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An efficient one-step-ahead optimal control for urban signalized traffic networks based on an averaged Cell-Transmission model

Pietro Grandinetti, Carlos Canudas de Wit and Federica Garin

Abstract—This paper presents a model for large urban traffic networks, based on the well-known macroscopic Cell Transmission Model. We start by describing the dynamics of traffic flow at signalized intersections. Then we develop an average-based approximation of such a system, that we use to build our control algorithm as a linear optimization problem. Simulation results validate the averaged approach and show the effectiveness of the proposed control strategy.

I. INTRODUCTION

Traffic congestion on urban roads is a problem of great interest nowadays since it strongly affects security and pollution; effective and easy-to-handle models are therefore needed to represent and control traffic behavior. The scientific community relies on macroscopic models of time-space evolution of the traffic. Such models describe traffic as a fluid, and are based on a mass conservation law [1].

With respect to microscopic models, macroscopic ones are preferred due to their simplicity in characterizing vehicles’ flows and densities. The Cell Transmission Model (CTM) [2] is an example, widely used, of this kind of representation.

In the last decades the major interest has been in analyzing and controlling highway traffic, and only recently the attention is moving on urban networks. The CTM’s extension to networks was proposed in [3], and in [4] it has been shown the existence of a macroscopic fundamental relation between flow and density even for urban roads, with experimental validation.

In [5, 6] outflow is defined independently from downstream supply and it is then equal to a link’s demand. This is unrealistic for traffic evolution, as pointed out in [7], since in such a way upstream flow is not affected by downstream congestion. Solutions to this problem are proposed in [7, 8].

Our contribution regarding traffic modelling consists first in the formulation of the traffic’s dynamics as an extended CTM with FIFO policy at intersections; based on this, we then define an approximated averaged model, provide numerical validation of its quality, and use it as a tool to design traffic lights’ control. Although the idea of representing traffic lights as the percentage of their green time (duty cycle) has been already employed in scientific works (e.g., [8, 9, 10]), authors employing this technique often base all their work on this approximated model, while we employ the averaged model as intermediate tool to design the control actions, which will then be actuated in the non-averaged network.

Our choice of traffic lights as control devices is motivated by a practical consideration: they are nowadays the main reason of more or less efficient network operations, even though they were originally installed only in order to guarantee the safe crossing.

Urban traffic control strategies are classified as fixed–time techniques [11, 12] and model–based algorithms [13, 14]. The main drawback of the former ones is that their settings are based on historical rather than real–time data, while the latter ones basic problem is that they require algorithms with exponential complexity for a global optimization. A survey about the existing techniques is [15]. More advanced schemes have been presented recently and they refers to different models for the network and for the chosen control actions, such as max pressure control [16] and cooperative green lights policies [9]. These works consider control of intersection stages that maximizes the throughput and analyzes stability property of the network for given static demands. Conversely, the contribution of this paper regarding traffic control is the design of a strategy that optimizes in real time traffic performance and deals with continuously changing demands. Furthermore, our algorithm is formulated as a linear optimization, and it is therefore computationally very efficient.

The rest of the paper is organized as follows: Section II introduces our definition of urban network and the proposed model, including an approximation of such a model based on the average theory. Section III shows the outcomes of our validation regarding the consistency of the averaged model; Section IV discusses the choices of performance index for urban traffic while Section V describes our control strategy and presents some result of software simulations. Conclusions and future works are outlined in Section VI.

II. SIGNALIZED NETWORK MODEL

We consider urban traffic networks as sets of roads, in which the vehicles’ flow passing from one road to another is regulated by traffic lights. These lights, where present, are located at the end of the roads. The crossing points, called intersections, have no capacity storage. We introduce here some preliminary notation necessary for further developments.

The roads of an urban network are a collection of three
sets $\mathcal{R}^{\text{in}}, \mathcal{R}, \mathcal{R}^{\text{out}}$, with associated the following relations:
\[
\begin{align*}
\text{prev}: & \quad \mathcal{R} \cup \mathcal{R}^{\text{out}} \rightarrow \mathcal{R}^{\text{in}} \cup \mathcal{R} \\
\text{next}: & \quad \mathcal{R}^{\text{in}} \cup \mathcal{R} \rightarrow \mathcal{R} \cup \mathcal{R}^{\text{out}},
\end{align*}
\]
where
- $\mathcal{R}$ is the set of the inner roads of the network;
- $\mathcal{R}^{\text{in}}(\mathcal{R}^{\text{out}})$ is the set of the roads entering (exiting) the network;
- $\text{next}(r) \ (\text{prev}(r))$ is the set of roads connected downstream (upstream) to the road $r$.

Inside the network every road in $\mathcal{R}^{\text{in}} \cup \mathcal{R}$ has at least one road downstream, while every road in $\mathcal{R} \cup \mathcal{R}^{\text{out}}$ has at least one road upstream. If $r \in \text{next}(q)$ we say that $r$ and $q$ are connected to the same intersection, where $r$ is incoming and $q$ is outgoing. Note that the $\text{prev}$ relation is undefined for the roads entering the network, while the next relation is undefined for the exiting ones.

A. CTM signalized model

To allow regulation of the traffic we introduce traffic lights as functions of the time, i.e., there exists a function
\[
\alpha : \mathbb{R}_+ \rightarrow \{0, 1\}^{|\mathcal{R}^{\text{in}} \cup \mathcal{R}|},
\]
where $\mathbb{R}_+$ is the set of all time instants. Thus, at each time instant $t$, $\alpha(t)$ is a vector of $|\mathcal{R}^{\text{in}} \cup \mathcal{R}|$ values which give red (0) or green (1) to each road in $\mathcal{R}^{\text{in}} \cup \mathcal{R}$. Traffic lights operation is time–cyclic, with cycle length $T$. Typical values of $T$ are usually around 2 minutes.

To ensure the right of way (r.o.w.) we assume that:
\[
\forall q \in \mathcal{R} \cup \mathcal{R}^{\text{out}}, \quad \forall t \in \mathbb{R}_+, \sum_{r \in \text{prev}(q)} \alpha_r(t) = 1. \tag{4}
\]

We will now model the behavior of traffic flow at intersections generalizing the diverge case discussed in [3]. There is no explicit need to model the merge of flows, since they are regulated by constraint (4). In our settings each road is considered as a cell of the CTM. We further assume that associated to the network there is an array of split ratios $\beta$, where each element $\beta_r \in (0, 1)$ represents the percentage of vehicles entered the intersection upstream to $r$ that want to go in $r$. It must be:
\[
\forall q \in \mathcal{R}^{\text{in}} \cup \mathcal{R}, \sum_{r \in \text{next}(q)} \beta_r = 1. \tag{5}
\]

**Remark 1.** The values $\beta$s are supposed to be given. They may change in time, according to the network’s status, but we do not consider them as variables for traffic regulation, since they are rather an indication of drivers’ intentions.

We now briefly recall the demand & supply paradigm [17], as it is necessary to describe the traffic evolution.

**Definition 1.** For a road $r$ we call:
- demand of $r$ ($D_r$) the flow of vehicles that want to exit $r$;
- supply of $r$ ($S_r$) the flow that can be received by $r$.

Analytical expressions for demand and supply are given through saturated functions of the density:
\[
\begin{align*}
D_r(\rho_r(t)) & = \min\{v_r\rho_r(t), \phi_r^{\text{max}}\} \tag{6a} \\
S_r(\rho_r(t)) & = \min\{\phi_r^{\text{max}}, w_r(\rho_r^{\text{max}} - \rho_r(t))\}, \tag{6b}
\end{align*}
\]
where road $r$ is characterized by the following parameters: $v_r$, the maximum speed in freeflow; $w_r$, the speed in congestion phase; $\rho_r^{\text{max}}$, the maximum density allowed; $\phi_r^{\text{max}}$, the maximum flow. The lowest value of density where the demand reaches its saturation value is called critical density ($\rho_c$).

According to the fundamental triangular diagram (see Figure 1) the relation between density and flow inside a road can be expressed as:
\[
f_r(t) = \min\{v_r\rho_r(t), w_r(\rho_r^{\text{max}} - \rho_r(t))\}. \tag{7}
\]

From the mass conservation law, the density (normalized w.r.t. the length of the road) evolves with the following rule:
\[
\dot{\rho}_r(t) = f_r^{\text{in}}(t) - f_r^{\text{out}}(t), \tag{8}
\]
where both $f_r^{\text{in}}$ and $f_r^{\text{out}}$ depend on the densities of different roads (the ones involved in the same intersection).

Consider now a road $r$ entering some intersection: the outflow of $r$ will be the maximum flow that respects the constraints given by $D_r$, as well as by the supply of the roads exiting the same intersection (we will introduce later the r.o.w.), i.e.,
\[
f_r^{\text{out}}(t) = \max \phi \\
\text{s.t. } \phi \leq D_r(\rho_r(t)) \\
\beta_j \phi \leq S_j(\rho_j(t)) \quad \forall j \in \text{next}(r). \tag{9}
\]

The solution of problem (9) is given by
\[
f_r^{\text{out}}(t) = \min \left\{ D_r(\rho_r(t)), \left\{ \frac{S_j(\rho_j(t))}{\beta_j} \right\}_{j \in \text{next}(r)} \right\}. \tag{10}
\]
Thus the outflow of road $r$ will be the flow given by (10) if $r$ has r.o.w. in the downstream intersection, zero otherwise. Equation (8) becomes:
\[
\dot{\rho}_r(t) = f_r^{\text{in}}(t) - \alpha_r(t)f_r^{\text{out}}(t). \tag{11}
\]

The inflow in $r$ from the upstream intersection will be:
\[
f_r^{\text{in}}(t) = \beta_r \sum_{q \in \text{prev}(r)} \alpha_q(t)f_q^{\text{out}}(t). \tag{12}
\]
Remark 2 (Boundary flows). Equation (10) is valid if \( r \in \mathbb{R}^{in} \cup \mathbb{R} \). If, instead, \( r \in \mathbb{R}^{out} \) then its outflow is bounded from the supply outside the network \((S^\text{out}_t)\), i.e.,
\[
f^\text{out}_r(t) = \min \{ D_r(\rho_r(t)), S^\text{out}_r(t) \}. \tag{13}
\]
Similarly, equation (12) is not applicable if \( r \in \mathbb{R}^{in} \), where the inflow is related to the external demand \((D^\text{in}_r)\):
\[
f^\text{in}_r(t) = \min \{ D^\text{in}_r(t), S_r(\rho_r(t)) \}. \tag{14}
\]
Note that the model we have defined so far is consistent with the property of positive invariance [7], thanks to the demand/supply paradigm and to the constraint (4). This guarantees the consistency of the model: it will never happen that the density evolves outside the prescribed physical limits.

B. Averaged model

An interesting simplification of the model presented can be obtained applying the average theory [18, Chapter 8]. This technique looks forward for both model simplification and control strategies design. As motivations for this consider that combinatorial problems (given by the binary behavior) are usually difficult to address; Moreover, regulation of smoother variables (e.g., duty cycle) fits better to traffic scenarios, rather than fast switching between binary values. The time-averaged of the density \((\alpha)\) is:
\[
\hat{\alpha}(t) = \frac{1}{T} \int_t^{t+T} \alpha(\tau) d\tau = \frac{1}{T} \int_t^{t+T} (f(\tau) - \alpha(\tau) f^\text{out}(\tau)) d\tau. \tag{15}
\]
In (15) the only \( T \)-cyclic signal is \( \alpha \), that can be averaged over the period as follows:
\[
\bar{\alpha}(t) \triangleq \frac{1}{T} \int_t^{t+T} \alpha(\tau) d\tau = \frac{G_\alpha}{T}, \tag{16}
\]
where \( G_\alpha \) represents the green time of the traffic light \( \alpha \). Thus \( \bar{\alpha}(t) \) is the duty cycle of traffic light \( \alpha \) during the cycle \([t, t+T]\).

We then choose to approximate the evolution of the real density with the following:
\[
\hat{\rho}_r(t) = f^\text{in}_r(t) - \bar{\alpha}(t) f^\text{out}_r(t). \tag{17}
\]
Equation (17) represents a different dynamical system which aims to approximate (11). The flows involved in (17) are then functions of the new state variable \( \bar{\rho} \). The respective of (12) will be:
\[
f^\text{in}_r(t) = \beta_r \sum_{q \in \text{prev}(r)} \bar{\alpha}(t) f^\text{out}_q(t). \tag{18}
\]
In the averaged approximation the oscillating behavior of the system (due to the green/red alternation) is lost. A limitation of this approach is that possible phase shifts are not captured; a study of representation (and choice) of traffic lights’ phase shifts is beyond the scope of this paper.

![Standard 4-ways intersection](image)

**Fig. 2:** Standard 4-ways intersection.

The full averaged model is given by the following set of equations:
\[
\begin{align*}
\hat{\rho}_r(t) &= f^\text{in}_r(t) - \bar{\alpha}(t) f^\text{out}_r(t), \quad r \in \mathbb{R}^{in} \cup \mathbb{R} \cup \mathbb{R}^{out} \quad (19a) \\
\hat{f}_r^\text{out}(t) &= \begin{cases} 
\min \{ D_r(t), S^\text{out}_r(t) \}, & r \in \mathbb{R}^{out} \\
\min \left\{ D_r(t), \left\{ \frac{S_r(t)}{\beta_r} \right\}_{j \in \text{next}(r)} \right\}, & \text{others.} \tag{19b}
\end{cases} \\
\hat{f}_r^\text{in}(t) &= \begin{cases} 
\min \{ D^\text{in}_r(t), S_r(t) \}, & r \in \mathbb{R}^{in} \\
\min \left\{ \beta_r \sum_{q \in \text{prev}(r)} \bar{\alpha}_q(t) f^\text{out}_q(t) \right\}, & \text{others.} \tag{19c}
\end{cases} \\
D_r(t) &= \min \{ v_r \bar{\rho}_r(t), \varphi_r^\text{max} \} \quad (19d) \\
S_r(t) &= \min \{ \varphi_r^\text{max}, w_r(\rho_r - \bar{\rho}_r(t)) \}, \tag{19e}
\end{align*}
\]
with the constraints:
\[
\begin{align*}
\forall q \in \mathbb{R} \cup \mathbb{R}^{out} & \quad \sum_{r \in \text{prev}(q)} \bar{\alpha}_q(t) = 1 \quad (20) \\
\forall q \in \mathbb{R}^{in} \cup \mathbb{R} & \quad \sum_{r \in \text{next}(q)} \beta_r = 1. \quad (21)
\end{align*}
\]

Remark 3. Note that the demand and supply paradigm is fulfilled even in the averaged system: in fact, thanks respectively to (19b) and (19c), outflow and inflow for a road \( r \) will satisfy the following for every \( t \):
\[
\bar{\alpha}_r(t) f^\text{out}_r(t) \leq D_r(t) \tag{19f}
\]
\[
f^\text{in}_r(t) \leq \beta_r \sum_{q \in \text{prev}(r)} \bar{\alpha}_q(t) f^\text{out}_q(t) \tag{19g}
\]

III. ILLUSTRATIVE EXAMPLE AND VALIDATION

We have built, via software simulation, a network with 40 roads connected by standard 4-ways intersections. A 4-way intersection (see Figure 2), is described by the averaged system with the following set of equations:
\[
\begin{align*}
\hat{\rho}_1 &= f_1^\text{in} - \bar{\alpha}_1 f_1^\text{out} \\
\hat{\rho}_2 &= f_2^\text{in} - \bar{\alpha}_2 f_2^\text{out} \\
\hat{\rho}_3 &= \beta_3 (\bar{\alpha}_1 f_1^\text{out} + \bar{\alpha}_2 f_2^\text{out}) - f_3^\text{out} \\
\hat{\rho}_4 &= \beta_4 (\bar{\alpha}_1 f_1^\text{out} + \bar{\alpha}_2 f_2^\text{out}) - f_4^\text{out},
\end{align*}
\]
where \( f_1^\text{in}, f_2^\text{in}, f_3^\text{out}, f_4^\text{out} \) are the flows to be connected to other roads in the network, and constraint (4) implies \( \bar{\alpha}_1, \bar{\alpha}_2 \). For ease of visualization we assume all roads have same features: \( \varphi_r^\text{max} = \varphi_r^\text{max}, \beta_r^\text{max} = \beta_r^\text{max}, v_r = v, w_r = w \), for every road \( r \). We are interested in:

- Capturing the modes of the network at each time instant: a road may be either free (i.e., with density lower than the critical density) or congested (with different levels of
congestion), and we want our approximation to capture rightly as many road–modes as possible;

- Overall precision: how much the averaged system is close to the actual one, knowing that it does not present any fast switching behavior.

The simulation has been run in the following scenario: each traffic light (at each intersection) is a periodic given signal and the split ratios cause some asymmetry in the traffic diffusion. Outside of the network, we have demands (supplies) varying in time (randomly generated in $[0, 0.5, 1]$) for all the entering (exiting) roads. Representative examples of our results are reported in Figures 3–4. Note that:

- The averaged system succeeds in capturing rightly an high percentage of modes. The mean error is around 10%, as it results from Figure 4c. Figures 3a and 3b show two instantaneous pictures of the modes over the entire network (for both actual and averaged systems);
- The precision of the averaged model results to be fully satisfying in approximating the density over all roads, as shown in Figure 4a and 4b, where the densities are very similar and the only significant difference is that oscillations in 4a are switched off in 4b, as expected.

IV. URBAN TRAFFIC PERFORMANCE METRICS

Traffic behavior needs to be evaluated with respect to performance indices properly defined. There exist several metrics in literature to address performance evaluation; in this paper we focused on the following two features.

A. Service of Demand (SoD)

An urban traffic network is an highly dynamical environment that continuously receives demand from outside. This demand cannot be ignored just to favour the inner quality of the system, because the external request will end up growing with several undesired effects, due to the bigger and bigger queues arising outside.

For this reason we define as quality of the service the number of vehicles (users) served:

$$\text{SoD}(t) = \int_0^t \sum_{r \in \mathcal{R}} f_r^\text{in}(\tau) d\tau = \int_0^t \sum_{r \in \mathcal{R}} \min\{D_r^\text{in}(\tau), \varphi_r^\text{max}, w_r(\rho_r^\text{max} - \rho_r(\tau))\},$$

where $f_r^\text{in}$ is the boundary flow defined by (14) and $D_r^\text{in}$ is the external demand for road $r$. The quantity expressed in (23) is a value that we would like to maximize, since this is equivalent to decrease the sum of the queue lengths.

B. Optimization of the infrastructures usage

In urban networks some road is preferred than others by the users. The infrastructure holder would like to set traffic lights as to diminish this usage disparity, to guarantee a more equilibrate diffusion of vehicles, thus reducing hard congestions in main streets as well as the possibility of accidents.

A standard metric that takes into account this behavior is the Total Travel Distance [19], a cumulative index defined as:

$$\text{TTD}(t) = \int_0^t \left( \sum_{r \in \mathcal{R} \cup \mathcal{R}^\text{out}} f_r(\tau) \right) d\tau,$$

where $f_r$ is the flow inside the road $r$, as it has been defined in (7).

TTD is a measure that should be maximized as well.
as at time $t$ both $\text{SoD}(t)$ and $\text{TTD}(t)$ are known values.

The control problem can be described with the following optimization:

$$\alpha(t) = \arg \max_{\alpha} \sigma_1 \text{SoD}^+(\alpha) + \sigma_2 \text{TTD}^+(\alpha)$$

$$\forall \alpha \in \mathcal{A}$$

$$\rho_r(t + \Delta t) = \rho_r(t) + \Delta t \left( f_r^\text{in}(t) - \rho_r f_r^\text{out}(t) \right)$$

for every $r \in \mathcal{R} \cup \mathcal{R}^\text{out}$,

where $f_r^\text{in}$ and $f_r^\text{out}$ are given by (19b)–(19e), while $\sigma_1, \sigma_2 \geq 0$ are weights for the involved indices.

By definition of duty cycle $\alpha(t) \in [0, 1]$ for every road. These constraints, together with (4), give the convex domain $\mathcal{A}$ in the problem (27).

Remark 4 (Control scheme). The control first computes the values of $f_r^\text{in}(t)$ and $f_r^\text{out}(t)$, using the measured densities; then it solves the optimization problem, so computing optimal duty cycles for the upcoming cycle and finally sets them in the network. This scheme is repeated at the beginning of every cycle and not more often, in order to avoid too many switching in the traffic lights’ signals.

We now show that problem (27) is truly a linear program. It can be rewritten substituting the explicit expressions of (25) and (26), and it results (we rename $\alpha(t) \equiv \rho_r(t + \Delta t)$ for sake of brevity):

$$\alpha(t) = \arg \max_{\alpha} \left( \sigma_1 \sum_{r \in \mathcal{R}^\text{in}} \min \{ D_r^\text{in}(t), \varphi_r^\text{max}, \nu_r(\rho_r^\text{max} - \rho_r^\text{in}(\alpha(r))) \} \right)$$

$$\sigma_2 \sum_{r \in \mathcal{R}^\text{in} \cup \mathcal{R}^\text{out}} \min \{ \nu_r \rho_r^\text{in}(\alpha(r)), \nu_r(\rho_r^\text{max} - \rho_r^\text{in}(\alpha(r))) \}$$

$$\forall \alpha \in \mathcal{A}$$

$$\alpha(t) = \rho_r(t) + \Delta t \left( f_r^\text{in}(t) - \rho_r f_r^\text{out}(t) \right)$$

(28)

which is a convex problem. We may rewrite it as a linear optimization, doubling the number of variables, in the following way:

$$\arg \max_{\pi, \mu, \nu} \left( \sigma_1 \sum_{j \in \mathcal{R}^\text{in}} \mu_j + \sigma_2 \sum_{k \in \mathcal{R}^\text{in} \cup \mathcal{R}^\text{out}} \nu_k \right)$$

$$\begin{cases} 
\mu_j \leq D_j^\text{in} \\
\mu_j \leq \varphi_j^\text{max} \\
\mu_j \leq w_j(\rho_j^\text{max} - \rho_j^\text{in}(\alpha)) \\
\nu_k \leq \nu_k \rho_k^\text{in}(\alpha) \\
\nu_k \leq \nu_k (\rho_k^\text{max} - \rho_k^\text{in}(\alpha)) \\
\alpha \in \mathcal{A} 
\end{cases}$$

(29)

$$\alpha(t) = \rho_r(t) + \Delta t \left( f_r^\text{in}(t) - \rho_r f_r^\text{out}(t) \right)$$

(28)

which is a convex problem. We may rewrite it as a linear optimization, doubling the number of variables, in the following way:

Proposition 1. Problem (28) is equivalent to problem (29).

Proof. Notice that problem (29) is a relaxation of problem (28) and its constraints concerning variables $\mu_j$ and $\nu_k$ can be written as

$$\mu_j \leq \min \{ D_j^\text{in}, \varphi_j^\text{max}, w_j(\rho_j^\text{max} - \rho_j^\text{in}(\alpha)) \}$$

$$\nu_k \leq \min \{ \nu_k \rho_k^\text{in}(\alpha), \nu_k (\rho_k^\text{max} - \rho_k^\text{in}(\alpha)) \}$$

(30)
Hence the statement is true if at the optimum both (30) and (31) are satisfied as equality, $\forall j, k$. We then complete the proof by contradiction. Suppose $\mu^*, \nu^*$ is optimal and there exist a real positive number $\epsilon$ and an integer $i$ such that

$$\mu_i^* + \epsilon = \min \left\{ \{\rho_i^\infty, \varphi_i^\max, w_i(\rho_i^\max - \rho_i^+ (\pi)) \right\} \right.$$

(32)

(the same reasoning applies if the variable in point is some $\nu_j^*$). Then we may add $\epsilon$ to the variable $\mu_i^*$ and we will obtain a new admissible solution (note that in (29) $\mu_i$ is uncorrelated with $\mu_j$, $i \neq j$). Because of the maximization purpose, such a new solution will be better than the previous, which, therefore, could not be optimal.

**B. Simulation and results**

To test the proposed control algorithm we have run simulations using the software [20] to solve the optimization problem. We stress again the fact that the control is actuated on the whole simulation horizon and TTD are evaluated over the whole simulation horizon individually. We stress again the fact that the control is actuated on the whole simulation horizon and TTD are evaluated over the whole simulation horizon. The metrics SoD and TTD are evaluated over the whole simulation horizon and compared with the ones obtained with a fixed-light policy based only on the split ratios (and not on real-time optimization). Representative results are shown in Figure 6 and in Table I. Note that:

- Our algorithm achieves good performance regarding the optimization of the infrastructures usage. Figure 6 shows how far each road’s density is from its best working point $\rho_c$, where lighter color means better performance;
- Table I gives the quantitative measures of the improvement, which is positive for both the choosen indices.

**VI. Conclusions and future works**

In this paper we presented a model for urban signalized traffic networks and its average-based approximation for control purposes. Numerical validations have confirmed the reliability of such approximation. We also proposed a numerically efficient control scheme based on global optimization, whose application in simulation has given good performance improvements.

Future research will aim to design multi-step-ahead control, combined with techniques for traffic lights phase shift.

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