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HyLAR+: improving Hybrid Location-Agnostic Reasoning with Incremental Rule-based Update

Mehdi Terdjimi
Université de Lyon, LIRIS
Université Lyon 1 - CNRS
UMR5205
F-69622, France
firstname.lastname@liris.cnrs.fr

Lionel Médini
Université de Lyon, LIRIS
Université Lyon 1 - CNRS
UMR5205
F-69622, France

Michael Mrissa
Université de Lyon, LIRIS
Université Lyon 1 - CNRS
UMR5205
F-69622, France

ABSTRACT
Web applications that rely on datasets of limited sizes to handle small but frequent updates and numerous queries have no simple way to define where data should be stored and processed. We propose a reasoning framework that can be integrated in Web applications and is able to perform the same reasoning tasks on both client or server sides. This framework embeds a rule-based reasoning engine that uses an algorithm relying on both incremental reasoning and named graphs. We evaluate the performance of our approach and compare the effects of incremental reasoning and named graphs in different experimental conditions. Results show that our reasoner can significantly reduce response times to INSERT and DELETE queries. During the demo we will exhibit how it can be used to perform reasoning tasks based on client-generated information and improve Web applications with location-agnostic reasoning.

Categories and Subject Descriptors
[Software and its engineering]: Contextual software domains—Software infrastructure[Middleware][Embedded middleware]; [Software and its engineering]: Software organization and properties[Software system structures][Distributed systems organizing principles][Organizing principles for web applications]; [Computing methodologies]: Artificial intelligence[Knowledge representation and reasoning]

General Terms
DESIGN, PERFORMANCE, EXPERIMENTATION

Keywords
semantic web, reasoning, client-side reasoning, rule-based reasoning

1. INTRODUCTION
More and more Web applications generate and use datasets of limited sizes. This is the case in domains such as personal small data [2] or Web of Things [3], and more generally for applications that keep track of individual user’s activities to adapt their behavior. Such datasets are subject to small but frequent updates and numerous queries. However, depending on the data volumes, process complexity and client capabilities, there is no simple way to define where data should be stored and processed and make appropriate choices at design time. Especially when semantic reasoning comes into play, the diverse reasoning tasks involved in bootstrapping a reasoner (classification), updating data (insert, delete and re-insert) and querying (select) must be taken into account wrt. the client processing capabilities. To partly cope with this issue, we designed and proposed the Hybrid Location-Agnostic Reasoner (HyLAR) architecture [8]. HyLAR is a JavaScript-based reasoning framework that can deploy different parts of a reasoning process either on the server or on the client.

In [8], we found out that in some cases, ontology classification could be performed on the server and the classified ontology could be shipped to the client so that querying could be performed locally. However, the limited OWL2-EL reasoner embedded in our HyLAR framework was only designed to perform these two tasks and data update was not taken into account. In this paper we propose an improvement of the HyLAR framework. It relies on an OWL2-RL reasoner that is able to perform (re)-insert and delete operations, thus providing a wider range of use to this framework. Moreover, this reasoner includes optimizations such as incremental reasoning [4] and named graphs support.

The paper is organized as follows. In Section 2, we review related work on reasoning optimization techniques and adaptive reasoning. We present our approach together with its implementation in Section 3 and our evaluation in 4. We summarize and discuss our results to give guidelines for future work in Section 6.

2. RELATED WORK
2.1 Reasoning in OWL profiles
OWL 2 profiles1 allow adjusting the trade-off between expressive power and efficiency. Each profile favors specific reasoning tasks in terms of complexity [7]. EL is suitable for very large TBoxes, such as in biomedical ontologies and performs well with ontology consistency, class expression

1http://www.w3.org/TR/owl2-profiles/
subsumption and instance checking. QL is appropriate for applications that manipulate high volumes of instances and targets query answering. RL relies on rule languages to offer a reasonable trade-off between expressive power and scalable reasoning for ontology consistency, class expression satisfiability, class expression subsumption, instance checking and conjunctive query answering [7]. As instance retrieval is the most important inference task in RL [6], most RL reasoners pre-compute and explicitly store inferences, so that queries can be answered by querying the knowledge base (KB) [5]. One optimization of OWL2-RL reasoners is the Incremental Reasoning (IR) approach presented in the DRed algorithm from [4, Section 7]. It avoids recalculating all consequences when partly updating the KB. To do so, facts in the KB are tagged as explicit (when directly inserted in the KB) or implicit (when derived from rules). At update time, the algorithm first performs an over-deletion step, to remove all explicit facts that are stated to be removed in the query, along with their inferred consequences. Then a rederivation step deduces all consequences of newly added explicit facts and inserts both new explicit and implicit facts in the KB. In order to reduce computation times, the algorithm performs these two steps over a limited selection of rules in the causes of which the explicit facts to be deleted/inserted appear.

2.2 Reasoning distribution

Several studies have been conducted to optimize reasoning tasks, from classical model in which all the semantic data are gathered, cleaned, lifted and stored on the server side to full client-side reasoning. On the server side, SPARQL federated queries provide a solution to handle the query answering step. The ANAPSID [1] query engine relies on several endpoints in order to split the data, decompose queries and select appropriate sources. The Triple Pattern Fragments (TPF) [9] approach is based on a principle close to federated querying but involves the client in the answering process. It allows “intelligent clients” to query TPF servers for triples that match a particular triple pattern and therefore reduce endpoints workload.

While several embedded reasoners exist to target mobile applications, most of them have been developed and optimized to fit certain languages and domain needs. To the best of our knowledge very few client-side reasoners are available for Web applications. EYE2 is a NodeJS3-compatible reasoner capable of inferring on FOL rules. EYE performs server-side reasoning and has not been ported onto the client side. Based on the JSW framework, OWLReasoner4 allows client-side processing of SPARQL queries on OWL2-EL ontologies. It can perform ontology classification and converts the TBox and ABox into an internal relational database. SPARQL queries sent to the reasoner are rewritten into SQL queries and processed on the database. Its SPARQL engine is limited to basic rule assertions. HyLAR [8] is a framework that allows hybrid OWL2-EL reasoning. That means that both server side (Node.js) and client side (AngularJS5) can execute reasoning tasks using the same parts of code. Reasoning in HyLAR relies on the OWLReasoner engine.

3. HyLAR+

We propose HyLAR+6, a new version of the HyLAR framework. We refactored its architecture and made it embed a JavaScript OWL2-RL reasoner that enables rule-based transformation, incremental reasoning and named graphs support. The HyLAR+ architecture is depicted in Figure 1 and comprises two main components: a parser and a reasoner. Both of them have been reimplemented to be faster and provide more reasoning capabilities (i.e. more constructs) than those of the JSW framework. The parser prepares the statements for the reasoner. It can now take two kinds of inputs: RDF/XML ontology files and SPARQL queries (INSERT DATA, DELETE DATA or SELECT). This allows bootstrapping the reasoner with an initial ontology, updating its contents when data sources change and querying it. For the needs of our evaluation, the reasoner embeds two different algorithms: a regular (a.k.a. “greedy”) algorithm that processes all available rules on the whole KB at each request and the IR algorithm described in Section 2.1. Each reasoning task can be performed with any of these two algorithms. The rule base is divided into OWL2-RL (see below) and domain-specific rules. The KB stores facts (i.e. triples) in a local relational database and converts SPARQL request into SQL using Trim Path7.

3.1 RL partial support and constructs

Our reasoner currently supports partial OWL2-RL semantics. The currently implemented rules are shown in Figure 2. In addition to the previously supported class subsumption and Object properties, HyLAR+ Parser and Reasoner modules now support Datatypes and Datatype properties, as well as additional OWL constructs: ObjectSomeValuesFrom, DataSomeValuesFrom, InverseObjectProperties, ObjectPropertyRange, ObjectPropertyDomain and DataPropertyRange. More rules can be added depending on the expressivity level needed by the application.

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4http://reasoning.restdesc.org/
5https://angularjs.org/
6Repo: https://github.com/ucbl/HyLAR/tree/IR-NG
7Online demo: http://dataconf.liris.cnrs.fr/hylar
8https://code.google.com/p/trimpath/
3.2 SPARQL and named graph support
HyLAR+ relies on SPARQL to allow updating and querying the KB all along the application lifecycle. While the previous HyLAR version only supported SELECT queries, our framework now additionally supports INSERT DATA and DELETE DATA queries. It also handles named graphs for these three types of queries. Query answering is performed as follows. The reasoner receives as input a SPARQL query parsed into facts (triples) and an already classified graph representing an ontology or a named graph. It applies OWL2-RL rules according to the chosen algorithm and iterates until no more facts can be inferred. The same operation is done with domain-specific rules, as processing inferences from the application domain must be done after deducing all OWL consequences of updated facts. Depending on the statement, it converts newly inferred implicit facts into an SQL query in order to update or select data from the KB.

4. EVALUATION
The location-agnostic feature of the HyLAR framework has already been evaluated in [8]. The reasoners embedded in the previous version do not follow the same OWL profile. Hence, as stated in Section 2.1, they do not cope with the same reasoning tasks and cannot be compared. In this section, we evaluate the performance of HyLAR+ w.r.t. two different aspects. First, we compare the execution times of the IR and greedy algorithms for ontology classification, as well as for INSERT DATA, DELETE DATA and SELECT statements. Second, we measure the impact of named graphs on these two algorithms while performing different SPARQL queries, in comparison to querying the whole KB.

In both evaluations, we use the FIPA-Device ontology\(^9\) to calculate the classification times for both reasoning algorithms, once the ontology file has been parsed into a set of axioms and triples. This ontology comprises 14 classes, 14 object properties and 18 datatype properties as well as 67 axioms to express subsumptions, domains and ranges. We ran our evaluation on a Dell Inspiron, i7-2670QM CPU @ 2.20GHz. Computation times are in milliseconds.

4.1 Incremental reasoning evaluation
As IR reduces computing using rule set restrictions, we chose the number of RL rules (see 3.1) as varying parameter. During this evaluation, we progressively added the following set of rules and applied them on the FIPA ontology:

\[
\begin{align*}
R0 &= \emptyset \\
R1 &= \{\text{cax-cls}\} \\
R2 &= R1 \cup \{\text{cax-sco}\} \\
R3 &= R2 \cup \{\text{cax-cls}\} \\
R4 &= R3 \cup \{\text{cax-eqc1}\} \\
R5 &= R4 \cup \{\text{cax-eqc2}\} \\
R6 &= R5 \cup \{\text{eq-trans}\} \\
R7 &= R6 \cup \{\text{eq-syn}\}
\end{align*}
\]

In Figure 3 we notice a gap between R2 and R3 which is due to our implementation choice: rules are currently represented with only one consequence. Thus, \text{cax-cls} splits into four sub-rules corresponding to its four consequences. Starting at R3, the IR classification time stabilizes, while the greedy algorithm increases as new rules are added to the set. We can assume that R3 is the point at which no more information in FIPA-Device is inferable from others rules. IR is more suitable in this situation as it does not compute over unnecessary rules.

4.2 Reasoning over named graphs
In order to measure the influence of named graphs on computation time, we evaluated the processing times for insertion, deletion and selection over the FIPA-Device ontology and subsumption rules from R2 (\text{cax-sco} and \text{cax-cls}). We consider four cases: IR (with no graph), IR (G) (on a named graph), GR (greedy reasoning with no graph) and GR (G) (greedy reasoning on a named graph). We evaluated these cases on 10, 20, 30, 40 and 50 triples. These data sets are available online, together with the framework.

The same behavior is observed on insertion and deletion. Without graphs, the query processing times increase faster as more triples are added/deleted, in comparison to updates using named graphs. At a certain point (50 triples when inserting and 20 triples when deleting), the greedy algorithm using named graphs is still better than the incremental without any graph. As expected, the best results are given by the IR with named graphs. For SELECT queries, processing times are similar with or without named graphs.

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\(\text{http://www.fipa.org/specs/fipa00091/PC00091A.html}\)


9http://chrjs.net/