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Stochastic Optimization for Large-scale Optimal Transport

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Abstract

Optimal transport (OT) defines a powerful framework to compare probability distributions in a geometrically faithful way. However, the practical impact of OT is still limited because of its computational burden. We propose a new class of stochastic optimization algorithms to cope with large-scale OT problems. These methods can handle arbitrary distributions (either discrete or continuous) as long as one is able to draw samples from them, which is the typical setup in high-dimensional learning problems. This alleviates the need to discretize these densities, while giving access to provably convergent methods that output the correct distance without discretization error. These algorithms rely on two main ideas: (a) the dual OT problem can be re-cast as the maximization of an expectation; (b) the entropic regularization of the primal OT problem yields a smooth dual optimization which can be addressed with algorithms that have a provably faster convergence. We instantiate these ideas in three different setups: (i) when comparing a discrete distribution to another, we show that incremental stochastic optimization schemes can beat Sinkhorn’s algorithm, the current state-of-the-art finite dimensional OT solver; (ii) when comparing a discrete distribution to a continuous density, a semi-discrete reformulation of the dual program is amenable to averaged stochastic gradient descent, leading to better performance than approximately solving the problem by discretization; (iii) when dealing with two continuous densities, we propose a stochastic gradient descent over a reproducing kernel Hilbert space (RKHS). This is currently the only known method to solve this problem, apart from computing OT on finite samples. We backup these claims on a set of discrete, semi-discrete and continuous benchmark problems.

1 Introduction

Many problems in computational sciences require to compare probability measures or histograms. As a set of representative examples, let us quote: bag-of-visual-words comparison in computer vision [18], color and shape processing in computer graphics [22], bag-of-words for natural language processing [12] and multi-label classification [9]. In all of these problems, a geometry between the features (words, visual words, labels) is usually known, and can be leveraged to compare probability distributions in a geometrically faithful way. This underlying geometry might be for instance the planar Euclidean domain for 2-D shapes, a perceptual 3D color metric space for image processing or a high-dimensional semantic embedding for words. Optimal transport (OT) [24] is the canonical

way to automatically lift this geometry to define a metric for probability distributions. That metric is
known as the Wasserstein or earth mover’s distance. As an illustrative example, OT can use a metric
between words to build a metric between documents that are represented as frequency histograms of
words (see [12] for details). All the above-cited lines of work advocate, among others, that OT is the
natural choice to solve these problems, and that it leads to performance improvement when compared
to geometrically-oblivious distances such as the Euclidean or $\chi^2$ distances or the Kullback-Leibler
divergence. However, these advantages come at the price of an enormous computational overhead.
This is especially true because current OT solvers require to sample beforehand these distributions
on a pre-defined set of points, or on a grid. This is both inefficient (in term of storage and speed)
and counter-intuitive. Indeed, most high-dimensional computational scenarios naturally represent
distributions as objects from which one can sample, not as density functions to be discretized.
Our goal is to alleviate these shortcomings. We propose a class of provably convergent stochastic
optimization schemes that can handle both discrete and continuous distributions through sampling.

Previous works. The prevalent way to compute OT distances is by solving the so-called Kantorovitch
problem [11] (see Section 2 for a short primer on the basics of OT formulations), which boils down
to a large-scale linear program when dealing with discrete distributions (i.e., finite weighted sums of
Dirac masses). This linear program can be solved using network flow solvers, which can be further
refined to assignment problems when comparing measures of the same size with uniform weights [3].
Recently, regularized approaches that solve the OT with an entropic penalization [6] have been shown
to be extremely efficient to approximate OT solutions at a very low computational cost. These regular-
ized approaches have supported recent applications of OT to computer graphics [22] and machine
learning [9]. These methods apply the celebrated Sinkhorn algorithm [21], and can be extended
to solve more exotic transportation-related problems such as the computation of barycenters [22].
Their chief computational advantage over competing solvers is that each iteration boils down to
matrix-vector multiplications, which can be easily parallelized, streams extremely well on GPU, and
enjoys linear-time implementation on regular grids or triangulated domains [22].

These methods are however purely discrete and cannot cope with continuous densities. The only
known class of methods that can overcome this limitation are so-called semi-discrete solvers [1],
that can be implemented efficiently using computational geometry primitives [13]. They can compute
distance between a discrete distribution and a continuous density. Nonetheless, they are restricted to
the Euclidean squared cost, and can only be implemented in low dimensions (2-D and 3-D). Solving
these semi-discrete problems efficiently could have a significant impact for applications to density
fitting with an OT loss [4] for machine learning applications, see [14]. Lastly, let us point out that
there is currently no method that can compute OT distances between two continuous densities, which
is thus an open problem we tackle in this article.

Contributions. This paper introduces stochastic optimization methods to compute large-scale optimal
transport in all three possible settings: discrete OT, to compare a discrete vs. another discrete measure;
semi-discrete OT, to compare a discrete vs. a continuous measure; and continuous OT, to compare
a continuous vs. another continuous measure. These methods can be used to solve classical OT
problems, but they enjoy faster convergence properties when considering their entropic-regularized
versions. We show that the discrete regularized OT problem can be tackled using incremental
algorithms, and we consider in particular the stochastic averaged gradient (SAG) method [20]. Each
iteration of that algorithm requires $N$ operations ($N$ being the size of the supports of the input
distributions), which makes it scale better in large-scale problems than the state-of-the-art Sinkhorn
algorithm, while still enjoying a convergence rate of $O(1/k)$, $k$ being the number of iterations. We
show that the semi-discrete OT problem can be solved using averaged stochastic gradient descent
(SGD), whose convergence rate is $O(1/\sqrt{k})$. This approach is numerically advantageous over the
brute force approach consisting in sampling first the continuous density to solve next a discrete OT
problem. Lastly, for continuous optimal transport, we propose a novel method which makes use of an
expansion of the dual variables in a reproducing kernel Hilbert space (RKHS). This allows us for
the first time to compute with a converging algorithm OT distances between two arbitrary densities,
under the assumption that the two potentials belong to such an RKHS.

Notations. In the following we consider two metric spaces $\mathcal{X}$ and $\mathcal{Y}$. We denote by $\mathcal{M}_+^1(\mathcal{X})$ the set
of positive Radon probability measures on $\mathcal{X}$, and $\mathcal{C}(\mathcal{X})$ the space of continuous functions on $\mathcal{X}$. Let
$\mu \in \mathcal{M}_+^1(\mathcal{X})$, $\nu \in \mathcal{M}_+^1(\mathcal{Y})$, we define

$$
\Pi(\mu, \nu) \overset{\text{def}}{=} \left\{ \pi \in \mathcal{M}_+^1(\mathcal{X} \times \mathcal{Y}) \mid \forall(A, B) \subset \mathcal{X} \times \mathcal{Y}, \pi(A \times \mathcal{Y}) = \mu(A), \pi(\mathcal{X} \times B) = \nu(B) \right\},
$$
the set of joint probability measures on $\mathcal{X} \times \mathcal{Y}$ with marginals $\mu$ and $\nu$. The Kullback-Leibler divergence between joint probabilities is defined as

$$\forall (\pi, \xi) \in M^1_+(\mathcal{X} \times \mathcal{Y})^2, \quad \text{KL}(\pi | \xi) \overset{\text{def}}{=} \int_{\mathcal{X} \times \mathcal{Y}} \left( \log \left( \frac{d\pi}{d\xi}(x,y) \right) - 1 \right) d\xi(x,y),$$

where we denote $\frac{d\pi}{d\xi}$ the relative density of $\pi$ with respect to $\xi$, and by convention $\text{KL}(\pi | \xi) \overset{\text{def}}{=} +\infty$ if $\pi$ does not have a density with respect to $\xi$. The Dirac measure at point $x$ is $\delta_x$. For a set $C$, $\iota_C(x) = 0$ if $x \in C$ and $\iota_C(x) = +\infty$ otherwise. The probability simplex of $N$ bins is $\Sigma_N = \{ \mu \in \mathbb{R}^N_+ : \sum_i \mu_i = 1 \}$. Element-wise multiplication of vectors is denoted by $\odot$ and $K^\top$ denotes the transpose of a matrix $K$. We define $1_N = (1, \ldots, 1)^\top \in \mathbb{R}^N$ and $0_N = (0, \ldots, 0)^\top \in \mathbb{R}^N$.

## 2 Optimal Transport: Primal, Dual and Semi-dual Formulations

We consider the optimal transport problem between two measures $\mu \in M^1_+(\mathcal{X})$ and $\nu \in M^1_+(\mathcal{Y})$, defined on metric spaces $\mathcal{X}$ and $\mathcal{Y}$. No particular assumption is made on the form of $\mu$ and $\nu$, we only assume that they both can be sampled from to be able to apply our algorithms.

### Primal, Dual and Semi-dual Formulations


$$\forall (\mu, \nu) \in M^1_+(\mathcal{X}) \times M^1_+(\mathcal{Y}), \quad W_\varepsilon(\mu, \nu) \overset{\text{def}}{=} \min_{\pi \in \Pi(\mu, \nu)} \int_{\mathcal{X} \times \mathcal{Y}} c(x,y)d\pi(x,y) + \varepsilon \text{KL}(\pi | \mu \otimes \nu). \quad (P_\varepsilon)$$

Here $c \in C(\mathcal{X} \times \mathcal{Y})$ and $c(x,y)$ should be interpreted as the “ground cost” to move a unit of mass from $x$ to $y$. This $c$ is typically application-dependent, and reflects some prior knowledge on the data to process. We refer to the introduction for a list of previous work where various examples (in imaging, vision, graphics or machine learning) of such costs are given.

When $\mathcal{X} = \mathcal{Y}$, $\varepsilon = 0$ and $c = d^p$ for $p \geq 1$, where $d$ is a distance on $\mathcal{X}$, then $W_0(\mu, \nu)^\frac{1}{p}$ is known as the $p$-Wasserstein distance on $M^1_+(\mathcal{X})$. Note that this definition can be used for any type of measure, both discrete and continuous. When $\varepsilon > 0$, problem $(P_\varepsilon)$ is strongly convex, so that the optimal $\pi$ is unique, and algebraic properties of the KL regularization result in computations that can be tackled using the Sinkhorn algorithm [6].

For any $c \in C(\mathcal{X} \times \mathcal{Y})$, we define the following constraint set

$$U_c \overset{\text{def}}{=} \{ (u, v) \in C(\mathcal{X}) \times C(\mathcal{Y}) : \forall (x,y) \in \mathcal{X} \times \mathcal{Y}, u(x) + v(y) \leq c(x,y) \},$$

and define its indicator function as well as its “smoothed” approximation

$$\iota_{U_c}(u, v) \overset{\text{def}}{=} \left\{ \begin{array}{ll} u_c(u, v) & \text{if } \varepsilon = 0, \\
\varepsilon \int_{\mathcal{X} \times \mathcal{Y}} \exp \left( \frac{u(x) + v(y) - c(x,y)}{\varepsilon} \right) d\mu(x)d\nu(y) & \text{if } \varepsilon > 0. \end{array} \right. \quad (1)$$

For any $v \in C(\mathcal{Y})$, we define its $c$-transform and its “smoothed” approximation

$$\forall x \in \mathcal{X}, \quad v^c(x) \overset{\text{def}}{=} \left\{ \begin{array}{ll} \min_{y \in \mathcal{Y}} c(x,y) - v(y) & \text{if } \varepsilon = 0, \\
-\varepsilon \log \left( \int_{\mathcal{Y}} \exp \left( \frac{v(y) - c(x,y)}{\varepsilon} \right) d\nu(y) \right) & \text{if } \varepsilon > 0. \end{array} \right. \quad (2)$$

The proposition below describes two dual problems. It is central to our analysis and paves the way for the application of stochastic optimization methods.

### Proposition 2.1 (Dual and semi-dual formulations)

For $\varepsilon \geq 0$, one has

$$W_\varepsilon(\mu, \nu) = \max_{u \in C(\mathcal{X}), v \in C(\mathcal{Y})} F_\varepsilon(u, v) \overset{\text{def}}{=} \int_{\mathcal{X}} u(x)d\mu(x) + \int_{\mathcal{Y}} v(y)d\nu(y) - \iota_{U_c}(u, v), \quad (D_\varepsilon)$$

$$= \max_{v \in C(\mathcal{Y})} H_\varepsilon(v) \overset{\text{def}}{=} \int_{\mathcal{X}} v^{c^\varepsilon}(x)d\mu(x) + \int_{\mathcal{Y}} v(y)d\nu(y) - \varepsilon, \quad (S_\varepsilon)$$

where $\iota_{U_c}$ is defined in [1] and $v^{c^\varepsilon}$ in [2]. Furthermore, $u$ solving $(D_\varepsilon)$ is recovered from an optimal $v$ solving $(S_\varepsilon)$ as $u = v^{c^\varepsilon}$. For $\varepsilon > 0$, the solution $\pi$ of $(P_\varepsilon)$ is recovered from any $(u, v)$ solving $(D_\varepsilon)$ as $d\pi(x,y) = \exp\left( \frac{u(x) + v(y) - c(x,y)}{\varepsilon} \right) d\mu(x)d\nu(y)$. 


Proof. Problem \( (D_2) \) is the convex dual of \( (P_2) \), and is derived using Fenchel-Rockafellar’s theorem. The relation between \( u \) and \( v \) is obtained by writing the first order optimality condition for \( v \) in \( (D_2) \). Plugging this expression back in \( (D_2) \) yields \( (S_2) \).

Problem \( (P_2) \) is called the primal while \( (D_2) \) is its associated dual problem. We refer to \( (S_2) \) as the “semi-dual” problem, because in the special case \( \varepsilon = 0 \), \( (S_2) \) boils down to the so-called semi-discrete OT problem \( \Pi \). Both dual problems are concave maximization problems. The optimal dual variables \((u, v)\)—known as Kantorovitch potentials—are not unique, since for any solution \((u, v)\) of \( (D_2) \), \((u + \lambda, v - \lambda)\) is also a solution for any \( \lambda \in \mathbb{R} \). When \( \varepsilon > 0 \), they can be shown to be unique up to this scalar translation \( \delta \).

A key advantage of \( (P_2) \) is that \((u, v)\), which does not ensure strong convexity of the semi-dual problem, as it impacts the convergence properties of the stochastic algorithms (stochastic averaged gradient and stochastic gradient descent) described below.

\[ \forall \varepsilon > 0, \quad F_\varepsilon(u, v) = \mathbb{E}_{X,Y} [f_\varepsilon(X,Y,u,v)] \quad \text{and} \quad \forall \varepsilon \geq 0, \quad H_\varepsilon(v) = \mathbb{E}_X [h_\varepsilon(X,v)], \]  

where the random variables \( X \) and \( Y \) are independent and distributed according to \( \mu \) and \( \nu \) respectively, and where, for \((x, y) \in \mathcal{X} \times \mathcal{Y} \) and \((u, v) \in \mathcal{C}(\mathcal{X}) \times \mathcal{C}(\mathcal{Y}) \),

\[ \forall \varepsilon > 0, \quad f_\varepsilon(x, y, u, v) \overset{\text{def.}}{=} u(x) + v(y) - \varepsilon \exp \left( \frac{u(x) + v(y) - c(x, y)}{\varepsilon} \right), \]

\[ \forall \varepsilon \geq 0, \quad h_\varepsilon(x, v) \overset{\text{def.}}{=} \int_Y v(y) \text{d}v(y) + \frac{c(x, y)}{\varepsilon} - \varepsilon. \]

This reformulation is at the heart of the methods detailed in the remainder of this article. Note that the dual problem \( (D_2) \) cannot be cast as an unconstrained expectation maximization problem when \( \varepsilon = 0 \), because of the constraint on the potentials which arises in that case.

When \( \nu \) is discrete, i.e. \( \nu = \sum_{j=1}^J \nu_j \delta_{y_j} \), the potential \( v \) is a \( J \)-dimensional vector \( (v_j)_{j=1}^J \) and we can compute the gradient of \( h_\varepsilon \). When \( \varepsilon > 0 \) the gradient reads \( \nabla_v h_\varepsilon(v, x) = \nu - \pi(x) \) and the hessian is given by \( \partial^2_v h_\varepsilon(v, x) = \frac{1}{\varepsilon} (\pi(x) \pi(x)^T - \text{diag}(\pi(x))) \) where \( \pi(x) = \text{exp}(\frac{c(x,y_j)}{\varepsilon} (\sum_{j=1}^J \text{exp}(\frac{c(x,y_j)}{\varepsilon})))^{-1} \). The eigenvalues of the hessian can be upper bounded by \( \frac{1}{\varepsilon} \), which guarantees a lipschitz gradient, and lower bounded by 0 which does not ensure strong convexity. In the unregularized case, \( h_0 \) is not smooth and a subgradient is given by \( \nabla_v h_0(v, x) = \nu - \tilde{\pi}(x) \), where \( \tilde{\pi}(x)_i = \chi_{i=j} \), and \( j^* = \arg \min_j c(x, y_j) - v_j \) (when several elements are in the argmin, we arbitrarily choose one of them to be \( j^* \)). We insist on the lack of strong convexity of the semi-dual problem, as it impacts the convergence properties of the stochastic algorithms (stochastic averaged gradient and stochastic gradient descent) described below.

3 Discrete Optimal Transport

We assume in this section that both \( \mu \) and \( \nu \) are discrete measures, i.e. finite sums of Diracs, of the form \( \mu = \sum_{i=1}^I \mu_i \delta_{x_i} \) and \( \nu = \sum_{j=1}^J \nu_j \delta_{y_j} \), where \((x_i)_i \subset \mathcal{X} \) and \((y_j)_j \subset \mathcal{Y} \), and the histogram vector weights are \( \mu \in \Sigma_I \) and \( \nu \in \Sigma_J \). These discrete measures may come from the evaluation of continuous densities on a grid, counting features in a structured object, or be empirical measures based on samples. This setting is relevant for several applications, including all known applications of the earth mover’s distance. We show in this section that our stochastic formulation can prove extremely efficient to compare measures with a large number of points.

Discrete Optimization and Sinkhorn. In this setup, the primal \( (P_2) \), dual \( (D_2) \) and semi-dual \( (S_2) \) problems can be rewritten as finite-dimensional optimization problems involving the cost matrix...
The state-of-the-art method to solve the discrete regularized OT (i.e. when $\varepsilon > 0$) is Sinkhorn’s algorithm [6, Alg.1], which has linear convergence rate [8]. It corresponds to a block coordinate maximization, successively optimizing $(\bar{D}_\varepsilon)$ with respect to either $u$ or $v$. Each iteration of this algorithm is however costly, because it requires a matrix-vector multiplication. Indeed, this corresponds to a “batch” method where all the samples $(x_i)$ and $(y_j)$ are used at each iteration, which has thus complexity $O(N^2)$ where $N = \max(I,J)$. We now detail how to alleviate this issue using online stochastic optimization methods.

**Incremental Discrete Optimization when $\varepsilon > 0$.** Stochastic gradient descent (SGD), in which an index $k$ is drawn from distribution $\mu$ at each iteration can be used to minimize the finite sum that appears in $(\bar{D}_\varepsilon)$. The gradient of that term $\bar{h}_\varepsilon(x_k, \cdot)$ can be used as a proxy for the full gradient in a standard gradient ascent step to maximize $\bar{H}_\varepsilon$.

When $\varepsilon > 0$, the finite sum appearing in $(\bar{S}_\varepsilon)$ suggests to use incremental gradient methods—rather than purely stochastic ones—which are known to converge faster than SGD. We propose to use the stochastic averaged gradient (SAG) [20]. As SGD, SAG operates at each iteration by sampling a point $x_k$ from $\mu$, to compute the gradient corresponding to that sample for the current estimate $v$. Unlike SGD, SAG keeps in memory a copy of that gradient. Another difference is that SAG applies a fixed length update, in the direction of the average of all gradients stored so far, which provides a better proxy of the gradient corresponding to the entire sum. This improves the convergence rate to $|\bar{H}_\varepsilon(v^*_\varepsilon) - \bar{H}_\varepsilon(v_k)| = O(1/k)$, where $v^*_\varepsilon$ is a minimizer of $\bar{H}_\varepsilon$, at the expense of storing the gradient for each of the $I$ points. This expense can be mitigated by considering mini-batches instead of individual points. Note that the SAG algorithm is adaptive to strong-convexity and will be linearly convergent around the optimum. The pseudo-code for SAG is provided in Algorithm 1, and we defer more details on SGD for Section 4 in which it will be shown to play a crucial role. Note that the Lipschitz constant of all these terms is upperbounded by $L = \max_i \mu_i / \varepsilon$.

**Numerical Illustrations on Bags of Word-Embeddings.** Comparing texts using a Wasserstein distance on their representations as clouds of word embeddings has been recently shown to yield state-of-the-art accuracy for text classification [12]. The authors of [12] have however highlighted that this accuracy comes at a large computational cost. We test our stochastic approach to discrete OT in this scenario, using the complete works of 35 authors (names in supplementary material). We use Glove word embeddings [15] to represent words, namely $X = Y = \mathbb{R}^{300}$. We discard all most frequent 1,000 words that appear at the top of the file glove.840B.300d provided on the authors’ website. We sample $N = 20,000$ words (found within the remaining huge dictionary of relatively rare words) from each authors’ complete work. Each author is thus represented as a cloud of 20,000 points in $\mathbb{R}^{300}$. The cost function $c$ between the word embeddings is the squared-Euclidean distance, re-scaled so that it has a unit empirical median on 2,000 points sampled randomly among all vector embeddings. We set $\varepsilon$ to 0.01 (other values are considered in the supplementary material). We compute all $(35 \times 34 / 2 = 595)$ pairwise regularized Wasserstein distances using both the Sinkhorn algorithm and SAG. Following the recommendations in [20], SAG’s stepsize is tested for $3$ different settings, $1/L, 3/L$ and $5/L$. The convergence of each algorithm is measured by computing the $\ell_1$ norm of the gradient of the full sum (which also corresponds to the marginal violation of the primal transport solution that can be recovered with these dual variables [9]), as well as the $\ell_2$ norm of the.

---

**Algorithm 1 SAG for Discrete OT**

**Input:** $C$

**Output:** $v$

$v \leftarrow 0, d \leftarrow 0, \forall i, \mathbf{g}_i \leftarrow 0, \forall i$

for $k = 1, 2, \ldots$ do

Sample $i \in \{1, 2, \ldots, I\}$ uniform.

$d \leftarrow d - \mathbf{g}_i$

$\mathbf{g}_i \leftarrow \mu_i \nabla v \bar{h}_\varepsilon(x_i, v)$

$d \leftarrow d + \mathbf{g}_i \quad v \leftarrow v + Cd$

end for

---

$c \in \mathbb{R}^{I \times J}_+$ defined by $c_{i,j} = c(x_i, y_j)$:

$$W_\varepsilon(\mu, \nu) = \min_{\pi \in \mathbb{R}^{I \times J}_+} \left\{ \sum_{i,j} c_{i,j} \pi_{i,j} + \varepsilon \sum_{i,j} \left( \log \frac{\pi_{i,j}}{\mu_i \nu_j} - 1 \right) \pi_{i,j} ; \pi 1 = \mu, \pi^T 1 = \nu \right\};$$

$$= \max_{u \in \mathbb{R}, v \in \mathbb{R}} \frac{1}{\varepsilon} \sum_{i,j} u_i \mu_i + v_j \nu_j - \varepsilon \sum_{i,j} \exp \left( \frac{u_i + v_j - c_{i,j}}{\varepsilon} \right) \mu_i \nu_j; \text{ (for } \varepsilon > 0 \) \quad (\bar{D}_\varepsilon)$$

$$= \max_{v \in \mathbb{R}^J} \bar{H}_\varepsilon(v) = \sum_{i \in I} \bar{h}_\varepsilon(x_i, v) \mu_i, \quad \text{where} \quad (\bar{S}_\varepsilon)$$

$$\bar{h}_\varepsilon(x, v) = \sum_{j \in J} v_j \nu_j + \left\{ -\varepsilon \log \left( \sum_{j \in J} \exp \left( \frac{v_j - c(x, y_j)}{\varepsilon} \right) \nu_j \right) - \varepsilon \right\} \quad \text{if } \varepsilon > 0,$$

$$\min_{j \in J} \left( c(x, y_j) - v_j \right) \quad \text{if } \varepsilon = 0,$$
deviation to the optimal scaling found after 4,000 passes for any of the three methods. Results are presented in Fig. 1 and suggest that SAG can be more than twice faster than Sinkhorn on average for all tolerance thresholds. Note that SAG retains exactly the same parallel properties as Sinkhorn: all of these computations can be streamlined on GPUs. We used 4 Tesla K80 cards to compute both SAG and Sinkhorn results. For each computation, all of these computations can be streamlined on GPUs. We used 4 Tesla K80 cards to compute both SAG and Sinkhorn results. For each computation, all of these computations can be streamlined on GPUs. We used 4 Tesla K80 cards to compute both SAG and Sinkhorn results. For each computation, all of these computations can be streamlined on GPUs. We used 4 Tesla K80 cards to compute both SAG and Sinkhorn results. For each computation, all of these computations can be streamlined on GPUs. We used 4 Tesla K80 cards to compute both SAG and Sinkhorn results. 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Figure 2 shows the evolution of $\|v_k - v_0^*\|_2 / \|v_0^*\|_2$ as a function of $k$. It highlights the influence of the regularization parameter $\varepsilon$ on the iterates of SGD. While the regularization iterates converge faster, they do not converge to the correct unregularized solution. This figure also illustrates the convergence theorem of solution of $(S_\varepsilon)$ toward those $(S_0)$ when $\varepsilon \to 0$, which can be found in the supplementary material.

Figure 2(b) shows the evolution of $\|v_k - v_0^*\|_2 / \|v_0^*\|_2$ as a function of $k$, for a fixed regularization parameter value $\varepsilon = 10^{-2}$. It compares SGD to SAG using different numbers $N$ of samples for the empirical measures $\hat{\mu}_N$. While SGD converges to the true solution of the semi-discrete problem, the solution computed by SAG is biased because of the approximation error which comes from the discretization of $\mu$. This error decreases when the sample size $N$ is increased, as the approximation of $\mu$ by $\hat{\mu}_N$ becomes more accurate.

## 5 Continuous optimal transport using RKHS

In the case where neither $\mu$ nor $\nu$ are discrete, problem $(\mathcal{D}_\varepsilon)$ is infinite-dimensional, so it cannot be solved directly using SGD. We propose in this section to solve the initial dual problem $(\mathcal{D}_\varepsilon)$, using expansions of the dual variables in two reproducing kernel Hilbert spaces (RKHS). Recall that contrarily to the semi-discrete setting, we can only solve the regularized problem here (i.e. $\varepsilon > 0$), since $(\mathcal{D}_0)$ cannot be cast as an expectation maximization problem when $\varepsilon = 0$. In the particular case where $c(x, y) = |x - y|$, note that the unregularized problem is equivalent to the maximum mean discrepancy problem tackled in [10].

### Stochastic Continuous Optimization

We consider two RKHS $\mathcal{H}$ and $\mathcal{G}$ defined on $\mathcal{X}$ and on $\mathcal{Y}$, with kernels $\kappa$ and $\ell$, associated with norms $\|\cdot\|_\mathcal{H}$ and $\|\cdot\|_\mathcal{G}$. Recall the two main properties of RKHS: (a) if $u \in \mathcal{H}$, then $u(x) = \langle u, \kappa(\cdot, x) \rangle_{\mathcal{H}}$ and (b) $\kappa(x, x') = \langle \kappa(\cdot, x), \kappa(\cdot, x') \rangle_{\mathcal{H}}$.

The dual problem $(\mathcal{D}_\varepsilon)$ is conveniently re-written in $(\mathcal{D})$ as the maximization of the expectation of $f^*(X, Y, u, v)$ with respect to the random variables $(X, Y) \sim \mu \otimes \nu$. The SGD algorithm applied to this problem reads, starting with $u_0 = 0$ and $v_0 = 0$,

$$
(u_k, v_k) \overset{\text{def.}}{=} (u_{k-1}, v_{k-1}) + \frac{C}{\sqrt{k}} \nabla f(x_k, y_k, u_{k-1}, v_{k-1}) \in \mathcal{H} \times \mathcal{G},
$$

where $(x_k, y_k)$ are i.i.d. samples from $\mu \otimes \nu$. The following proposition shows that these $(u_k, v_k)$ iterates can be expressed as finite sums of kernel functions, with a simple recursion formula.

### Proposition 5.1

The iterates $(u_k, v_k)$ defined in $(\mathcal{D})$ satisfy

$$
(u_k, v_k) \overset{\text{def.}}{=} \sum_{i=1}^{k} \alpha_i(\kappa(\cdot, x_i), \ell(\cdot, y_i)), \quad \text{where} \quad \alpha_i \overset{\text{def.}}{=} \Pi_B\left(\frac{C}{\sqrt{i}} \left(1 - e^{-\alpha_{i-1}(x_i) + \nu_{i-1}(y_i) - \epsilon(x_i, y_i)}\right)\right),
$$

where $(x_i, y_i)_{i=1}^k$ are i.i.d samples from $\mu \otimes \nu$ and $\Pi_B$ is the projection on the centered ball of radius $r$. If the solutions of $(\mathcal{D}_\varepsilon)$ are in the $\mathcal{H} \times \mathcal{G}$ and if $r$ is large enough, the iterates $(u_k, v_k)$ converge to a solution of $(\mathcal{D}_\varepsilon)$.

The proof of proposition 5.1 can be found in the supplementary material.
We have shown in this work that the computations behind (regularized) optimal transport can be considerably alleviated, or simply enabled, using a stochastic optimization approach. In the discrete case, we have shown that incremental gradient methods can surpass the Sinkhorn algorithm in terms of efficiency, taken for granted that the (constant) stepsize has been correctly selected, which should be possible in practical applications. We have also proposed the first known methods that can address the challenging semi-discrete and continuous cases. All of these three settings can open new perspectives for the application of OT to high-dimensional problems.

Conclusion

We have shown in this work that the computations behind (regularized) optimal transport can be considerably alleviated, or simply enabled, using a stochastic optimization approach. In the discrete case, we have shown that incremental gradient methods can surpass the Sinkhorn algorithm in terms of efficiency, taken for granted that the (constant) stepsize has been correctly selected, which should be possible in practical applications. We have also proposed the first known methods that can address the challenging semi-discrete and continuous cases. All of these three settings can open new perspectives for the application of OT to high-dimensional problems.
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References


A Convergence of \((\mathcal{S}_\varepsilon)\) as \(\varepsilon \to 0\)

The convergence of the solution of \((\mathcal{P}_\varepsilon)\) toward a solution of \((\mathcal{P}_0)\) as \(\varepsilon \to 0\) is proved in [4]. The convergence of solutions of \((\mathcal{P}_\varepsilon)\) toward solutions of \((\mathcal{D}_0)\) as \(\varepsilon \to 0\) is proved for the special case of discrete measures in [5]. To the best of our knowledge, the behavior of \((\mathcal{S}_\varepsilon)\) has not been studied in the literature, and we propose a convergence result in the case where \(\nu\) is discrete, which is the setting in which this formulation is most advantageous.

**Proposition A.1.** We assume that \(\forall y \in \mathcal{Y}, c(\cdot, y) \in L^1(\mu)\), that \(\nu = \sum_{j=1}^J \nu_j \delta_{y_j}\), and we fix \(x_0 \in \mathcal{X}\). For all \(\varepsilon > 0\), let \(v^\varepsilon\) be the unique solution of \((\mathcal{S}_\varepsilon)\) such that \(v^\varepsilon(x_0) = 0\). Then \((v^\varepsilon)_\varepsilon\) is bounded and all its converging sub-sequences for \(\varepsilon \to 0\) are solutions of \((\mathcal{S}_0)\).

We first prove a useful lemma.

**Lemma A.2.** If \(\forall y, x \mapsto c(x, y) \in L^1(\mu)\) then \(H_\varepsilon\) converges pointwise to \(H_0\).

**Proof.** Let \(\alpha_j(x) \overset{\text{def}}{=} v_j - c(x, y_j)\) and \(j^* \overset{\text{def}}{=} \arg\max_{j} \alpha_j(x)\).

On the one hand, since \(\forall j, \alpha_j(x) \leq \alpha_{j^*}(x)\) we get
\[
\varepsilon \log\left(\sum_{j=1}^J e^{\frac{\alpha_j(x)}{\varepsilon}} \nu_j\right) = \varepsilon \log\left(e^{\frac{\alpha_{j^*}(x)}{\varepsilon}} \sum_{j=1}^J e^{\frac{\alpha_j(x) - \alpha_{j^*}(x)}{\varepsilon}} \nu_j\right) \leq \alpha_{j^*}(x) + \varepsilon \log\left(\sum_{j=1}^J \nu_j\right) = \alpha_{j^*}(x) \tag{7}
\]

On the other hand, since log is increasing and all terms in the sum are non negative we have
\[
\varepsilon \log\left(\sum_{j=1}^J e^{\frac{\alpha_j(x)}{\varepsilon}} \nu_j\right) \geq \varepsilon \log\left(e^{\frac{\alpha_{j^*}(x)}{\varepsilon}} \nu_{j^*}\right) = \alpha_{j^*}(x) + \varepsilon \log(\nu_{j^*}) \xrightarrow{\varepsilon \to 0} \alpha_{j^*}(x) \tag{8}
\]

Hence \(\varepsilon \log\left(\sum_{j=1}^J e^{\frac{\alpha_j(x)}{\varepsilon}} \nu_j\right) \xrightarrow{\varepsilon \to 0} \alpha_{j^*}(x)\) and \(\varepsilon \log\left(\sum_{j=1}^J e^{\frac{\alpha_j(x)}{\varepsilon}} \nu_j\right) \leq \alpha_{j^*}(x)\).

Since we assumed \(x \mapsto c(x, y_j) \in L^1(\mu)\), then \(\alpha_{j^*} \in L^1(\mu)\) and by dominated convergence we get that \(\varepsilon \log\left(\sum_{j=1}^J e^{\frac{\alpha_j(x)}{\varepsilon}} \nu_j\right) \xrightarrow{\varepsilon \to 0} H_0(v)\).

**Proof of Proposition A.1.** First, let’s prove that \((v^\varepsilon)_\varepsilon\) has a converging subsequence. With similar computations as in Proposition 2.1 we get that \(v^\varepsilon(y_i) = -\varepsilon \log\left(\int X e^{\frac{c(x, y_i) - c(x, y_j) - \varepsilon}{\varepsilon}} \mu(\omega(x))\right)\). Denote by \(\tilde{v}^\varepsilon\) the c-transform of \(u^\varepsilon\) such that \(\tilde{v}^\varepsilon(y_i) = \min_{x \in \mathcal{X}} c(x, y_i) - u^\varepsilon(x)\). From standard results on optimal transport (see [19], p.11) we know that \(|\tilde{v}^\varepsilon(y_i) - \tilde{v}^\varepsilon(y_j)| \leq \omega(\|y_i - y_j\|)\) where \(\omega\) is the modulus of continuity of the cost \(c\). Besides, using once again the soft-max argument we can bound \(|v^\varepsilon(y) - \tilde{v}^\varepsilon(y)|\) by some constant \(C\).

Thus we get that:
\[
|v^\varepsilon(y_i) - v^\varepsilon(y_j)| \leq |v^\varepsilon(y_i) - \tilde{v}^\varepsilon(y_i)| + |\tilde{v}^\varepsilon(y_i) - \tilde{v}^\varepsilon(y_j)| + |\tilde{v}^\varepsilon(y_j) - v^\varepsilon(y_j)| \tag{9}
\]
\[
\leq C + \omega(\|y_i - y_j\|) + C \tag{10}
\]

Besides, the regularized potentials are unique up to an additive constant. Hence we can set without loss of generality \(v^\varepsilon(y_0) = 0\). So from the previous inequality yields:
\[
v^\varepsilon(y_i) \leq 2C + \omega(\|y_i - y_0\|) \tag{11}
\]

So \(v^\varepsilon\) is bounded on \(\mathbb{R}^J\) which is a compact set and thus we can extract a subsequence which converges to a certain limit that we denote by \(\tilde{v}\).

Let \(v^* \in \arg\max_{\nu} H_0\). To prove that \(\tilde{v}\) is optimal, it suffices to prove that \(H_0(v^*) \leq H_0(\tilde{v})\).

By optimality of \(v^\varepsilon\),
\[
H_\varepsilon(v^*) \leq H_\varepsilon(v^\varepsilon) \tag{12}
\]

The term on the left-hand side of the inequality converges to \(H_0(v^*)\) since \(H_\varepsilon\) converges pointwise to \(H_0\). We still need to prove that the right-hand term converges to \(H_0(\tilde{v})\).

By the Mean Value Theorem, there exists \(\tilde{v}^\varepsilon \overset{\text{def}}{=} (1 - t^\varepsilon)v^\varepsilon + t^\varepsilon \tilde{v}\) for some \(t^\varepsilon \in [0, 1]\) such that
\[
|H_\varepsilon(v^\varepsilon) - H_\varepsilon(\tilde{v})| \leq \|\nabla H_\varepsilon(\tilde{v}^\varepsilon)\| \|v^\varepsilon - \tilde{v}\| \tag{13}
\]
Figure 4: Comparisons between the Sinkhorn algorithm and SAG, tuned with different stepsizes, using different regularization strengths. The setting is identical to that used in Figure 1. Note that to prevent numerical overflow when using very small regularizations, the metric is thresholded such that rescaled costs \( c(x, y_j)/\varepsilon \) are not bigger than \( \log(10^{200}) \).

The gradient of \( H_\varepsilon \) reads

\[
\nabla_\nu H_\varepsilon(\nu) = \nu - \pi(\nu)
\]

where \( \pi_i(\nu) = \frac{\int_X e^{\nu_i - c(x, y_i)} \nu_j d\mu(x)}{\int_X \sum_{j=1}^J e^{\nu_i - c(x, y_j)} \nu_j d\mu(x)} \)

It is the difference of two elements in the simplex thus it is bounded by a constant \( C \) independently of \( \varepsilon \).

Using this bound in (13) get

\[
H_\varepsilon(\bar{\nu}) - C \| \nu_e - \bar{\nu} \| \leq H_\varepsilon(\nu_e) \leq H_\varepsilon(\bar{\nu}) + C \| \nu_e - \bar{\nu} \|
\]

By pointwise convergence of \( H_\varepsilon \) we know that \( H_\varepsilon(\bar{\nu}) \rightarrow H_0(\bar{\nu}) \), and since \( \bar{\nu} \) is a limit point of \( \nu_e \) we can conclude that the left and right hand term of the inequality converge to \( H_0(\bar{\nu}) \). Thus we get \( H_\varepsilon(\nu_e) \rightarrow H_0(\bar{\nu}) \).

B Discrete-Discrete Setting

The list of authors we consider is: KEATS, CERVANTES, SHELLEY, WOOLF, NIETZSCHE, PLUTARCH, FRANKLIN, COLERIDGE, MAUPASSANT, NAPOLEON, AUSTEN, BIBLE, LINCOLN, PAINE, DELAFONTAINE, DANTE, VOLTAIRE, MOORE, HUME, BURROUGHS, JEFFERSON, DICKENS, KANT, ARISTOTLE, DOYLE, HAWTHORNE, PLATO, STEVENSON, TWAIN, IRVING, EMERSON, POE, WILDE, MILTON, SHAKESPEARE.
C Proof of convergence of stochastic gradient descent in the RKHS

The SGD algorithm applied to the regularized continuous OT problem reads, starting with \( u_0 = 0 \) and \( v_0 = 0 \),

\[
(u_k, v_k) \overset{\text{def.}}{=} (u_{k-1}, v_{k-1}) + \frac{C}{\sqrt{k}} \nabla f_\varepsilon(x_k, y_k, u_{k-1}, v_{k-1}) \in \mathcal{H} \times \mathcal{G},
\]

where \((x_k, y_k)\) are i.i.d. samples from \( \mu \otimes \nu \). The following proposition shows that these \((u_k, v_k)\) iterates can be expressed as finite sums of kernel functions, with a simple recursion formula.

**Proposition C.1.** The iterates \((u_k, v_k)\) defined in (16) satisfy

\[
(u_k, v_k) \overset{\text{def.}}{=} \sum_{i=1}^{k} \alpha_i (\kappa(\cdot, x_i), \ell(\cdot, y_i)), \quad \text{where} \quad \alpha_i \overset{\text{def.}}{=} \Pi_B r \left( \frac{C}{\sqrt{i}} \left( 1 - e^{-\frac{u_{i-1}(x_i) + v_{i-1}(y_i) - c(x_i, y_i)}{\varepsilon}} \right) \right),
\]

where \((x_i, y_i)_{i=1\ldots k}\) are i.i.d samples from \( \mu \otimes \nu \) and \( \Pi_B \) is the projection on the centered ball of radius \( r \). If the solutions of \((D_\varepsilon)\) are in the \( \mathcal{H} \times \mathcal{G} \) and if \( r \) is large enough, the iterates \((u_k, v_k)\) converge to a solution of \((D_\varepsilon)\).

**Proof.** Rewriting \( u(x) \) and \( v(y) \) as scalar products in \( f_\varepsilon(X, Y, u, v) \) yields

\[
f_\varepsilon(x, y, u, v) = \langle u, \kappa(x, \cdot) \rangle_\mathcal{H} + \langle v, \ell(y, \cdot) \rangle_\mathcal{G} - \varepsilon \exp \left( \frac{\langle u, \kappa(x, \cdot) \rangle_\mathcal{H} + \langle v, \ell(y, \cdot) \rangle_\mathcal{G} - c(x, y)}{\varepsilon} \right).
\]

Plugging this formula in iteration (16) yields \((u_k, v_k) = u_{k-1} + \alpha_k (\kappa(\cdot, x_k), \ell(\cdot, y_k))\), where \( \alpha_k \) is defined in (17). Since the parameters \((\alpha_i)_{i<k}\) are not updated at iteration \( k \), we get the announced formula. Putting a bound on the iterates of \( \alpha \) by projecting on \( B_r \) ensures convergence [7]. \( \square \)