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Towards Modular Instrumentation of Interpreters in JavaScript

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Abstract
With an initial motivation based on the security of web applications written in JavaScript, we consider the instrumentation of an interpreter for a dynamic analysis as a crosscutting concern. We define the instrumentation problem – an extension to the expression problem with a focus on modifying interpreters. We then illustrate how we can instrument an interpreter for a simple language using only the bare language features provided by JavaScript.

Categories and Subject Descriptors D.3.3 [Programming Languages]: Language Constructs and Features

Keywords JavaScript, Instrumentation, Interpreter, Expression Problem, Modularity

1. Introduction
Today, JavaScript, with its special connection to HTML5 [6], is considered the language of choice for building web applications. It has made it possible to turn the static pages of the early web into dynamic pages, providing effective interactions with remote sites. This context has created new security challenges, in particular, with respect to information flow security and its dual aspects of confidentiality (local private information should not flow to untrusted remote sites) and integrity (remote sites should not corrupt local information). Although these properties are traditionally enforced through static analysis [12], this new context has created the need for considering dynamic analyses as well as hybrid analyses with various theoretical and practical results [3]. In practice, the corresponding prototypes, when they exist and have been made accessible, turn out to be built each in a very specific way, making it hard to compare and reuse approaches. In particular, the integration of the analysis is often invasive, which means that it is even difficult to relate the principles of the analysis and the implementation or evolve either the analysis or the base JavaScript runtime to which it applies.

This can be seen as a problem of modular instrumentation: what are the basic principles making it possible to modularly compose a JavaScript runtime and a dynamic or hybrid analysis dealing with information flow? In this paper, we study this problem starting from a full-blown case study: the extension of Narcissus, a JavaScript interpreter, written by Mozilla, to faceted evaluation [2], a novel and expressive information flow analysis technique able to dynamically handle implicit flows (flows depending on conditionals).

On the basis of this case study, we define the four criteria that would enable the modular instrumentation of interpreters. We frame these criteria as the instrumentation problem – an extension to the expression problem (section 2). We then illustrate how to build an interpreter incrementally, and how to instrument it in a way that satisfies the requirements of the instrumentation problem using only simple features of the JavaScript language (section 3). Finally, we discuss how our proposal relates to other approaches to modular instrumentation (section 4).

2. Instrumenting Interpreters For Dynamic Analyses
Instrumenting an interpreter for a dynamic analysis is straightforwardly achieved by modifying the source code of the interpreter. The dynamic analysis logic may require to extend or disable parts of the existing interpreter code, in several, separate places. When the code of the instrumentation is split across the interpreter code, the concerns of the dynamic analysis and the concerns of the language interpretation become tangled. As a result, extensibility, maintainability and reusability are all decreased; and modular reasoning is lost.

2.1 Case Study: Narcissus Instrumentation for Faceted Evaluation
Narcissus is a meta-circular interpreter for JavaScript, written by Mozilla, that is used for prototyping improvements to the language. Tom Austin and Cormac Flanagan used Narcissus as a base to implement their dynamic information flow analysis, called faceted evaluation [2], in JavaScript. This instrumentation was done on top of Narcissus in a straightforward way, that is, favoring ease of implementation. As a result, getting all the modifications made by the instrumentation requires extracting the difference between two well-chosen commits. If we extract such a diff, we can see that the instrumentation touches three files out of seven, and nearly all the changes are in the file concerned with executing the parsed code. Looking more closely at these changes, two patterns are apparent:

1. The addition of a \( \text{pc} \) parameter (Program Counter) at nine distinct locations. In the faceted evaluation strategy, sensible values are wrapped into pairs: one public value and one visible only to a set of principals. The program counter records the branches taken by the program, allowing implicit flows to be tracked. As the program counter is part of the state used by the faceted evaluation strategy, it has to be present in the execution context of Narcissus. Here are two occurrences of this pattern (‘<’ are lines removed, ‘>’ are lines added):

We think they apply broadly to other dynamic analyses, other
The instrumentation should be easily activated dynamically. Multiple, non-interfering instrumentations can be
Pluggability
Intercession
The instrumentation may extend, partially alter or entirely replace the behavior of code present in the base interpreter.
Modularity
The instrumentation should be defined as a module; all its code must be part of this module. In particular, no instrumentation code should appear in the base interpreter source.
Intercession
The instrumentation may extend, partially alter or entirely replace the behavior of code present in the base interpreter.
Local state
The instrumentation may refer to state that concerns only the dynamic analysis it implements.
Pluggability
The instrumentation should be easily activated dynamically. Multiple, non-interfering instrumentations can be activated together without extra effort.

2.2 The Instrumentation Problem
Based on the observations made in the case study, we define the requirements for a modular instrumentation of interpreters for dynamic analyses. They are:

Modularity
The instrumentation should be defined as a module; all its code must be part of this module. In particular, no instrumentation code should appear in the base interpreter source.
Intercession
The instrumentation may extend, partially alter or entirely replace the behavior of code present in the base interpreter.
Local state
The instrumentation may refer to state that concerns only the dynamic analysis it implements.
Pluggability
The instrumentation should be easily activated dynamically. Multiple, non-interfering instrumentations can be activated together without extra effort.

These requirements are not specific to the faceted evaluation analysis, to the Narcissus interpreter or to the JavaScript language. We think they apply broadly to other dynamic analyses, other interpreters, and any language where modules can be defined.

2.3 Relation to the Expression Problem
Extending an interpreter can be seen as an instance of the expression problem coined by Wadler [14]. In fact, the instrumentation problem can be thought of as the expression problem with additional requirements. Wadler defines the expression problem as:

The Expression Problem is a new name for an old problem. The goal is to define a datatype by cases, where one can add new cases to the datatype and new functions over the datatype, without recompiling existing code, and while retaining static type safety (e.g., no casts).

The lack of recompilation of existing code can be roughly translated as the modularity requirement of the instrumentation problem. Both aim for a clean separation between base code and extension code. The addition of cases and functions can be absorbed into the intercession requirement. Static type safety is not applicable in our case, since JavaScript is dynamically typed. More importantly, the expression problem in the context of interpreters is concerned only with extending interpreters via additional operations (like a pretty printer), rather than modifying existing operations (the evaluation function in particular).

2.4 A Minimal Example
While the instrumentation of faceted evaluation on Narcissus is far from simple, the changes it expresses are all covered by the requirements of the instrumentation problem. To focus on the underlying language mechanisms needed for instrumentation, we exhibit a minimal example. To that effect, we follow Wadler and Odersky [10] and define a language of arithmetic expressions. We show that we can add new data variants and new operations straightforwardly in JavaScript, thereby solving the expression problem. We then illustrate how we can modify existing operations and pass state, thus satisfying the requirements of the instrumentation problem. Along the way, we mention how the examples relate back to faceted evaluation.

3. Building an Interpreter with Modules
In all the following code examples, we use a subset of the EcmaScript 6 standard of JavaScript as implemented in Mozilla Firefox 37a, for the shorthand notations reduce syntax noise. All the examples can be translated back to standard EcmaScript 5 supported by all major JavaScript implementations. In the examples, the result of the preceding expression is indicated by the comment syntax ‘//’.

3.1 Solving the Expression Problem
The simplest language has terms only for numbers, and an operation to evaluate a number to a numeric value.

```javascript
< function ExecutionContext(type, version) { >
 function ExecutionContext(type, pc, version) { }
< var x = new ExecutionContext(GLOBAL_CODE, Narcissus.options.version); >
 var x = new ExecutionContext(GLOBAL_CODE, new ProgramCounter(), Narcissus.options.version);

2. The addition of tests for FacetedValue objects. New cases are created to handle faceted values, and tests are needed to distinguish these values from other runtime values, such as references.

```
```javascript
The fact that these patterns have multiple occurrences hints at a possible factorization.
Another observation is that all the changes made for the instrumentation are scattered throughout the code. Without the diff or prior knowledge of the base interpreter, it is impossible to tell the instrumentation parts from the base interpreter. The dynamic analysis and the interpreter are effectively tangled together.

The Expression Problem is a new name for an old problem.

1 var num = {
2 new(n) { return {__proto__: this, n} },
3 eval() { return this.n }
4 }
5
6 var e1 = num.new(3)
7 e1.eval() //: 3
```

We adopt an object-oriented decomposition where data variants are objects, and operations on the datatype are methods of these objects. The following code in JavaScript implements this datatype using objects and prototypical inheritance. The num object has two functions: new for returning an object that has this as prototype and holds the numeric value n, and eval for returning the numeric value held by such objects. On line 6, e1 is an instance of the num term, and we can call eval on it to get its value.

All the code examples from this paper can be found online at https://gist.github.com/fmdkdd/6f0faa7d105dbd8c514
3.1.1 Adding a Variant

The ‘plus’ data variant takes two terms and evaluates to the sum of these terms.

<table>
<thead>
<tr>
<th>Num</th>
<th>Plus</th>
</tr>
</thead>
<tbody>
<tr>
<td>eval</td>
<td>eval</td>
</tr>
</tbody>
</table>

In our object-oriented implementation, a new data variant is easily added by creating a new object `plus` with an `eval` function. The evaluation of `plus` recursively calls `eval` on the left and right sub-expressions.

```javascript
var plus = {
  new(l, r) { return {__proto__: this, l, r} },
  eval() { return this.l.eval() + this.r.eval() }
}
var e2 = plus.new(num.new(1), num.new(2))
e2.eval() // 3
```

In the instrumentation of faceted evaluation, `instanceof` tests were added to distinguish a `FacetedValue` from a `Reference`. Both objects are data variants of the datatype of runtime values manipulated by the interpreter. By dispatching on these data variants instead of using `instanceof` tests, we can write the `FacetedValue` like the `plus` object, and extend the interpreter incrementally.

3.1.2 Adding an Operation

We now wish to add another operation to print expressions instead of evaluating them.

```javascript
var show = function(base) {
  var num = {...proto__: base.num,
    show() { return this.n.toString() }}
var plus = {...proto__: base.plus,
  show() { return this.l.show() + `+` + this.r.show() })
var s = show({num, plus});
e2.show //: undefined
plus_new(s.num.new(1), s.num.new(2)).show() //: "1+2"
s.plus_new(s.num.new(1), s.num.new(2)).eval() //: 3
```

As lines 23 and 25 attest, existing expression objects are not affected by the definition or activation of `show`. But on line 26, a new expression is created using the derivatives of `plus` and `num`, and supports `show`.

On line 27, we see that it is possible to mix expression objects that support show, and those that do not: the second is created from `s.num`, while the first one is not. When calling `eval`, an operation that both have in common, this is not an issue. But since JavaScript is not statically typed, nothing would prevent us to call `show` on an expression of mixed languages, and raise an error only at runtime.

It is possible to prevent the writing of expressions of mixed languages by shadowing the bindings for the original `num` and `plus`. By using the `with` construct on an object, all the properties of this object are in scope of the code delimited by `with`. On line 29, `plus` and `num` refer to the derivatives created by the ‘show’ module. The original `num` and `plus` are shadowed by these bindings; it is impossible to refer to them inside the `with` construct. In addition, the extension of the scope does not persist when exiting `with`, so the `num` and `plus` bindings only exist inside, restricting the activation of the `show` module.

```javascript
with((show({num, plus}))) {
  plus_new(num.new(1), num.new(2)).show() //: "1+2"
}
```

The `with` construct does not allow us to select which names to import from the `show` module; every property of the object returned by `show` on line 28 is put into scope. If we would rather selectively import a subset of names from the ‘show’ module, then we can use an immediately-invoked function expression (IIFE) and specify which names we want out as arguments to this function. For instance, here we only import `num`.

```javascript
(function(num) {
  num.new(1).show() //: "1"
})(show({num, plus}))
```

The last example shows that the `with` construct is not necessary to our scheme, since an IIFE is more expressive. Nevertheless, it provides an interesting use-case for this controversial construct, and since examples using `with` are easier to read, we will stick to it throughout the rest of the paper.

3.1.3 Summary

We have seen how we could:

- add variants to a datatype by defining new objects (3.1.1),
- add operations that act on all the existing variants by extending the prototype of objects (3.1.2),
- without modifying the existing code.

Thereby, we have solved the expression problem as it applies to a dynamic language such as JavaScript.
Furthermore, we have seen that we could add operations to the
datatype in two ways: by directly modifying the prototype and
affecting all the existing and future objects of the extended data
variants, or by creating derivatives of the data variants. With the latter
it is possible to mix data variants from separate languages in unsafe
ways. We have seen how we can restrict the use of only one language
at a time by shadowing and creating new scopes with function
expressions and the `with` construct; this dynamic activation satisfies
the pluggability requirement of the instrumentation problem.

Interestingly, in all the examples we have seen so far, we used
only a fraction of the features provided by the language, without
resorting to extra libraries or framework. We develop this point
further in section 4.

3.2 Solving the Instrumentation Problem

We have seen how to extend the interpreter for the language of
arithmetic expressions by adding variants and operations increment-
ally as modules – without modifying the extended code. Now we
tackle the remaining requirements of the instrumentation problem:
intercession and local state.

3.2.1 Modifying Operations

Let us imagine that we want to modify the ‘eval’ operation on the
data variant ‘num’ to return twice the value it holds.

<table>
<thead>
<tr>
<th>Num</th>
<th>Plus</th>
<th>double.Num</th>
</tr>
</thead>
<tbody>
<tr>
<td>eval</td>
<td>Plus</td>
<td>eval</td>
</tr>
</tbody>
</table>

This is straightforwardly achieved in JavaScript, since functions
held by objects are like other properties: replaceable by assignment.
We used the same mechanism in the addition of `show`. Here we erase
the previous definition of `eval` and replace it with the new one.

```javascript
num.eval = function() { return this.n * 2 }

num.new(1).eval() // 2
plus.new(num.new(1), num.new(2)).eval() // 6
```

Now we change the specification of our modification to ‘num’: it
should return twice the value of the result of the original ‘eval’
operation on ‘num’. The solution requires the ability to call the
previous version of `num.eval`, so we use a function expression to
save this version in the closure of the new `eval`. Then we need to
call this previous version with the current object as receiver, which
is done by invoking the `call` function with the argument `this`.

```javascript
(function(previousEval){
    num.eval = function() { return previousEval.call(this) * 2 }
}(num.eval))

num.new(1).eval() // 2
plus.new(num.new(1), num.new(2)).eval() // 6
```

<table>
<thead>
<tr>
<th>Num</th>
<th>Plus</th>
<th>double.Num</th>
</tr>
</thead>
<tbody>
<tr>
<td>eval</td>
<td>eval</td>
<td>eval</td>
</tr>
</tbody>
</table>

Modifying the original `num` object is destructive: the original
`eval` is lost, only accessible by the closure created by the IIFE. In
this instance, we can create a new data variant that refers to the
original `eval` to avoid duplication. To that effect, we create a function
parameterized by the original `num` object to serve as a module.

```javascript
var double = function(base) {
    var num = {__proto__: base.num,
        eval() { return base.num.eval.call(this) * 2 }}
}
```

On line 33 we define the new `eval` as a wrapper around the previ-
ous functionality. This enables us to compose extensions by passing
a modified base as parameter to `double`. The following example
illustrate how we can combine such extensions by cascading calls
to `with`.

```javascript
with (double(num)) {
    with (double(num)) {
        plus.new(num.new(1), num.new(2)).eval() // 24
    }
}
```

3.2.2 Passing State to Operations

In the instrumentation of Narcissus for faceted evaluation, the
“program counter” is an object that must be passed down to recursive
calls of the interpreter, representing state needed by the dynamic
analysis. To mimic this behavior, here we add a program counter
to the ‘eval’ operation that is incremented each time a data variant
calls `eval`.

<table>
<thead>
<tr>
<th>Num</th>
<th>Plus</th>
<th>state.Num</th>
<th>state.Plus</th>
</tr>
</thead>
<tbody>
<tr>
<td>eval</td>
<td>eval</td>
<td>eval</td>
<td>eval</td>
</tr>
</tbody>
</table>

The `state` function creates derivatives of the `num` and `plus` data
variants that increments the local program counter `pc` when calling
`eval`. The value of the program counter is exposed via the exported
`getPC` function.

```javascript
var state = function(base, pc = 0) {
    var num = {__proto__: base.num,
        eval() { pc++; return base.num.eval.call(this) }}
    var plus = {__proto__: base.plus,
        eval() { pc++; return base.plus.eval.call(this) }}
    var getPC = () => pc
    return [num, plus, getPC] }

with (state(num, plus)) {
    getPC() //: 0
    plus.new(num.new(1), num.new(2)).eval() //: 3
    getPC() //: 3
}
```

On line 52 we see that the program counter is zero, and is later
incremented to 3 on line 54 after the evaluation of the expression on
line 53.

On lines 42 and 45, the new `eval` functions are defined as
wrappers around `base.num.eval` and `base.plus.eval`, which allows
for composition of multiple extensions. For instance, we can add
the `double` modification of the previous subsection.

```javascript
with (double(num)) {
    with (state(num, plus)) {
        getPC() //: 0
        plus.new(num.new(1), num.new(2)).eval() //: 6
        getPC() //: 3
    }
}
```

4 2015/7/29
We have seen how to extend the solution to the expression problem

3.2.3 Summary

Pluggability

module objects.

instanceof

The object modules, satisfy all the criteria of the instrumentation prob-

sion and interces-

of subsection 3.1 to tackle the additional requirements of

modules is irrelevant.

data variant, hence they commute.

a side effect that does not interfere with the doubling of the 'num'

semantics of each operation. In this instance, the program counter is

is irrelevant. This is not always the case, as it depends on the

both the double modification and state modification are in effect.

Lines 63-68 attest that the order of activation of these modifications

is irrelevant. This is not always the case, as it depends on the

semantics of each operation. In this instance, the program counter is

a side effect that does not interfere with the doubling of the ‘num’

data variant, hence they commute.

Finally, we can also add the show module without modifying

any of the existing definitions. Here again the order of activation of

modules is irrelevant.

with (state([num, plus])) {
    getPC() //: 0
    plus.new(num.new(1), num.new(2)).eval() //: 6
    getPC() //: 3
}
}

On line 59 we see that the expression evaluates to 6, and lines

58 and 60 indicate that only three expressions were evaluated. Thus

both the double modification and state modification are in effect.

Lines 63-68 attest that the order of activation of these modifications

is irrelevant. This is not always the case, as it depends on the

semantics of each operation. In this instance, the program counter is

a side effect that does not interfere with the doubling of the ‘num’

data variant, hence they commute.

Finally, we can also add the show module without modifying

any of the existing definitions. Here again the order of activation of

modules is irrelevant.

with (state([num, plus])) {
    with (double([num])) {
        with (show([num, plus])) {
            getPC() //: 0
            var n = plus.new(num.new(1), num.new(2))
            n.eval() //: 6
            getPC() //: 3
            n.show() //: "1+2"
        }
    }
}

3.2.3 Summary

We have seen how to extend the solution to the expression problem

of subsection 3.1 to tackle the additional requirements of interces-

sion and local state. The solution we exposed, based on simple

object modules, satisfy all the criteria of the instrumentation prob-

lem we defined in subsection 2.2:

Modularity

The show, double and state extensions were all ex-

pressed as modules, without any need to alter or duplicate the

code of the parts they augment. The plus data variant, expressed

simply as an object, as well as the first definition of num, could

have been built as a simple modules with no import.

Intercession

The plus and show examples were two cases of ex-

tensions to the base language. The double and state examples

illustrated replacement of the existing eval operation. In the

double case, we first saw a complete replacement of the eval

functionality. In all cases, the extension or alteration of the base

language was done incrementally – without touching the ex-

tended code.

Local state

The state example served to illustrate how we could

add state that is local to an extension, without having to pass

it as an argument to the modified eval operations, and without

altering the base code.

Pluggability

The show, double and state examples were all easily

activated by the with construct, or an immediately-invoked

function expression. Each time, the extract same line of code was

used to add 1 and 2, and each time the effect varied depending

on the active modules. We have shown how these extensions

could all be activated at the same time, without modification

needed to the base code or to the client code inside with.

In addition, we have illustrated how to solve two patterns

we highlighted from the instrumentation of Narcissus for faceted

evaluation in subsection 2.1: (1) the addition of instanceof tests can

be transformed into dispatching on new data variants, and (2) the

addition of local state (the program counter) can be encapsulated

in module objects.

4. Discussion and Related Work

A combination of simple language features

To solve the expression problem and instrumentation problem, we used only features

provided by the JavaScript language; we did not extend JavaScript

with a class system, a reflection framework or even a module system.

Several key features of JavaScript made this possible.

The flexibility of the object system is an advantage here. Objects

are essentially dictionaries with the ability to inherit keys from a

prototype link. Since objects are open, it is easy to add or modify

methods. Objects can also be created at any point at runtime, and

this allowed us to define derivatives that delegate to base objects

with a prototype link. To distinguish the derivatives from the original

objects, we only needed a simple object as a dictionary to serve as

namespace.

Our module objects are parameterized by the objects they extend

just because they are defined as functions. Functions in JavaScript

are first-class values, and can be used as expressions for defining ob-

ject methods, or the module objects we presented. When activating

modules, the shadowing of names in the outer scope was a desirable

feature, provided by the with statement and functions expressions.

In JavaScript, there are no methods strictly speaking: the value of

the receiver, accessible inside a function via the this keyword, is
determined only when a function is called, not when it is defined.

We can even specify its value to be an arbitrary object using the call

construct. We used this feature whenever we needed to invoke an

overridden behavior on the current object. Dynamic typing is also a benefit in this case. The absence of

a static type system allows us to compose objects in a way that

would be difficult to type in languages like Java, Haskell or Scala.

Odersky and Zenger [10] show how they can use Scala to solve

the expression problem and retain static type safety, but they do

not consider the modification of existing operations. While it is

one thing to safely extend interpreters, it is quite another to safely

change a function signature at some point in runtime. In fact, the

requirement of intercession in the instrumentation problem seems at

odds with static type safety; but that remains to be shown.

The design philosophy behind our scheme is inspired by Findler

and Flatt [5]. They exhibit a solution to the expression problem in

MzScheme using mixins and a form of modules called units. Units

and mixins follow a single design principle: “specify connections

between modules or classes separately from their definitions”. In our

case, this is translated as adding indirection by names (arguments

and namespaces). This loose coupling can be found in our module

objects: they are parameterized by the objects they extend, rather

than being bound to it prematurely. Late binding is why we can

compose module objects in any order in the later examples.

Aspect-Oriented Programming (AOP)

The instrumentation of Narcissus for faceted evaluation is a crosscutting concern. AOP

is well-suited for expressing such concerns in a modular way,

using predicates (pointcuts) to target specific parts of the code

(joinpoints) where instrumentation should take over. For instance,

FlowR [11] is a library that provides information flow control for

the Ruby programming language. The authors used an AOP toolkit

to write FlowR and successfully lowered the complexity of their

implementation.

AspectScript [13] is an implementation of AOP for JavaScript

with advanced scoping strategies, and the ability to capture join-

points inside function bodies. As it wraps all the client code, not

necessarily code that is targeted by pointcuts, AspectScript multi-

plies the execution time by a factor of 5 to 13. Achenbach and

Ostermann define a meta-aspect protocol for developing dynamic

program analyses [1], and provide a Ruby implementation.

While we have no doubts that the instrumentation problem we

consider can be solved using AOP, the scheme we present here shows
that the flexibility of a few JavaScript constructs are sufficient. By leveraging the right parts of the language, we are able to define extensions as modules without having to refer to a sophisticated model.

Context-Oriented Programming (COP) Our way of activating modules via with is reminiscent of the activation of layers in COP systems. ContextJS [7] is a COP extension to JavaScript, enabling sideline composition of layers on a base class. In ContextJS, a layer refines a class with partial methods that override the behavior of methods of the base class; extra methods can be added as well. Layers can then be activated and deactivated dynamically, on all or specific instances of the class, for a delimited extent or globally. Layers can stack, and partial methods can refer to the behavior of upper layers indirectly by using the proceed special argument, allowing the gracefully composition of layers similar to our final examples in 3.2.2.

We can view our module objects as rudimentary layers. Our composition of modules could be improved syntactically by using a proceed keyword. The layer activation possibilities of ContextJS are more expressive, but they go beyond the needs of the instrumentation problem we defined. Our module objects are parameterized by the base object they extend, and as such are not tied to a particular class. The main difference is, again, that the simplicity of our scheme lies in combining the raw parts of the JavaScript language, and does not require external definitions.

Other approaches to modular instrumentation Marek et al. propose DiSL [8], a Java framework for writing dynamic program analyses using the AOP model of pointcuts and joinpoints, focusing on ease of instrumentation and efficiency of weaving. Polyglot [9] is an extensible compiler front-end for Java used notably for implementing the Jif language [4], which extends Java with information flow control functionality. These approaches share our goals of extending languages with dynamic analyses while preserving modular reasoning. However, they target the Java language, which apart from inspiring the name “JavaScript”, bear little resemblance to the language and platform we address.

Modular instrumentation of Narcissus The remaining question is: how does the solution sketched here apply to the instrumentation of Narcissus for the faceted evaluation analysis? Our examples apply to the interpreter for arithmetic expressions that was ostensibly developed for context-oriented decomposition of data variants and operations, instead favoring switch-cases and instanceof tests. Refactoring Narcissus is clearly needed as a first step to allow extensions to be written. The refactoring effort should be moderate, as the interpreter part of Narcissus is only 1500 lines, and the rest of the program is already split into modules. Once Narcissus is in a form more amenable to extension, we expect the scheme presented here to allow the modular expression of faceted evaluation and other dynamic information flow analyses.

Furthermore, since the instrumentation problem we defined is not particularly tied to Narcissus, we expect our solution to be applicable to other interpreters written in JavaScript.

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References


