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# PINT: a Static Analyzer for Transient Dynamics of Qualitative Networks with IPython Interface

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**Abstract.** The software PINT is devoted to the scalable analysis of the traces of automata networks, which encompass Boolean and discrete networks. PINT implements formal approximations of transient reachability-related properties, including mutation prediction and model reduction. PINT is distributed with command line tools, as well as a Python module *pypint*. The latter provides a seamless integration with the Jupyter IPython notebook web interface, which allows to easily save, reuse, reproduce, and share workflows of model analysis. PINT can address networks with hundreds to thousands interacting components, which are typically intractable with standard approaches.

## 1 Introduction

The computational analysis of the qualitative dynamics of biological networks faces the state space explosion problem, limiting the tractability of detailed models. Many studies have to use reduced models which often lose important properties and may lead to approximative results.

PINT provides formal and scalable analysis for the transient discrete dynamics (traces/trajectories) of automata networks, which subsume Boolean and multi-valued networks. PINT implements an abstract interpretation of traces based on a static analysis of causality of transitions. It results in over- and under-approximations of PSPACE-complete problems by  $P \cdot \exp(k-1)$  and  $NP \cdot \exp(k-1)$  problems, where  $k$  is the number of qualitative levels of network nodes (2 for Boolean networks). PINT then relies on Boolean constraint satisfaction (SAT) and Answer-Set Programming (ASP, [3]) for their efficient resolution.

Besides simple transient reachability analysis (from state  $s_0$  there exists a succession of transitions leading, even briefly, to a state satisfying a given property), PINT features include the prediction of mutations to control the reachability properties, the identification of *bifurcation* transitions responsible for differentiation processes, and model reduction which preserves transient reachability properties. For each case, returned results have formal guarantees on their correctness (under-approximations, satisfying sufficient conditions) or completeness (over-approximations, satisfying necessary conditions).

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Most of PINT analysis can typically handle networks with several hundreds of components. PINT also provides interfaces with exact model-checkers, such as NUSMV [5], ITS [16] and MOLE [25], taking advantage of implemented static model reduction to enhance their tractability on large models. Usual explicit reachable state graph analysis are also available, although other tools dedicated to Boolean or multi-valued networks already provide them, e.g., [10,14,18].

*User Interfaces* PINT can be invoked either using command line executables, suited for batch deployments, or through a programmable python interface. Moreover, its embedding in the Jupyter IPython notebook allows a user-friendly web interface to ease the management of models and calls to PINT. Jupyter notebooks provide a convenient environment for editing, saving, sharing, and reproducing model analysis workflows. It is a common framework for data-oriented bioinformatics tools [2,6,12], and has promising suitability for computational systems biology, where reproducibility is very important as well.

*Distribution* PINT is written mainly with the OCaml programming language and is actively developed since 2011. It is distributed under the free software licence CeCILL, and is available at <http://loicpauleve.name/pint> where binary packages are provided for Ubuntu Linux and Mac OS X.

The Docker<sup>1</sup> image [pauleve/pint](https://hub.docker.com/r/pauleve/pint) provides a ready-to-use PINT environment for usual operating systems (Windows, Mac OS X, Linux), and notably the Jupyter web interface. Such a kind of distribution becomes standard for providing accessible and reproducible analyses in bioinformatics, e.g., *BioContainers* [15].

## 2 Input model

PINT takes as input *asynchronous automata networks* specified in plain text. Automata networks are sets of finite-state machines having local transitions conditioned by the state of other automata in the network. The global state space of the network is the produce of the local states of individual automata, and transitions are applied non-deterministically.

Fig. 1 shows an example of automata networks with its plain text representation in PINT format. By convention, the file names end with `.an`.

Automata networks are expressive enough to encode the asynchronous semantics of Boolean and multi-valued networks. The main difference with these latter frameworks is the explicit specification of local transitions for each automaton (node) of the network, compared to a function-centred specification for Boolean and multi-valued networks [7,19].

PINT can automatically convert models expressed as Boolean or multi-valued networks using the `pint-import` command or `pypint.load()` python function. Most of the conversions are performed using GINSIM [10], enabling the support for SBML-qual, GINsim, as well as various text formats. Models can be directly imported from URLs and from CellCollective database [13]. BIOCHAM reaction networks are also supported, following their Boolean semantics [4].

<sup>1</sup> <http://docker.com>

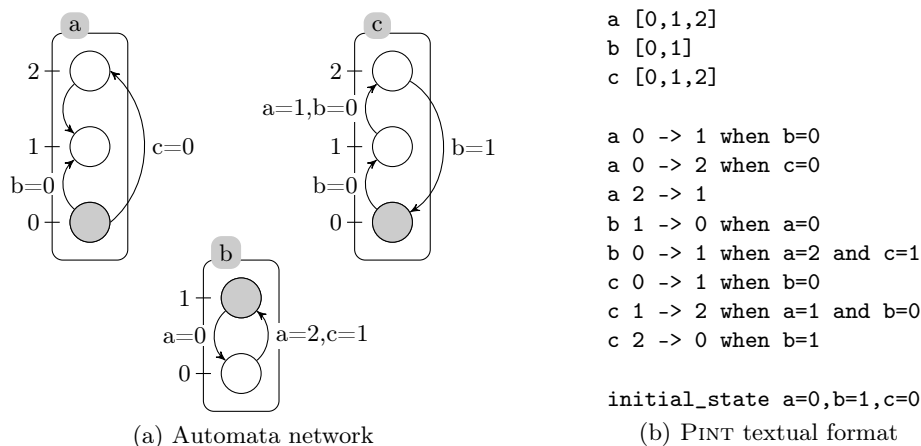


Fig. 1: (a) graphical representation of an automata network: automata are labelled boxes and their local states by circles where ticks are their identifier within the automaton. The initial state is composed of the local states in gray. A local transition is a directed edge between two local states of an automaton. Transitions can be labelled with states of other automata which are necessary to trigger the transition. (b) equivalent PINT plain text representation

### 3 Main features and benchmarks

The main originality of PINT resides in the static analysis for transient reachability properties: such an approach avoids building the reachable state transition graph, neither explicitly nor symbolically. Therefore, the analysis aims at being tractable on large networks, at the price of giving possibly incomplete results.

We present the related features, illustrated in Fig. 2, with benchmarks to support their tractability on large biological networks. Computation times have been obtained on an Intel® Core™ i7-4770 3.40GHz CPU with 16GiB RAM.

**Reachability analysis: formal approximation and model reduction** — Given an initial state, a usual problem is to determine the existence of a sequence of transitions which leads to the activation or de-activation of key components (e.g., transcription factors) or to a particular attractor. Reachability verification is a PSPACE-complete problem and its resolution often explodes on large networks. PINT implements over- and under-approximation of reachability [21,9] which allow tackling large models, although being potentially inconclusive when the over-approximation is satisfied but not the under-approximation. In such cases, one should fall back to classical model-checking. To that aim, the *goal-oriented reduction* [19] identifies transitions that do not contribute to the goal reachability, and hence can be removed prior to the reachability analysis. This model reduction preserves *all* minimal traces to the goal, and can enhance greatly the tractability of model-checking. See Table 1 and [19] for benchmarks.

Model ( nodes )	T	states	Verification of goal reachability		
			NuSMV (EF $g$ )	ITS-REACH	PINT
TCell-d (101) [1] profile 1	381	$\approx 2.4 \cdot 10^8$	2s 40Mb	0.5s 26Mb	0.02s
	<b>0</b>	<b>1</b>			
TCell-d (101) profile 2	381	KO	KO	960s 1.6Gb	4.5s
	<b>221</b>	<b>75,947,684</b>	<b>470s 270Mb</b>	<b>15s 160Mb</b>	
RBE2F (370) [22]	742	KO	KO	KO	0.2s
	<b>56</b>	<b>2,350,494</b>	<b>3s 37Mb</b>	<b>4s 13Mb</b>	
MAPK (309) [24]	1251	KO	KO	KO	48s
	<b>429</b>	<b>KO</b>	<b>KO</b>	<b>KO</b>	

Table 1: Benchmark<sup>†</sup> of goal reachability verification with two exact methods (NuSMV and ITS-REACH) and PINT, before (normal font) and after (bold font) goal-oriented model reduction; |T| is the number of local transitions in automata networks; |state| is the number of reachable global states, when computable. KO indicates an out-of-memory/time computation. In all cases PINT is conclusive.

Goal	TCell-d (101) FOXP3=1		Egf-r (104) [23] AP1=1		MAPK (309) ERK-PP=1		PID (10,229) [20] SNAIL=1	
	3-cut sets	0.06s	35	0.02s	34	0.06s	24	1.2s
4-cut sets	0.10s	101	0.02s	34	0.1s	48	5s	37
6-cut sets	0.60s	495	0.03s	34	1s	60	10m	907
3-mutations	0.30s	15	0.30s	20	5s	222	50m	7
4-mutations	0.30s	15	0.30s	22	10s	1896	50m	67
6-mutations	0.30s	15	0.30s	22	KO		50m	367

Table 2: Performance<sup>†</sup> of cut sets and mutations under-approximations with PINT depending on the maximal cardinality of returned sets.

	T	states	goal	NuSMV		PINT	
				$ t_b $	time	$ t_b $	time
EGF/TNF (28) [17]	53	3968	NFkB = 0	5	0.2s	2	0.1s
MAPK (53) [11]	173	KO	Proliferation = 1	KO		13	40s
TCell-d (101)	381	KO	FOXP3 = 1	KO		4	58s

Table 3: Performance<sup>†</sup> of exact and approximated *identification of bifurcation transitions* with NuSMV and PINT, respectively;  $|t_b|$  is the number of identified bifurcation transitions.

<sup>†</sup> Scripts and models available at <http://loicpauleve.name/pint-benchmarks.tbz2>

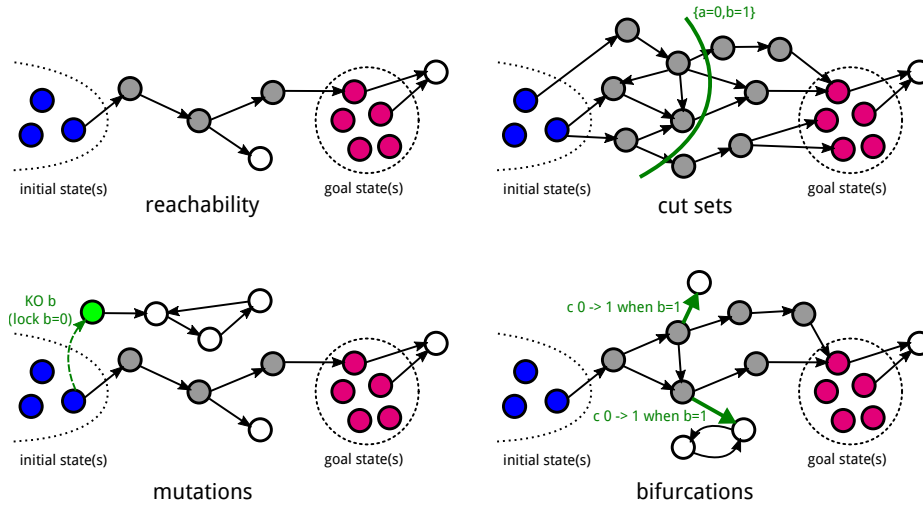


Fig. 2: Illustration of main features of PINT related to the transient reachability of a set of goal states from (a set of) initial state(s). Circles represent global states of the network and plain arrows dynamical transitions. Gray (resp. white) states are states which are (resp. are not) connected to a goal state.

**Prediction of mutations for controlling reachability** — Given an initial state and a goal state of interest, PINT provides several methods to control the transient reachability of the goal.

The most scalable approach identifies cut sets of all the paths of transitions leading to the goal. A cut set consists in one or several local states of automata which are necessary for the goal reachability: if one prevents the transitions involving these local states, the goal is disconnected from the initial state. PINT provides extremely scalable under-approximation of cut sets [20], which is tractable on Boolean networks with thousands of nodes (Table 2). Cut sets can thus be implemented as mutations which lock automata to its initial local state.

An alternative approach relies on a combination of static analysis and SAT solving and allows to directly infer mutations (gain or loss of function) which prevent the goal reachability. Whereas less scalable than cut set computations, it provides in general complementary solutions to cut sets, notably by identifying mutations which modify the initial state of the network.

**Identification of bifurcation transitions** — PINT implements static analysis for identifying so-called *bifurcation transitions* [8] after which the systems loses the capability to reach a given goal. Bifurcation transitions correspond to local transitions of the automata network which turn out to be important decision steps during differentiation processes. They can be fully identified by model-checking, but the static analysis in PINT allows tackling larger models, at the price of returning incomplete results (Table 3).

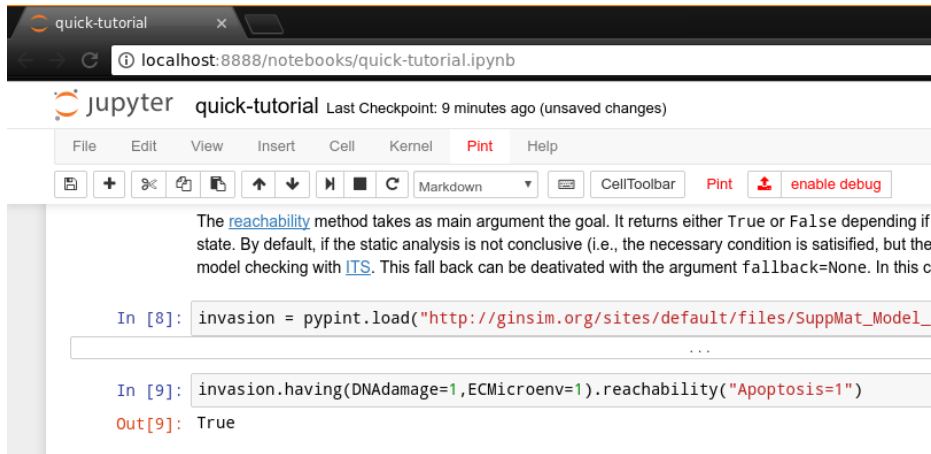


Fig. 3: Screen capture of Jupyter web interface running *pypint* in a notebook.

## 4 Integration with Jupyter IPython Web Notebook

Jupyter (<http://jupyter.org>) provides an interactive web interface for creating documents, named *notebooks*, which contain code, equations, and formatted texts. A notebook typically describes a full workflow of analysis, both with textual explanations and the full code and parameters to reproduce the results. It is a very popular framework in data science, including in bioinformatics [6,12]. A notebook is a single file which can be easily modified, shared, re-executed, and visualized online. For instance, the companion quick tutorial is available at <http://nbviewer.jupyter.org/github/pauleve/pint/blob/master/notebook/quick-tutorial.ipynb>.

The *pypint* module provides custom integration within the Jupyter IPython notebook, with custom menus and actions for loading models and executing PINT commands, as well as direct visualization of data structures. See Fig. 3 and the companion quick tutorial for a preview.

## 5 Conclusion

In this paper, we presented the prominent features of PINT on the static analysis for transient reachability of automata networks, from property verification to inference, which are tractable on large biological networks. PINT also implements classical state transition graph analysis, from fixpoint computation (using SAT solving) to explicit state space exploration, with a limited scalability. A tour of features is given at <https://loicpauleve.name/pint/doc/#Tutorial>.

In the next major release, we plan to add full support for synchronized local transitions, i.e., transitions that modify simultaneously the state of several automata. This improvement will allow to import any safe (1-bounded) Petri nets, broadening the class of supported dynamical models.

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