R_e \) is satisfied.

\[ \square \]

**Claim.** The resulting function \( f \) is \( \Delta^0_2 \).

**Proof.** Consider a particular element \( a \). Because of Claim 1, if \( e > a \) then \( Z_{e,a} = \emptyset \).

We have then two cases: Either \( Z_{e,a} \) is finite for all \( e \leq a \), in which case for all but finitely many \( s \), \( f_s(a) = 0 \), or \( Z_{e,a} \) is cofinite for some \( e \). Then there is a stage \( s \) at which requirement \( R_e \) has committed \( f_s(a) = \Phi_e^D(a) \) for assignment and has never been injured. Thus \( f \) is \( \Delta^0_2 \).

\[ \square \]

**Acknowledgements.** The author is thankful to his PhD advisor Laurent Bienvenu for interesting comments and discussions. The author is funded by the John Templeton Foundation (’Structure and Randomness in the Theory of Computation’ project). The opinions expressed in this publication are those of the author(s) and do not necessarily reflect the views of the John Templeton Foundation.

**References**


---

LAboratoire PPS, Université Paris Diderot - Paris 7, Case 7014, 75205 PARIS CEDEX 13, FRANCE
E-mail address: ludovic.patey@computability.fr