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Predictive control algorithm for spacecraft rendezvous hovering phases

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Abstract—The present article proposes a predictive control law for orbital rendezvous hovering phases. The control consists of impulsive thrust corrections that maintain the relative motion between spacecraft bounded. The proposed predictive algorithm accounts for thrusters saturation and a propellant budget on top of the state constraints. It is shown that computing an admissible control is equivalent to determine a point in the intersection between two closed convex sets. To resolve this problem, an alternating projections algorithm is used. The efficiency of the proposed predictive control algorithm is bench tested on a nonlinear simulator.

Keywords—predictive control, rendezvous, periodic relative motion, constrained trajectories, non-negative polynomials, alternating projections algorithm

I. INTRODUCTION

In the future missions of on-orbit inspection and on-orbit servicing, the ability to maintain a relative periodic motion is a crucial issue for any rendezvous spacecraft. The purpose of this research is to develop a predictive control law that enables the control of periodic relative orbits that are included in a given linear region for sake of safety.

Constrained relative trajectories have been studied under the term of "hovering". Hovering has been defined in [1] as follow

To develop a control strategy to place a deputy satellite inside a specific lobe defined in the chief body-fixed frame and keep it there in a fuel-optimal manner.

In most of works on the subject, as [1]–[7], the hovering strategy is based on two different strategies. First, the "teardrop" strategy consists in computing an impulsive thrust such that the initial and final relative positions on the bounds of the tolerance subset are equal for a given time of flight [1], [8]. Second, the "pogo" strategy, where the impulsive control is computed each time the vehicle hit the bounds of the lobe to maintain it as long as possible while minimizing the fuel consumption [1], [4], [7].

However, some contributions have recently taken advantage of the naturally periodic relative orbits to design hovering control laws [9]–[11]. In [9], a methodology to compute the coordinates of periodic relative orbits included in a polytopic region is exposed. This method has been extended to the non periodic relative trajectories in [11].

A control law proposed in [10] consists of two impulsive maneuvers that stabilize the spacecraft with respect to a given constrained periodic relative orbit previously computed. Even though this control law has been proven stable for any eccentric target orbit and requires very few computational effort, its optimality in terms of fuel consumption is not guaranteed.

The present article proposes a different strategy to hover inside a given polytopic subset of the relative state space. The purpose is to steer the vehicle to the set of the periodic relative orbits enclosed in a hovering zone.

To achieve such a goal, a predictive methodology is applied to provide the impulsive control. At each call, the controller computes, if it exists, the coordinates of a relative periodic orbit included in the hovering zone that go through the current position. The impulsive control is then deduced from these coordinates. Moreover, during the control computation, spacecraft abilities and propellant resource are also taken into account.

Although the present work extends the results from the conference article [12], it differs in several points. In the present paper, the chaser is steered to a periodic and admissible orbit not defined a priori, contrary to what is done in [12]. Consequently, the control consists of one impulse instead of two for the techniques from [12]. Finally, the proposed methodology permits to consider in the problem saturation and propellant budget constraints.

As it is exposed thereafter, the set of the orbits coordinates admissible with respect to the above-mentioned constraints is described as the intersection between a semi-algebraic set and a hyperplane. Obtaining the control mainly consists in computing one point of this intersection set. This is achieved using an alternating projections algorithm. Note that this algorithm by means of its properties make the use of a complete optimization procedure unnecessary as explained in this paper.

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The paper is structured as follows. Section II gives a brief description of the relative motion and recall the trajectories parametrization used in the sequel. Section III describes, in terms of the relative trajectories coordinates, the constraints on the relative orbits and the controls. Section IV presents the alternating projection algorithm that permits to compute the coordinates of an admissible relative trajectory and the control. Finally, section VI proposes numerical examples.

II. RELATIVE MOTION MODELING

The Fig. 1 presents the frames used to model the relative motion between the leader \( S_l \) and the follower \( S_f \) spacecraft.
The ECI\(^1\) and the LVLH\(^2\) frames are respectively given by \(\{O, \hat{I}, \hat{J}, \hat{K}\}\) and \(\{S_l, \hat{x}, \hat{y}, \hat{z}\}\). The true anomaly \(\nu\) (the angle between the direction of perigee and the leader’s position) is also shown. Considering that \(R = ||\hat{O}S_l|| \gg ||S_l\hat{S}'_f||\) and assuming Keplerian orbits, the relative dynamics can be given by the linearized and simplified Tschauner-Hempele equations [13] (time is replaced by true anomaly as independent variable):
\[
\begin{align*}
\ddot{x}' &= 2\dot{\chi}' + \dot{u}_x \\
\ddot{y}' &= -\dot{\chi}' + \dot{u}_y \\
\ddot{z}' &= \frac{3}{1 + e \cos \nu} \dot{\chi}' - 2\dot{\nu}' + \dot{u}_z 
\end{align*}
\]  
where \(e\) is the leader orbit eccentricity. Defining the vector \(\dot{X}(\nu) = [\dot{x}(\nu) \; \dot{y}(\nu) \; \dot{z}(\nu)]^T\), a state-space representation for the non-autonomous system (1) is:
\[
\dot{X}(\nu) = A(\nu)X(\nu) + B\ddot{u}
\]  
Note that the state expressed in time and expressed in true anomaly are linked by means of the following similar transformation:
\[
X(\nu) = \left[\begin{array}{c}
(1 + e \cos \nu)I_3 \\
-\dot{e} \sin \nu I_3
\end{array}\right] \left[\begin{array}{c}
0_3 \\
\dot{e}
\end{array}\right] X(t)
\]  
Using the transition matrix, \(\Phi(\nu, \nu_0)\), proposed by Yamanka and Ankersen in [14] to propagate an initial state \(X(\nu_0)\) and considering an unforced motion:
\[
\dot{X}(\nu) = \Phi(\nu, \nu_0)X(\nu_0), \quad \nu \geq \nu_0
\]  
Expanding eq. (4), the following relative trajectories description can be deduced:
\[
\begin{align*}
\dot{x}(\nu) &= (2 + e c_\nu)(d_1 s_\nu - d_2 c_\nu) + d_3 + 3d_0 J'_{me}(1 + e c_\nu)^2 \\
\dot{y}(\nu) &= d_4 c_\nu + d_5 s_\nu \\
\dot{z}(\nu) &= (1 + e c_\nu)(d_2 s_\nu + d_1 c_\nu - 3e d_0 J'_{me} s_\nu) + 2d_0
\end{align*}
\]  
where \(c_\nu = \cos \nu, \; s_\nu = \sin \nu\), \(J'_{me}\) is given by:
\[
J'_{me} = \int_{0}^{\nu} \left(\frac{1}{1 + e \cos \tau}\right)^2 d\tau
\]  
and \(D(\nu) = [d_0(\nu) \; d_1(\nu) \; d_2(\nu) \; d_3(\nu) \; d_4(\nu) \; d_5(\nu)]^T\), the parameter vector, is expressed as a linear function of \(\dot{X}(\nu)\):
\[
D(\nu) = C(\nu)\dot{X}(\nu)
\]  
Using the variable change (12) and the rational expressions (13) in (11), we can obtain the following polynomial inequalities:

\[
\begin{align*}
P_x(w) &\leq ((-1-e)w^4 + 2w^2 + 1 + e)x_{\max} \\
P_y(w) &\leq y_{\max} \\
P_z(w) &\leq ((-1-e)w^4 + 2w^2 + 1 + e)z_{\max} \\
P_x(w) &\geq ((-1-e)w^4 + 2w^2 + 1 + e)x_{\min}, \quad \forall w \geq w_0 \\
P_y(w) &\geq y_{\min} \\
P_z(w) &\geq ((-1-e)w^4 + 2w^2 + 1 + e)z_{\min}
\end{align*}
\]  

that can be reformulated as:

\[
\Gamma_i(w, D) \geq 0, \quad \forall w \geq w_0, \quad i = 1..6 \quad (16)
\]

The non-negativity conditions in (16) result in an infinite number of constraints to be respected. Deaconu presents in [9] a way to transform this infinite number of constraints in a finite number of convex constraints using the following theorem presented by Nesterov in [15]:

**Theorem 1:** Let \( P(w) \) be a polynomial of degree \( 2n \) on \( \mathbb{R} \) and \( p \in \mathbb{R}^{2n+1} \) the vector containing its coefficients. Then:

\[
\begin{align*}
P(w) &\geq 0, \quad \forall w \in \mathbb{R} \iff \exists Y \in S^+_{n+1}|p = \Lambda^*(Y) \\
p_l = \Lambda_l^*(Y) &= \text{trace}(Y H_{n,l}), \quad l = 1..2n+1
\end{align*}
\]

where \( S^+_{n+1} \) is the cone of symmetric positive semidefinite matrices of size \( n+1 \) and \( H_{n,l} \) are the Henkel matrices:

\[
H_{n,l}(j, k) = \begin{cases} 
1, & \text{if } j + k = l + 1 \\
0, & \text{otherwise}
\end{cases} \quad (18)
\]

This is equivalent to affirm that the cone of non-negative univariate polynomials is the image of the cone of positive semi-definite matrices by the linear operator \( \Lambda^*(\cdot) \).

Let \( \gamma_i \) be the vectors containing the coefficients of the polynomials \( \Gamma_i(w) \). Using (17), we can transform conditions in (16) into:

\[
\begin{align*}
H_{i, \tilde{X}_{\text{pos}}}(v, D) &\leq \tilde{K}_i(v) \\
\Gamma_i(w, D) &\geq 0, \quad \forall w \geq w_0, \quad i = 1..6 \\
\exists Y_i \in S_+ &\text{ such that } \gamma_i = \Lambda^*(Y_i)
\end{align*}
\]

In the end of this section, constraints on the control are given: the thrusters saturation and the propellant budget. Considering saturation constraints is one of the main advantages of the predictive control law. Moreover, the propellant is a rare resource for spacecraft systems since they can not be refueled. Thus, propellant budget constraints is also proposed to take into account the scarceness of the propellant.

**C. Thrusters saturation**

Let \( \Delta V = [\Delta V_x \Delta V_y \Delta V_z]^T \) be the vector representing the impulsive velocity correction executed by the follower. Considering that the spacecraft have a pair of equivalent thrusters symmetrically and oppositely disposed on each axis and assuming that the saturation limit for each thruster is \( \Delta V_{\text{max}} > 0 \), the saturation constraint can be written as:

\[
|\Delta V_s| \leq \Delta V_{\text{max}}, \quad s = x, y, z \quad (20)
\]

These inequalities can be transformed in a matrix positive semi-definite constraint:

\[
|\Delta V_s| \leq \Delta V_{\text{max}} \iff \begin{bmatrix} \Delta V_{\text{max}} & \Delta V_s \\ \Delta V_s & \Delta V_{\text{max}} \end{bmatrix} \geq 0, \quad s = x, y, z \quad (21)
\]

**D. Propellant budget**

An maximum budget of propellant \( \sigma \) is available for each impulse executed. This constraint is expressed as:

\[
||\Delta V||_1 = |\Delta V_x| + |\Delta V_y| + |\Delta V_z| \leq \sigma \quad (22)
\]

To check condition (22), we must verify if:

\[
\begin{align*}
Z_x, Z_y, Z_z &\geq 0 \\
W^-_x, W^+_x, W^-_y, W^+_y, W^-_z, W^+_z &\geq 0 \quad \text{s.t.} \quad W^- = Z_s - \Delta V_s \\
W^+_x = Z_s + \Delta V_s, \quad W^+_y = Z_s + \Delta V_s, \quad W^+_z = Z_s + \Delta V_z \\
W_0 &\geq 0 \\
\sigma - \sum Z_s = W_0
\end{align*}
\]  

**IV. Computing the control law**

Accordingly to the model predictive control scheme, the control is computed periodically in order to bring the system on a periodic trajectory inside a given polytope respect to actuator saturation and budget constraints. Using the previous equations, links between the variables of the problem and its restrictions can be stated. The Fig. 2 presents these links, each ellipse representing a constraint.

![Fig. 2: Variables and its constraints](image)

Given a set of constraints \( \{ H, K, \Delta V_{\text{max}}, \sigma \} \) and an initial relative state \( X_0 \), we want to determine an impulse \( \Delta V \) respecting both saturation and budget constraints such that the state after the impulse, \( X^+ = X_0 + \Delta V \), equivalent to the state \( D^+ \) (by means of (7)) whose first component \( d_0 \) is zero to stop the periodicity and that generates (by means of (14), (15) and (16)) six non-negative polynomials to ensure the polytopic inclusion of the orbits.

**A. Closed convex sets intersection problem**

The constraints can be summarized as follows:

\[
\begin{align*}
\text{Periodicity:} & \quad d_0 = 0 \\
\text{Polytope constraints:} \\
\exists Y_i \geq 0 & \quad \text{s.t.} \quad \gamma_i = \Lambda^*(Y_i), \quad i = 1..6 \\
\text{Thrusters saturation:} \\
\exists M_s \geq 0 & \quad \text{s.t.} \quad \begin{bmatrix} \Delta V_{\text{max}} & \Delta V_s \\ \Delta V_s & \Delta V_{\text{max}} \end{bmatrix} = M_s, \quad s = x, y, z \\
\text{Propellant budget:} \\
Z_s &\geq 0 \\
W^-_s &= Z_s - \Delta V_s \\
W^+_s &= Z_s + \Delta V_s, \quad s = x, y, z \\
W_0 &\geq 0 \\
\sigma - \sum Z_s = W_0
\end{align*}
\]  

(24)
We remark that these constraints are convex, since they are mathematically represented by matrices positive semi-definiteness (on the left) and equalities (on the right).

Let \( Q \) be the symmetric block diagonal matrix given by:
\[
Q = diag(M_x, M_y, M_z, Z_x, Z_y, Z_z, W^+_x, W^+_y, W^+_z, W^-_x, W^-_y, W^-_z, W_0, Y_1, \ldots, Y_6)
\] (25)

This matrix is sparse by construction. Table I shows that taking symmetry into account we only need 49 variables to completely describe it.

<table>
<thead>
<tr>
<th>( M_x )</th>
<th>#</th>
<th>Size</th>
<th>Total</th>
<th>#Vars</th>
<th>Total</th>
</tr>
</thead>
<tbody>
<tr>
<td>( Z_x, W^+_x, W^-_x, W_0 )</td>
<td>10</td>
<td>1 x 1</td>
<td>10 x 10</td>
<td>1</td>
<td>10</td>
</tr>
<tr>
<td>( Y_{(1,3,4,6)} )</td>
<td>4</td>
<td>3 x 3</td>
<td>12 x 12</td>
<td>6</td>
<td>24</td>
</tr>
<tr>
<td>( Y_{(2,5)} )</td>
<td>2</td>
<td>2 x 2</td>
<td>4 x 4</td>
<td>3</td>
<td>6</td>
</tr>
<tr>
<td><strong>Total</strong></td>
<td><strong>32 x 32</strong></td>
<td><strong>49</strong></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

We can use \( Q \) to reformulate the positive semi-definiteness conditions and the equations linking the variables in (24) as follows:
\[
\begin{align*}
\{ & \text{tr}(A_i Q) = b_i, \quad A_i \in S^{32}, \quad b_i \in \mathbb{R}, \quad i = 1, m \\
& Q \succeq 0
\end{align*}
\] (26)

where \( S^{32} \) is the set of \( 32 \times 32 \) symmetric matrices and \( m \) is the number of equations.

Equations \( \text{tr}(A_i Q) = b_i \) describe an affine subset of \( S^{32} \) and \( Q \succeq 0 \) represents the symmetric semi-definite matrices cone \( S^{32}_+ \). Therefore, solving the problem is equivalent to find an element \( Q \) belonging to the intersection of these sets.

B. Alternating projections algorithm

For two closed and convex sets, we can apply the alternating projections algorithm proposed by Boyd and Dattorro in [16] to determine the existence of an intersection:

Let \( C \) and \( D \) be two closed convex subsets of \( \mathcal{E} \), let \( \| \cdot \|_E \) be a norm on \( \mathcal{E} \) and let \( P_C \) and \( P_D \) denote the projections on \( C \) and \( D \), respectively. The algorithm consists in:

**Algorithm 1 Alternating Projections**

Choose an arbitrary \( Q_0 \) and compute \( Q^C_1 = P_C(Q_0) \)

for \( k = 0, 1, 2, \ldots \) do

\[
\begin{align*}
Q^D_k &= P_D(Q^C_k) \\
Q^C_{k+1} &= P_C(Q^D_k)
\end{align*}
\]

end for

It has been proven that [17]:

- if \( C \cap D \neq \emptyset \), \( \exists Q_* \in C \cap D \) s.t. \( Q^*_{k \to \infty} \rightharpoonup Q_* \) and \( Q^D_{k \to \infty} \rightharpoonup Q_* \); \\
- if \( C \cap D = \emptyset \), \( \exists Q^C_* \in C, Q^D_* \in D \) s.t. \( Q^C_{k \to \infty} \rightharpoonup Q^C_* \) and \( Q^D_{k \to \infty} \rightharpoonup Q^D_* \) and \( \| Q^C_* - Q^D_* \|_E = \text{dist}(C, D) \), i.e. a pair of points in \( C \) and \( D \) that have minimum distance.

Thereafter, let \( C \) be the affine set described by the equations from system (26), let \( D \) be the symmetric semi-definite matrices cone \( S^{32}_+ \) and let \( \| \cdot \|_E \) be the Frobenius norm. Since these sets are closed and convex, the alternating projections algorithm can be implemented to determine an element belonging to their intersection. In this case, projection \( P_C \) is given by:
\[
P_C(Q) = Q - \sum_{i=1}^m a_i A_i
\] (27)

where \( a_i \) are solution of the following equation:
\[
G a = g
\] (28)

where
\[
G(j,k) = \text{trace}(A_j A_k)
\] (29)

\[
g = [\text{trace}(A_1 Q) - b_1, \ldots, \text{trace}(A_m Q) - b_m]^T
\]

The projection \( P_D \) is given by:
\[
P_D(Q) = U \text{diag}(\max \{0, \lambda_1\}, \ldots, \max \{0, \lambda_{32}\}) U^T
\] (30)

by considering the spectral decomposition is written as
\[
Q = U \text{diag}(\lambda_1, \ldots, \lambda_{32}) U^T
\] (31)

C. Initialization and stopping criteria

As any iterative algorithm, the projection algorithm’s behavior depends on two items: the initial guess and the stopping criteria. The initial guess to provide to the alternating projection algorithm is the matrix \( Q_0 \). In fact, the alternating projection algorithm provides to the initial guess the nearest element in the intersection of the subsets (if exists) [18]. Thus, the choice of \( Q_0 \) can have a great impact the solution \( Q_* \) as can be seen of 3. In this paper, two strategies are tested: the initial guess \( Q_0 \) is either the null matrix or the solution \( Q_* \) of the previous call of the projection algorithm. As it is exposed in section VI, different strategies for choosing the initial iterate will lead to different controlled trajectories.

![Fig. 3: Alternating projection algorithm](image-url)
distance between the convex subsets so that the intersection is an empty set.

V. Stability by Means of Invariance

The stability around a periodic trajectory is understood in the sense given in [19], which requires the characterization of the evolution of the trajectories that start arbitrarily close to the periodic one. According to [20, theorem 3.5], the stability of periodic system is characterized by the eigenvalues of its monodromy matrix $M = \Phi(\nu + 2\pi, \nu)$. In the case of the linearized relative motion, $M$ has at least two eigenvalues $\mu_i$ such that $|\mu_i| > 1$. This leads to the conclusion that the autonomous system $\dot{X}^i(\nu) = \tilde{A}(\nu)\dot{X}(\nu)$ is naturally unstable. Thus, without control, the trajectories, described by (4), that started arbitrary close to any periodic trajectory will diverge from it.

The proposed predictive control law has for goal to steer the system in the vicinity of the set of admissible periodic orbits. Thus, the stability will be defined in the sequel as the invariance of given subsets described thereafter. Now, two sets are given in order to define the stability of the system controlled by means of the proposed predictive algorithm:

- $S^0 \{H, K\}$, the set of vectors $D(\nu) \in \mathbb{R}^6$ that describes periodic trajectories satisfying the polytope constraints;
- $S^1 \{H, K, \Delta V_{\text{max}}, \sigma\}$, the set of vectors $D(\nu) \in \mathbb{R}^6$ for which it is possible to find one impulse $\Delta V$, respecting the saturation constraints and the propellant budget, such that $D^+(\nu)$, the state right after the impulse, belongs to $S^0 \{H, K\}$.

The set $S^0 \{H, K\}$ naturally invariant because of the periodicity respected by each of its elements.

The set $S^1 \{H, K, \Delta V_{\text{max}}, \sigma\}$ is also a class of equivalence under the developed control law, since it is possible to compute an impulsive correction via the alternating projections algorithm that generates a new state belonging to $S^0 \{H, K\}$. Consequently, the set $S^0$ is invariant and attractive under the developed control law with $S^1$ as convergence set. If the current state does not belong to $S^1$, an ad hoc control law must be used to bring the system in $S^1$. For instance, an optimal impulsive plan computed using techniques from [11] can bring the chaser satellite to the vicinity of the hovering box from any relative state.

The Figure 4 represents the sets $S^0$ and $S^1$ and the stability and instability zones in $\mathbb{R}^6$.

VI. Numerical Examples

To illustrate our methodology, the example of a given mission is proposed in table II.

<table>
<thead>
<tr>
<th>Semi major axis [km]</th>
<th>Eccentricities</th>
</tr>
</thead>
<tbody>
<tr>
<td>20 000</td>
<td>{0.1, 0.2, 0.3, 0.4, 0.5, 0.6}</td>
</tr>
</tbody>
</table>

This mission has been implemented on a simulator based on Gauss nonlinear equations for the relative motion, taking into account the atmospheric drag and Earth's oblateness effects (see [21] for details).

In order to assess the impact of the initialization techniques described in IV-C on the resultant trajectory and the consumption, a warm start initialization, taking the last iteration solution as initial guess, and a cold start initialization, using a constant $Q_0$ matrix for all the iterations, were performed.

Simulations with different eccentricities and time-intervals between control calls are conducted to test the algorithm with both initialization procedures. The aim of these simulations is first to verify the designed control law efficiency in respecting the different constraints while rejecting the unmodeled dynamics, and second to highlight the importance of the algorithm initialization procedure. The results obtained for these simulations are presented in Figure 5. We can remark that, on one hand, the smaller eccentricity, the lower consumption is. On the other hand, the consumption increases with the period between controls. Another conclusion can be drawn from Figure 5: the control algorithm using a cold start technique generates greater consumptions compared to the one using a warm start technique. This can be explained by the fact that cold start leads to computation of parameters of admissible orbit that

![Graph showing mission simulation results](image-url)
We studied the influence of the initial guess choice on the intersection subset of convex sets when it is not empty. The control algorithm provides to any initial guess the nearest point in the admissible orbits. Moreover, budget and saturation constraints are also accounted in the control computation. One of the admissible orbits. Furthermore, budget and saturation constraints are also satisfied. These facts highlight that the admissible orbit computed using a warm start initialization is closer to current position than the one computed from a cold start. In fact, after every call of the control algorithm, the chaser will evolve close to the computed orbit. With the warm start initialization, the next computed orbit will remain in the vicinity of the previously computed orbit. In the contrary, the cold start procedure leads to the computation of a completely different admissible orbit.

On Figure 6, controlled trajectories obtained by means of the nonlinear simulator previously described are exposed. The control algorithm is called every 20 second for this specific simulation. First, the simulated trajectories remains admissible during the 5-orbit period of hovering. However, it can be observed that both initialization procedures lead to different geometries respecting the tolerance box constraints despite the orbital perturbation. The total $\Delta V$ consumption over the 5 orbital periods is of 4.4 mm/s and 18 mm/s for the warm start and cold start control respectively. Note that, the $\Delta V$ budget and saturation constraints are also satisfied. These facts highlight that the admissible orbit computed using a warm start initialization is closer to current position than the one computed from a cold start. In fact, after every call of the control algorithm, the chaser will evolve close to the computed orbit. With the warm start initialization, the next computed orbit will remain in the vicinity of the previously computed orbit. In the contrary, the cold start procedure leads to the computation of a completely different admissible orbit.

The present paper develops a predictive control law for the orbital rendezvous. This control law aims to make the chaser satellite hover in a given tolerance set in order to maintain a way point situation. As the relative motion has a periodic dynamics, the predictive control algorithm take advantage of the existence of periodic orbits for this system to develop a law that maintain the hovering while consuming propellant essentially to reject perturbation, contrary to pogo or tear drop strategies available in the literature. This control law consists in steering periodically the chaser spacecraft into the subset of the admissible orbits. Moreover, budget and saturation condition are also accounted in the control computation. One of the main features of the proposed control law is that it is based on the alternating projection algorithm. This iterative algorithm provides to any initial guess the nearest point in the intersection subset of convex sets when it is not empty. We studied the influence of the initial guess choice on the controlled trajectories and the consumption through nonlinear simulations. Further studies should also focus on the numerical issues like convergence time as function of the stopping criteria tolerance. Another problem of interest is the characterization of the invariant set $S^1$, that could speed up the decision of whether or not to apply the projection algorithm, once it should be possible to determine the algorithm convergence by verifying the belonging of the current state to $S^1$.

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