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The Hopf algebra of Fliess operators and its dual prelie algebra

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ABSTRACT. We study the Hopf algebra $H$ of Fliess operators coming from Control Theory in the one-dimensional case. We prove that it admits a graded, finite-dimensional, connected gradation. Dually, the vector space $\mathbb{R}\langle x_0, x_1 \rangle$ is both a prelie algebra for the prelie product dual to the coproduct of $H$, and an associative, commutative algebra for the shuffle product. These two structures admit a compatibility which makes $\mathbb{R}\langle x_0, x_1 \rangle$ a Com-Prelie algebra. We give a presentation of this object as a Com-Prelie algebra and as a prelie algebra.

KEYWORDS. Fliess operators; prelie algebras; Hopf algebras.

AMS CLASSIFICATION. 16T05, 17B60, 93B25, 05C05.

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Introduction

Right prelie algebras, or shortly prelie algebras [4, 1], are vector spaces with a bilinear product \( \cdot \) satisfying the following axiom:

\[
(x \cdot y) \cdot z - x \cdot (y \cdot z) = (x \cdot z) \cdot y - x \cdot (z \cdot y).
\]

Consequently, the antisymmetrization of \( \cdot \) is a Lie bracket. These objects are also called right-symmetric algebras or Vinberg algebra [12, 17]. If \( A \) is a prelie algebra, the symmetric algebra \( S(A) \) inherits a product \(*\) making it a Hopf algebra, isomorphic to the enveloping algebra of the Lie algebra \( A \) [13, 14]. Whenever it is possible, we can consider the dual Hopf algebra \( S(A)^* \) and its group of characters \( G \), which is the exponentiation, in some sense, of the Lie algebra \( A \).

We here consider an inverse construction, departing from a group used in Control Theory, namely the group of Fliess operators [3, 5, 6]; this group is used to study the feedback product. We limit ourselves here to the one-dimensional case. This group is the set \( \mathbb{R}\langle\langle x_0, x_1 \rangle\rangle \) of noncommutative formal series in two indeterminates, with a certain product generalizing the composition of formal series (definition 1). The Hopf algebra \( H \) of coordinates of this group is described in [5], where it is also proved that it is graded by the length of words; note that this gradation is not connected and not finite-dimensional. We first give a way to describe the composition in the group \( \mathbb{R}\langle\langle x_0, x_1 \rangle\rangle \) and the coproduct of \( H \) by induction on the length of words (lemma 2 and proposition 3). We prove that \( H \) admits a second gradation, which is connected; the dimensions of this gradation are given by the Fibonacci sequence (proposition 8). As the product of \( \mathbb{R}\langle\langle x_0, x_1 \rangle\rangle \) is left-linear, \( H \) is a commutative, right-sided combinatorial Hopf algebra [10], so, dually, \( \mathbb{R}\langle x_0, x_1 \rangle \) inherits a prelie product \( \cdot \), which is inductively defined in proposition 11. We prove that the words \( x_1^n, n \geq 0 \), form a minimal subset of generators of this prelie algebra (theorem 12).

The prelie algebra \( \mathbb{R}\langle x_0, x_1 \rangle \) has also an associative, commutative product, namely the shuffle product \( \shuffle \) [15]. We prove that the following axiom is satisfied (proposition 14):

\[
(x \shuffle y) \cdot z = (x \cdot z) \shuffle y + x \shuffle (y \cdot z).
\]

So \( \mathbb{R}\langle x_0, x_1 \rangle \) is a Com-Prelie algebra [11] (definition 15). We give a presentation of this Com-Prelie algebra in theorem 27. We use for this a description of free Com-Prelie algebras in terms of partitioned trees (definition 17), which generalizes the construction of prelie algebras in terms of rooted trees in [1]. We deduce a presentation of \( \mathbb{R}\langle x_0, x_1 \rangle \) as a prelie algebra in theorem 30. This presentation induces a new basis of \( \mathbb{R}\langle x_0, x_1 \rangle \) in terms of words with letters in \( \mathbb{N}^* \), see corollary 31. The prelie product of two elements of this basis uses a dendriform structure [2, 9] on the algebra of words with letters in \( \mathbb{N}^* \) (theorem 34). The study of this dendriform structure is postponed to the appendix, as well as the enumeration of partitioned trees; we also prove that free Com-Prelie algebras are free as prelie algebras, using Livernet’s rigidity theorem [7].

Aknowledgment. The research leading these results was partially supported by the French National Research Agency under the reference ANR-12-BS01-0017.

Notation. We denote by \( \mathbb{K} \) a commutative field of characteristic zero. All the objects (algebra, coalgebras, prelie algebras...) in this text will be taken over \( \mathbb{K} \).

1 Construction of the Hopf algebra

1.1 Definition of the composition

Let us consider an alphabet of two letters \( x_0 \) and \( x_1 \). We denote by \( \mathbb{K}\langle\langle x_0, x_1 \rangle\rangle \) the completion of the free algebra generated by this alphabet, that is to say the set of noncommutative formal
series in \(x_0\) and \(x_1\). Note that \(\mathbb{K}\langle\langle x_0, x_1\rangle\rangle\) is an algebra for the concatenation product and for the shuffle product, which we denote by \(\shuffle\).

**Exemps.** If \(a, b, c, d \in \{x_0, x_1\}\):

\[
\begin{align*}
abc\shuffle d &= abcd + abdc + adbc + dabc, \\
ab\shuffle cd &= abcd + acbd + cabd + acdb + cdab, \\
a\shuffle bcd &= abcd + bacd + bcad + bced.
\end{align*}
\]

The unit for both these products is the empty word, which we denote by \(\emptyset\). The algebra \(\mathbb{K}\langle\langle x_0, x_1\rangle\rangle\) is given its usual ultrametric topology.

**Definition 1** [5].

1. For any \(d \in \mathbb{K}\langle\langle x_0, x_1\rangle\rangle\), we define a continuous algebra map \(\varphi_d\) from \(\mathbb{K}\langle\langle x_0, x_1\rangle\rangle\) to \(\operatorname{End}(\mathbb{K}\langle\langle x_0, x_1\rangle\rangle)\) in the following way: for all \(X \in \mathbb{K}\langle\langle x_0, x_1\rangle\rangle\),

\[
\varphi_d(x_0)(X) = x_0X, \quad \varphi_d(x_1)(X) = x_1X + x_0(d\shuffle X).
\]

2. We define a composition \(\circ\) on \(\mathbb{K}\langle\langle x_0, x_1\rangle\rangle\) in the following way: for all \(c, d \in \mathbb{K}\langle\langle x_0, x_1\rangle\rangle\),

\[
c \circ d = \varphi_d(c)(\emptyset) + d.
\]

It is proved in [5] that this composition is associative.

**Notation.** For all \(c, d \in \mathbb{K}\langle\langle x_0, x_1\rangle\rangle\), we put \(c\tilde{o}d = c \circ d - d = \varphi_d(c)(\emptyset)\).

**Remark.** If \(c_1, c_2, d \in \mathbb{K}\langle\langle x_0, x_1\rangle\rangle\), \(\lambda \in \mathbb{K}\):

\[
(c_1 + \lambda c_2)\tilde{o}d = \varphi_d(c_1 + \lambda c_2)(\emptyset) = (\varphi_d(c_1) + \lambda \varphi_d(c_2))(\emptyset) = \varphi_d(c_1)(\emptyset) + \lambda \varphi_d(c_2)(\emptyset) = c_1\tilde{o}d + \lambda c_2\tilde{o}d.
\]

So the composition \(\tilde{o}\) is linear on the left. As \(\varphi_d\) is continuous, the map \(c \mapsto c\tilde{o}d\) is continuous for any \(d \in \mathbb{K}\langle\langle x_0, x_1\rangle\rangle\). Hence, it is enough to know how to compute \(c\tilde{d}d\) for any word \(c\), which is made possible by the next lemma, using an induction on the length:

**Lemma 2** For any word \(c\), for any \(d \in \mathbb{K}\langle\langle x_0, x_1\rangle\rangle\):

1. \(\emptyset\tilde{o}d = \emptyset\).
2. \((x_0c)\tilde{o}d = x_0(c\tilde{o}d)\).
3. \((x_1c)\tilde{o}d = x_1(c\tilde{o}d) + x_0(d\shuffle (c\tilde{o}d))\).

**Proof.**

1. \(\emptyset\tilde{o}d = \varphi_d(\emptyset)(\emptyset) = 1d(\emptyset) = \emptyset\).

2. \((x_0c)\tilde{o}d = \varphi_d(x_0c)(\emptyset) = \varphi_d(x_0) \circ \varphi_d(c)(\emptyset) = \varphi_d(x_0)(c\tilde{o}d) = x_0(c\tilde{o}d)\).

3. \((x_1c)\tilde{o}d = \varphi_d(x_1c)(\emptyset) = \varphi_d(x_1) \circ \varphi_d(c)(\emptyset) = \varphi_d(x_1)(c\tilde{o}d) = x_1(c\tilde{o}d) + x_0(d\shuffle (c\tilde{o}d))\). \(\square\)

### 1.2 Dual Hopf algebra

We here give a recursive description of the Hopf algebra of the coordinates of the group \(\mathbb{K}\langle\langle x_0, x_1\rangle\rangle\) of [5].

For any word \(c\), let us consider the map \(X_c \in \mathbb{K}\langle\langle x_0, x_1\rangle\rangle^*\), such that for any \(d \in \mathbb{K}\langle\langle x_0, x_1\rangle\rangle\), \(X_c(d)\) is the coefficient of \(c\) in \(d\). We denote by \(V\) the subspace of \(A^*\) generated by these maps.
Let $H = S(V)$, or equivalently the free associative, commutative algebra generated by the $X_c$’s. The elements of $H$ are seen as polynomial functions on $\mathbb{K}\langle\langle x_0, x_1 \rangle\rangle$; the elements of $H \otimes H$ are seen as polynomial functions on $\mathbb{K}\langle\langle x_0, x_1 \rangle\rangle \times \mathbb{K}\langle\langle x_0, x_1 \rangle\rangle$. Then $H$ is given a multiplicative coproduct defined in the following way: for any word $c$, for any $f, g \in \mathbb{K}\langle\langle x_0, x_1 \rangle\rangle$,$$
abla(X_c)(f, g) = X_c(f \circ g).$$As $\circ$ is associative, $\nabla$ is coassociative, so $H$ is a bialgebra.

Notations.
1. The space of words is a commutative algebra for the shuffle product $\mathbf{shuffle}$. Dually, the space $V$ inherits a coassociative, cocommutative coproduct, denoted by $\Delta_{\mathbf{shuffle}}$. For example, if $a, b, c \in \{x_0, x_1\}$:

   \[
   \Delta_{\mathbf{shuffle}}(X_0) = X_0 \otimes X_0,
   \]

   \[
   \Delta_{\mathbf{shuffle}}(X_a) = X_a \otimes X_0 + X_0 \otimes X_a,
   \]

   \[
   \Delta_{\mathbf{shuffle}}(X_{ab}) = X_{ab} \otimes X_0 + X_a \otimes X_b + X_b \otimes X_a + X_0 \otimes X_{ab},
   \]

   \[
   \Delta_{\mathbf{shuffle}}(X_{abc}) = X_{abc} \otimes X_0 + X_a \otimes X_{bc} + X_b \otimes X_{ac} + X_c \otimes X_{ab}
   \]

   \[
   + X_{ab} \otimes X_c + X_{ac} \otimes X_b + X_{bc} \otimes X_a + X_0 \otimes X_{abc}.
   \]

2. We define two linear endomorphisms $\theta_0, \theta_1$ of $V$ by $\theta_i(X_c) = X_{x_i c}$ for any word $c$.
   The following proposition allows to compute $\nabla(X_c)$ for any word $c$ by induction on the length.

   **Proposition 3** For all $x \in V$, we put $\tilde{\Delta}(x) = \Delta(x) - 1 \otimes x$.
   
   1. $\tilde{\Delta}(X_0) = X_0 \otimes 1$.
   2. $\tilde{\Delta} \circ \theta_0 = (\theta_0 \otimes \text{Id}) \circ \tilde{\Delta} + (\theta_1 \otimes m) \circ (\tilde{\Delta} \otimes \text{Id}) \circ \Delta_{\mathbf{shuffle}}$.
   3. $\tilde{\Delta} \circ \theta_1 = (\theta_1 \otimes \text{Id}) \circ \tilde{\Delta}$.

   **Proof.** For any word $c$, for any $f, g \in \mathbb{K}\langle\langle x_0, x_1 \rangle\rangle$:

   \[
   \tilde{\Delta}(X_c)(f, g) = \Delta(X_c)(f, g) - (1 \otimes X_c)(f, g) = X_c(f \circ g) - X_c(g) = X_c(f \otimes g - g) = X_c(f \circ g).
   \]

   As $\circ$ is linear on the left, $\tilde{\Delta}(X_c) \in V \otimes H$, so formulas in points 2 and 3 make sense.

   Let $f \in \mathbb{K}\langle\langle x_0, x_1 \rangle\rangle$. It can be uniquely written as $f = x_0 f_0 + x_1 f_1 + \lambda 0$, with $f_0, f_1 \in \mathbb{K}\langle\langle x_0, x_1 \rangle\rangle, \lambda \in K$. For all $g \in \mathbb{K}\langle\langle x_0, x_1 \rangle\rangle$:

   \[
   f \circ g = (x_0 f_0) \circ g + (x_1 f_1) \circ g + \lambda 0 \circ g = x_0 (f_0 \circ g + g \mathbf{shuffle}(f_1 \circ g)) + x_1 (f_1 \circ g) + \lambda 0.
   \]

   1. We obtain:

   \[
   \tilde{\Delta}(X_0)(f, g) = X_0(x_0(f_0 \circ g + g \mathbf{shuffle}(f_1 \circ g)) + x_1(f_1 \circ g) + \lambda 0) = 0 + 0 + \lambda = (X_0 \otimes 1)(f, g),
   \]

   so $\tilde{\Delta}(X_0) = X_0 \otimes 1$.

   2. Let $c$ be a word.

   \[
   \tilde{\Delta} \circ \theta_0(X_c)(f, g) = \tilde{\Delta}(X_{x_0 c})(f, g)
   \]

   \[
   = X_{x_0 c}(x_0(f_0 \circ g + g \mathbf{shuffle}(f_1 \circ g)) + x_1(f_1 \circ g) + \lambda 0)
   \]

   \[
   = X_c(f_0 \circ g + g \mathbf{shuffle}(f_1 \circ g)) + x_1(f_1 \circ g) + \lambda 0
   \]

   \[
   = X_c(f_0 \circ g + (f_1 \circ g) \mathbf{shuffle}) + x_1(f_1 \circ g) + \lambda 0
   \]

   \[
   = \tilde{\Delta}(X_c)(f_0, g) + (\tilde{\Delta} \otimes \text{Id}) \circ \Delta_{\mathbf{shuffle}}(X_c)(f_1, g)
   \]

   \[
   = \tilde{\Delta}(X_c)(f_0, g) + (\text{Id} \otimes m) \circ (\tilde{\Delta} \otimes \text{Id}) \circ \Delta_{\mathbf{shuffle}}(X_c)(f_1, g)
   \]

   \[
   = (\theta_0 \otimes \text{Id}) \circ \tilde{\Delta}(X_c)(f, g) + (\theta_1 \otimes \text{Id}) \circ (\text{Id} \otimes m) \circ (\tilde{\Delta} \otimes \text{Id}) \circ \Delta_{\mathbf{shuffle}}(X_c)(f, g),
   \]
so $\tilde{\Delta} \circ \theta_0(X_c) = (\theta_0 \otimes \text{Id}) \circ \tilde{\Delta}(X_c) + (\theta_1 \otimes \text{Id}) \circ (\text{Id} \otimes m) \circ (\tilde{\Delta} \otimes \text{Id}) \circ \Delta_w(X_c)$.

3. Let $c$ be a word.

$$\tilde{\Delta} \circ \theta_1(X_c)(f, g) = \tilde{\Delta}(X_{x^0c})(f, g) = X_{x^1c}(x_0(f_0 \tilde{\circ} g + g \mu(f_1 \tilde{\circ} g)) + x_1(f_1 \tilde{\circ} g) + \lambda 0) = 0 + X_c(f_1 \tilde{\circ} g) + 0 = \tilde{\Delta}(X_c)(f_1, g) = (\theta_1 \otimes \text{Id}) \circ \tilde{\Delta}(X_c)(f, g),$$

so $\tilde{\Delta} \circ \theta_1(X_c) = (\theta_1 \otimes \text{Id}) \circ \tilde{\Delta}(X_c)$.

**Examples.**

$\Delta(X_{x^0}) = X_{x^0} \otimes 1 + 1 \otimes X_{x^0} + X_{x^1} \otimes X_{\emptyset},$

$\Delta(X_{x^2}) = X_{x^2} \otimes 1 + 1 \otimes X_{x^2} + X_{x^0 x^1} \otimes X_{\emptyset} + X_{x^1 x_0} \otimes X_{\emptyset} + X_{x^1 x_1} \otimes X_{\emptyset}^2 + X_{x^1} \otimes X_{x^0},$

$\Delta(X_{x^0 x^1}) = X_{x^0 x^1} \otimes 1 + 1 \otimes X_{x^0 x^1} + X_{x^1 x_0} \otimes X_{\emptyset} + X_{x^1} \otimes X_{x^1},$

$\Delta(X_{x^1 x_0}) = X_{x^1 x_0} \otimes 1 + 1 \otimes X_{x^1 x_0} + X_{x^1 x_1} \otimes X_{\emptyset}.$

**Corollary 4** For all $n \geq 1$, $\tilde{\Delta}(X_{x^1^n}) = X_{x^1^n} \otimes 1$ and $\Delta(X_{x^1^n}) = X_{x^1^n} \otimes 1 + 1 \otimes X_{x^1^n}$.

**Proof.** It comes from an easy induction on $n$.

### 1.3 gradation

It is proved in [5] that the Hopf algebra $H$ is graded by the length of words, but this gradation is not connected, that is to say that the homogeneous component of degree 0 is not $(0)$, as it contains $X_{\emptyset}$. Moreover, the homogeneous components of $H$ are not finite-dimensional, as for example $X_{\emptyset}^n X_{x^0}$ is homogeneous of degree $k$ for all $n \geq 0$. We now define another gradation on $H$, which is connected and finite-dimensional.

**Definition 5** 1. Let $c = c_1 \ldots c_n$ be a word. We put:

$$\text{deg}(c) = n + 1 + \sharp \{ i \in \{1, \ldots, n\} \mid c_i = x_0 \}.$$

2. For all $k \geq 1$, we put:

$$V_k = \text{Vect}(X_c \mid \text{deg}(x) = k).$$

This define a connected gradation of $V$, that is to say:

$$V = \bigoplus_{k \geq 1} V_k.$$

3. This gradation induces a connected gradation of the algebra $H$:

$$H = \bigoplus_{k \geq 0} H_k, \text{ and } H_0 = \mathbb{K}.$$

**Lemma 6** If $c$ is a word of degree $n$, then:

$$\tilde{\Delta}(X_c) \in \bigoplus_{i+j=n} V_i \otimes H_j.$$

**Proof.** Let us start by the following observations:
1. Let \( c \) be a word of degree \( k \). Then \( x_0c \) is a word of degree \( k + 2 \). Hence, \( \theta_0 \) is homogeneous of degree 2 on \( V \).

2. Let \( c \) be a word of degree \( k \). Then \( x_1c \) is a word of degree \( k + 1 \). Hence, \( \theta_1 \) is homogeneous of degree 1 on \( V \).

3. Let \( c \) and \( d \) be two words of respective degrees \( k \) and \( l \). Then any word obtained by shuffling \( c \) and \( d \) is of degree \( k + l - 1 \): its length is the sum of the length of \( c \) and \( d \), and the number of \( x_0 \) in it is the sum of the numbers of \( x_0 \) in \( c \) and \( d \). Hence, dually, the coproduct \( \Delta_w \) is homogeneous of degree 1 from \( V \) to \( V \otimes V \).

Let us prove the result by induction on the length \( k \) of \( c \). If \( k = 0 \), then \( c = \emptyset \) so \( n = 1 \), and \( \tilde{\Delta}(X_c) = X_c \otimes 1 \in V_1 \otimes H_0 \). Let us assume the result for all words of length \( < k - 1 \). Two cases can occur.

1. If \( c = x_0d \), then \( \text{deg}(d) = n - 2 \). we put \( \Delta_w(X_d) = \sum x'_i \otimes x''_i \). By the preceding third observation, we can assume that for all \( i \), \( x'_i, x''_i \) are homogeneous elements of \( V \), with \( \text{deg}(x'_i) + \text{deg}(x''_i) = n - 2 + 1 = n - 1 \). Then:

\[
\tilde{\Delta}(X_c) = (\theta_0 \otimes \text{Id}) \circ \tilde{\Delta}(X_d) + \sum_i (\theta_1 \otimes m) \circ (\tilde{\Delta}(x'_i) \otimes x''_i).
\]

By the induction hypothesis, \( \tilde{\Delta}(X_d) \in (V \otimes H)_{n-1} \). By the second observation, \( (\theta_0 \otimes \text{Id}) \circ \tilde{\Delta}(X_d) \in (V \otimes H)_n \). By the induction hypothesis applied to \( x'_i \), for all \( i \), \( (\tilde{\Delta}(x'_i) \otimes x''_i) \in (V \otimes H \otimes V)_{n-1} \), so by the first observation, \( (\theta_1 \otimes m) \circ (\Delta(x'_i) \otimes x''_i) \in (V \otimes H)_{n+1} \subseteq (V \otimes H)_n \). So \( \Delta(X_c) \in (V \otimes H)_n \).

2. \( c = x_1d \), then \( \text{deg}(d) = n - 1 \). Moreover, \( \tilde{\Delta}(X_c) = (\theta_1 \otimes \text{Id}) \circ \tilde{\Delta}(X_d) \). By the induction hypothesis, \( \tilde{\Delta}(X_d) \in (V \otimes H)_{n-1} \). By the second observation, \( \Delta(X_c) \in (V \otimes H)_n \).

So the result holds for any word \( c \).

\[\square\]

**Proposition 7** With this gradation, \( H \) is a graded, connected Hopf algebra.

**Proof.** We have to prove that for all \( n \geq 0 \):

\[\Delta(H_n) \subseteq \bigoplus_{i+j=n} H_i \otimes H_j.\]

This comes from the multiplicativity of \( \Delta \).

\[\square\]

Let us now study the formal series of \( V \) and \( H \).

**Proposition 8** \( 1. \) For all \( k \), let us put \( p_k = \text{dim}(V_k) \) and \( F_V = \sum_{k=1}^{\infty} p_k X^k \). Then:

\[F_V = \frac{X}{1 - X - X^2},\]

and for all \( k \geq 1 \):

\[p_k = \frac{1}{\sqrt{5}} \left( \left( \frac{1 + \sqrt{5}}{2} \right)^k - \left( \frac{1 - \sqrt{5}}{2} \right)^k \right).\]

This is the Fibonacci sequence (A000045 in [16]).
2. We put $F_H = \sum_{k=0}^{\infty} \dim(H_k)X^k$. Then:

$$F_H = \prod_{k=1}^{\infty} \frac{1}{(1 - X^k)^{p_k}}.$$ 

**Proof.** Let us consider the formal series:

$$F(X_0, X_1) = \sum_{i,j \geq 0} \# \{\text{words in } x_0, x_1 \text{ with } i x_0 \text{ and } j x_1 \} X_i X_j.$$ 

Then $F(X_0, X_1) = \frac{1}{1 - X_0 - X_1}$. Moreover, by definition of the degree of a word:

$$F_V = XF(X^2, X) = \frac{X}{1 - X - X^2}.$$ 

As $H$ is the symmetric algebra generated by $V$, its formal series is given by the second point. □

**Examples.** We obtain:

<table>
<thead>
<tr>
<th>$k$</th>
<th>0</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
<th>5</th>
<th>6</th>
<th>7</th>
<th>8</th>
<th>9</th>
<th>10</th>
</tr>
</thead>
<tbody>
<tr>
<td>$\dim(V_k)$</td>
<td>0</td>
<td>1</td>
<td>2</td>
<td>3</td>
<td>4</td>
<td>5</td>
<td>6</td>
<td>7</td>
<td>8</td>
<td>13</td>
<td>21</td>
</tr>
<tr>
<td>$\dim(H_k)$</td>
<td>1</td>
<td>1</td>
<td>2</td>
<td>3</td>
<td>4</td>
<td>8</td>
<td>15</td>
<td>30</td>
<td>56</td>
<td>108</td>
<td>203</td>
</tr>
</tbody>
</table>

The third row is sequence A166861 of [16].

**Remark.** Consequently, the space $V$ inherits a bigradation:

$$V_{k,n} = \text{Vect}(X_c \mid \deg(c) = k \text{ and } lg(c) = n).$$

If $c$ is a word of length $n$ and of degree $k$, denoting by $a$ the number of its letters equal to $x_0$ and by $b$ the number of its letters equal to $x_1$, then:

$$\begin{cases} a + b = n, \\ 2a + b + 1 = k, \end{cases}$$

so $a = k - n - 1$. Hence:

$$\dim(V_{k,n}) = \binom{n}{k - n - 1},$$

and the formal series of this bigradation is:

$$\sum_{k,n \geq 0} \dim(V_{k,n})X^kY^n = \frac{X}{1 - XY - X^2Y}.$$ 

## 2 Prelie structure on $\mathbb{K}\langle x_0, x_1 \rangle$

### 2.1 Prelie coproduct on $V$

As the composition $\circ$ is linear on the left, the dual coproduct satisfies $\tilde{\Delta}(V) \subseteq V \otimes H$, so $H$ is a commutative right-sided Hopf algebra in the sense of [10], and $V$ inherits a right prelie coproduct: if $\pi$ is the canonical projection from $H = S(V)$ onto $V$, 

$$\delta = (\pi \otimes \pi) \circ \Delta = (\text{Id} \otimes \pi) \circ \tilde{\Delta}.$$ 

It satisfies the right prelie coalgebra axiom:

$$(23).((\delta \otimes \text{Id}) \circ \delta - (\text{Id} \otimes \delta) \circ \delta) = 0.$$ 

The following proposition allows to compute $\delta(X_c)$ by induction on the length of $c$. 


Proposition 9  
1. $\delta(X_{\emptyset}) = 0$. 
2. $\delta \circ \theta_0 = (\theta_0 \otimes Id) \circ \delta + (\theta_1 \otimes Id) \circ \Delta_w$. 
3. $\delta \circ \theta_1 = (\theta_1 \otimes Id) \circ \delta$.

Proof. The first point comes from $\Delta(X_{\emptyset}) = X_{\emptyset} \otimes 1 + 1 \otimes X_{\emptyset}$. Let $x \in V$. We put $\Delta_w(x) = x' \otimes x'' \in V \otimes V$. For any $y \in V$, we put $\Delta(y) - y \otimes 1 = y^{(1)} \otimes y^{(2)} \in V \otimes H_+$. Then:

$$
(\theta_1 \otimes m) \circ (\Delta \otimes Id) \circ \Delta_w(x) = (\theta_1 \otimes m)(x' \otimes 1 \otimes x'' + x^{(1)} \otimes x^{(2)} \otimes x'')
$$

$$
= \theta_1(x') \otimes x'' + x^{(1)} \otimes x^{(2)} \otimes x'' 
$$

Applying $Id \otimes \pi$, it remains:

$$
(Id \otimes \pi) \circ (\theta_1 \otimes m) \circ (\Delta \otimes Id) \circ \Delta_w(x) = (\theta_1 \otimes Id) \circ \Delta_w(x).
$$

Let $i = 0$ or 1. Then:

$$(Id \otimes \pi) \circ (\theta_i \otimes Id) \circ \tilde{\Delta} = (\theta_i \otimes Id) \circ (Id \otimes \pi) \circ \tilde{\Delta} = (\theta_i \otimes Id) \circ \delta.$$

The result is induced by these remarks, combined with proposition 3. \qed

Examples.

$$
\delta(X_{x_0}) = X_{x_1} \otimes X_{\emptyset}, \\
\delta(X_{x_0^2}) = X_{x_0 x_1} \otimes X_{\emptyset} + X_{x_1 x_0} \otimes X_{\emptyset} + X_{x_1} \otimes X_{x_0}, \\
\delta(X_{x_0 x_1}) = X_{x_1 x_1} \otimes X_{\emptyset} + X_{x_1} \otimes X_{x_1}, \\
\delta(X_{x_1 x_0}) = X_{x_1 x_1} \otimes X_{\emptyset}.
$$

Proposition 10 $Ker(\delta) = Vect(X_{x_1^k}, n \geq 0)$.

Proof. The inclusion $\supseteq$ is trivial by corollary 4. Let us prove the other inclusion.

First step. Let us prove the following property: if $x \in V_k$ is such that

$$
\delta(x) = \lambda \sum_{i+j=k-2} \frac{(k-2)!}{i! j!} X_{x_{i}^1} \otimes X_{x_{j}^1},
$$

then there exists $\mu \in K$ such that $x = \mu x_1^{k-1}$. It is obvious if $k = 1$, as then $x = \mu \emptyset$. Let us assume the result at all ranks $< k$. We put $x = x_1^\alpha (x_0 f_0 + x_1 f_1)$, where $\alpha \geq 0$, $f_0$ is homogeneous of degree $k - 2 - \alpha$ and $f_1$ is homogeneous of degree $k - 1 - \alpha$.

$$
\delta(x) = (\theta_1^\alpha \otimes Id) ((\theta_0 \otimes Id) \circ \delta(f_0) + (\theta_1 \otimes Id) \circ \delta(f_1) + (\theta_1 \otimes Id) \circ \Delta_w(f_0)).
$$

Let us consider the terms in this expression of the form $X_{\emptyset} \otimes X_e$, with $c$ a word. This gives:

$$
\lambda X_{\emptyset} \otimes X_{x_1^{k-2}} = 0,
$$

so $\lambda = 0$ and $\delta(x) = 0$. Let us now consider the terms of the form $X_{x_1^c} \otimes X_{x_1^d}$, with $c, d$ words. We obtain:

$$
(\theta_1^\alpha \otimes \theta_0 \otimes Id) \circ \delta(f_0) = 0.
$$

As both $\theta_0$ and $\theta_1$ are injective, we obtain $\delta(f_0) = 0$. By the induction hypothesis, $f_0 = \nu X_{x_1^l}$, with $l = k - 2 - \alpha < k$. Hence:

$$
\Delta_w(f_0) = \nu \sum_{i+j=l} \frac{l!}{i! j!} X_{x_{i}^1} \otimes X_{x_{j}^1},
$$

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and:
\[
(\theta_1^{a+1} \otimes \text{Id}) \left( \delta(f_1) + \nu \sum_{i+j=l} \frac{(k-a-2)!}{i!j!} X_{x_i} \otimes X_{x_j} \right) = 0.
\]

As \( \theta_1 \) is injective, we obtain with the induction hypothesis that \( f_1 = \mu X_{x_i^{k-2}} \), so:
\[
x = \mu X_{x_i^{k-1}} + \nu X_{x_i x_0 x_i^{k-2}}.
\]

This gives:
\[
\delta(x) = \nu (\theta_1^{a+1} \otimes \text{Id}) \left( \sum_{i+j=k-a-2} \frac{(k-a-2)!}{i!j!} X_{x_i} \otimes X_{x_j} \right)
\]
\[
= \nu \sum_{i+j=k-a-2} \frac{(k-a-2)!}{i!j!} X_{x_i} \otimes X_{x_j}
\]
\[
= 0,
\]
so necessarily \( \nu = 0 \) and \( x = \mu X_{x_i^{k-1}} \).

**Second step.** Let \( x \in \text{Ker}(\delta) \). As \( \delta \) is homogeneous of degree 0, the homogeneous components of \( x \) are in \( \text{Ker}(\delta) \). By the first step, with \( \lambda = 0 \), these homogeneous components, hence \( x \), belong to \( \text{Vect}(X_{x_i^{k}}, k \geq 0) \). \( \square \)

### 2.2 Dual prelie algebra

As \( V \) is a graded prelie coalgebra, its graded dual is a prelie algebra. We identify this graded dual with \( \mathbb{K}\langle x_0, x_1 \rangle \subseteq \mathbb{K}(\langle x_0, x_1 \rangle) \); for any words \( c, d \), \( X_c(d) = \delta_{c,d} \). The prelie product of \( \mathbb{K}\langle x_0, x_1 \rangle \) is denoted by \( \bullet \). Dualizing proposition 9, we obtain:

**Proposition 11**

1. For all word \( c \), \( \emptyset \bullet c = 0 \).
2. For all words \( c, d \), \( (x_0 c) \bullet d = x_0 (c \bullet d) \).
3. For all words \( c, d \), \( (x_1 c) \bullet d = x_1 (c \bullet d) + x_0 (c \sqcup d) \).

**Proof.** Let \( u, v, w \) be words. Then \( X_w(u \bullet v) = \delta(X_w)(u \otimes v) \). Hence, if \( d \) is a word:
\[
X_{\emptyset}(u \bullet v) = 0,
\]
\[
X_{x_0 d}(u \bullet v) = (\theta_0 \otimes \text{Id}) \circ \delta(X_d)(u \otimes v) + (\theta_1 \otimes \text{Id}) \circ \Delta_{x_0} (X_d)(u \otimes v)
\]
\[
= X_d(\theta_0(u) \bullet v) + \theta_1(u) \sqcup v,
\]
\[
X_{x_1 d}(u \bullet v) = (\theta_1 \otimes \text{Id}) \circ \delta(X_d)(u \otimes v)
\]
\[
= X_d(\theta_1(u) \bullet v).
\]
Moreover, for all word \( c \):
\[
\theta_0^c(\emptyset) = 0, \quad \theta_0^c(x_0 c) = c, \quad \theta_0^c(x_1 c) = 0,
\]
\[
\theta_1^c(\emptyset) = 0, \quad \theta_1^c(x_0 c) = 0, \quad \theta_1^c(x_1 c) = c.
\]
Hence, for any words \( c, d \):
\[
X_{x_0 d}(x_0 c \bullet v) = X_d(c \bullet v) = X_{x_0 d}(x_1 c \bullet v) = X_d(c \sqcup v) = X_{x_0 d}(x_1 c \bullet v) + x_0 (c \sqcup v),
\]
\[
X_{x_1 d}(x_0 c \bullet v) = 0 = X_{x_1 d}(x_1 c \bullet v) = X_d(c \bullet v) = X_{x_1 d}(x_1 c \bullet v) + x_0 (c \sqcup v).
\]
Hence, for any \( w \), \( X_w(x_0 c \bullet v) = X_w(x_0 (x \bullet v)) \) and \( X_w(x_1 c \bullet v) = X_w((x_1 (c \bullet v) + x_0 (c \bullet w)) \).

Examples.

\[
\begin{align*}
    x_0 \bullet x_0 &= 0 & x_0 \bullet x_0 x_0 &= 0 & x_1 \bullet x_0 x_0 &= x_0 x_0 x_0 \\
    x_0 \bullet x_1 &= 0 & x_0 \bullet x_0 x_1 &= 0 & x_1 \bullet x_0 x_1 &= x_0 x_0 x_1 \\
    x_1 \bullet x_0 &= x_0 x_0 & x_0 \bullet x_1 x_0 &= 0 & x_1 \bullet x_1 x_0 &= x_0 x_1 x_0 \\
    x_1 \bullet x_1 &= x_0 x_1 & x_0 \bullet x_1 x_1 &= 0 & x_1 \bullet x_1 x_1 &= x_0 x_1 x_1 \\
\end{align*}
\]

\[
\begin{align*}
    x_0 x_0 \bullet x_0 &= 0 & x_0 x_0 \bullet x_1 &= 0 \\
    x_0 x_1 \bullet x_0 &= x_0 x_0 x_0 & x_0 x_1 \bullet x_1 &= x_0 x_0 x_1 \\
    x_1 x_0 \bullet x_0 &= 2 x_0 x_0 x_0 & x_1 x_0 \bullet x_1 &= x_0 x_0 x_1 + x_0 x_1 x_0 \\
    x_1 x_1 \bullet x_0 &= x_1 x_0 x_0 + x_0 x_1 x_0 + x_0 x_0 x_1 & x_1 x_1 \bullet x_1 &= x_1 x_0 x_1 + 2 x_0 x_1 x_1 \\
\end{align*}
\]

Dualizing proposition 10:

**Theorem 12** \( \mathbb{K}(x_0, x_1) = Vect(x^n_0, n \geq 0) \oplus (\mathbb{K}(x_0, x_1) \bullet \mathbb{K}(x_0, x_1)) \). Hence, \( (x^n_0)_{n \geq 0} \) is a minimal system of generators of the prelie algebra \( \mathbb{K}(x_0, x_1) \).

**Proof.** As \( \bullet = \delta^* \), \( \text{Im}(\bullet) = \text{Ker}(\delta)^\perp = Vect(X_{x^n_0}, n \geq 0)^\perp \). The first assertion is then immediate. As \( \mathbb{K}(\langle x_0, x_1 \rangle) \) is a graded, connected prelie coalgebra, \( \mathbb{K}(x_0, x_1) \) is a graded, connected prelie algebra. The result then follows from the next lemma. \( \square \)

**Lemma 13** Let \( A \) be a graded, connected prelie algebra, and \( V \) be a graded subspace of \( A \).

1. \( V \) generates \( A \) if, and only if, \( A = V + A \bullet A \).

2. \( V \) is a minimal subspace of generators of \( A \) if, and only if, \( A = V \oplus A \bullet A \).

**Proof.** 1. \( \implies \). Let \( x \in A \). Then it can be written as an element of the prelie subalgebra generated by \( v \), so as the sum of an element of \( V \) and of iterated prelie products of elements of \( V \). Hence, \( x \in V + A \bullet A \). Note that we did not use the gradation of \( A \) to prove this point.

1. \( \iff \). Let \( B \) be the prelie subalgebra generated by \( V \). Let \( x \in A_n \), let us prove that \( x \in B \) by induction on \( n \). As \( A_0 = (0) \), it is obvious if \( n = 0 \). Let us assume the result at all ranks \( < n \). We obtain, by the gradation:

\[
A_n = V_n \oplus \sum_{i=1}^{n-1} A_i \bullet A_{n-i}.
\]

So we can write \( x = \lambda x_1^{n-1} + \sum x_i \bullet y_i \), where \( x_i, y_i \) are homogeneous of degree \( < n \). By the induction hypothesis, these elements belong to \( B \), so \( x \in B \).

2. \( \implies \). By 1. \( \implies \), \( A = V + A \bullet A \). If \( V \cap A \bullet A \neq (0) \), we can choose a graded subspace \( W \subseteq V \), such that \( A = W \oplus A \bullet A \). By 1. \( \iff \), \( W \) generates \( A \), so \( V \) is not a minimal system of generators of \( A \): contradiction. So \( A = V \oplus A \bullet A \).

2. \( \iff \). By 1. \( \iff \), \( V \) is a space of generators of \( A \). If \( W \subseteq V \), then \( W \oplus A \bullet A \subseteq A \). By 1. \( \implies \), \( W \) does not generate \( V \). So \( V \) is a minimal subspace of generators. \( \square \)

**Proposition 14** For all \( x, y, z \in \mathbb{K}(x_0, x_1) \), \( (x \bullet y) \bullet z = (x \bullet z) \bullet y + x \bullet (y \bullet z) \).
Proof. We prove it if \(x, y, z\) are words. If \(x = \emptyset\), then:

\[
(\emptyset \cdot y) \cdot z = y \cdot z = (\emptyset \cdot z) \cdot y + \emptyset \cdot (u \cdot z).
\]

If \(y = \emptyset\), the result is also true, using the commutativity of \(\cdot\). We can now consider that \(x, y\) are nonempty words.

Let us proceed by induction on \(k = l_2(x) + l_2(y)\). If \(k = 0\) or 1, there is nothing to prove. Let us assume the result at all rank \(< k\). Four cases can occur.

First case. \(x = x_0a\) and \(y = x_0b\). Then:

\[
(x \cdot y) \cdot z = (x_0(a \cdot x_0b)) \cdot z + (x_0(x_0a \cdot b)) \cdot z,
\]

\[
= x_0((a \cdot x_0b) \cdot z) + x_0((x_0a \cdot b) \cdot z).
\]

Second case. \(x = x_1a\) and \(y = x_0b\). This gives:

\[
(x \cdot y) \cdot z = (x_1(a \cdot x_0b)) \cdot z + (x_0(x_1a \cdot b)) \cdot z
\]

\[
= x_1((a \cdot x_0b) \cdot z) + x_0((x_1a \cdot b) \cdot z) + x_0((a \cdot x_0b) \cdot z).
\]

Third case. \(x = x_0a\) and \(y = x_1b\). This is a consequence of the second case, using the commutativity of \(\cdot\).

Last case. \(x = x_1a\) and \(y = x_1b\). Similar computations give:

\[
(x \cdot y) \cdot z = x_1((a \cdot x_1b) \cdot w) + x_1(a \cdot x_1b \cdot z).
\]

So the result holds in all cases. □
3 Presentation of $\mathbb{K}\langle x_0, x_1 \rangle$ as a Com-Prelie algebra

Proposition 14 motivates the following definition:

**Definition 15** [11] A Com-Prelie algebra is a triple $(V, \bullet, \triangleright)$, such that:

1. $(V, \bullet)$ is a prelie algebra.
2. $(V, \triangleright)$ is a commutative, associative algebra (non necessarily unitary).
3. For all $a, b, c \in V$, $(a \triangleright b) \bullet c = (a \bullet c) \triangleright b + a \triangleright (b \bullet c)$.

For example, $\mathbb{K}\langle x_0, x_1 \rangle$ is a Com-Prelie algebra. See [11] for an example of Com-Prelie algebra based on rooted trees.

3.1 Free Com-Prelie algebras

**Definition 16** 1. A partitioned forest is a pair $(F, I)$ such that:

(a) $F$ is a rooted forest (the edges of $F$ being oriented from the leaves to the roots).
(b) $I$ is a partition of the vertices of $F$ with the following condition: if $x, y$ are two vertices of $F$ which are in the same part of $I$, then either they are both roots, or they have the same direct descendant.

2. We shall say that a partitioned forest is a partitioned tree if all the roots are in the same part of the partition.

3. Let $D$ be a set. A partitioned tree decorated by $D$ is a pair $(t, d)$, where $t$ is a partitioned tree and $d$ is a map from the set of vertices of $t$ into $D$. For any vertex $x$ of $t$, $d(x)$ is called the decoration of $x$.

4. The set of isoclasses of partitioned trees will be denoted by $\mathcal{PT}$. For any set $D$, the set of isoclasses of partitioned trees decorated by $D$ will be denoted by $\mathcal{PT}(D)$.

**Examples.** We represent partitioned trees by the Hasse graph of the underlying rooted forest, the partition being represented by horizontal edges, of different colors. Here are all the partitioned trees with $\leq 4$ vertices:

$\psi, 1, \psi : V, V, V, 1, 1 = J, \cdots ; V, V ; V, V, 1 = \psi, \psi, 1, 1 = 1, 1 = 1, 1 = 1, 1 = 1, 1 = 1, 1 = 1, 1 = 1, 1 = 1$.

**Definition 17** Let $t = (t, I)$ and $t' = (t', J) \in \mathcal{PT}$.

1. Let $s$ be a vertex of $t'$. The partitioned tree $t \bullet_s t'$ is defined as follows:

(a) As a rooted forest, $t \bullet_s t'$ is obtained by grafting all the roots of $t'$ on the vertex $s$ of $t$.
(b) We put $I = \{I_1, \ldots, I_k\}$ and $J = \{J_1, \ldots, J_l\}$. The partition of the vertices of this rooted forest is $I \sqcup J = \{I_1, \ldots, I_k, J_1, \ldots, J_l\}$.

2. The partitioned tree $t \triangleright t'$ is defined as follows:

(a) As a rooted forest, $t \triangleright t'$ is $tt'$.
(b) We put $I = \{I_1, \ldots, I_k\}$ and $J = \{J_1, \ldots, J_l\}$ and we assume that the set of roots of $t$ is $I_1$ and the set of roots of $t'$ is $J_1$. The partition of the vertices of $t \triangleright t'$ is $\{I_1 \sqcup J_1, I_2, \ldots, I_k, J_2, \ldots, J_l\}$.
Examples.

1. Here are the three possible graftings $\nabla \cdot_s \vdash \nabla$, $\nabla$, and $\nabla$.

2. Here are the two possible graftings $1 \cdot_s \vdash \nabla$ and $\nabla$.

These operations can also be defined for decorated partitioned trees.

**Proposition 18** Let $\mathcal{D}$ be a set. We denote by $\mathfrak{g}_{\mathcal{PT}(\mathcal{D})}$ the vector space generated by $\mathcal{PT}(\mathcal{D})$. We extend $\mathfrak{w}$ by bilinearity on $\mathfrak{g}_{\mathcal{PT}(\mathcal{D})}$ and we define a second product $\bullet$ on $\mathfrak{g}_{\mathcal{PT}(\mathcal{D})}$ in the following way: if $t, t' \in \mathcal{PT}(\mathcal{D})$,

$$t \bullet t' = \sum_{s \in V(t)} t \bullet_s t'.$$

Then $(\mathfrak{g}_{\mathcal{PT}(\mathcal{D})}, \bullet, \mathfrak{w})$ is a Com-Prelie algebra.

**Proof.** Let $t, t', t''$ be three partitioned trees.

If $s', s''$ are two vertices of $t$, we define by $t \bullet_{s', s''} (t', t'')$ the partitioned trees obtained by grafting the roots of $t'$ on $s'$, the roots of $t''$ on $s''$, the partition of the vertices of the obtained rootes forest being the union of the partitions of $t, t'$ and $t''$. Then:

$$(t \bullet t') \bullet t'' = \sum_{s' \in V(t)} (t \bullet_{s'} t') \bullet t''$$

$$= \sum_{s', s'' \in V(t)} (t \bullet_{s'} t') \bullet_{s''} t'' + \sum_{s' \in V(t), s'' \in V(t')} (t \bullet_{s'} t') \bullet_{s''} t''$$

$$= \sum_{s', s'' \in V(t)} t \bullet_{s', s''} (t', t'') + \sum_{s'' \in V(t')} t \bullet_{s''} (t', t'')$$

$$= \sum_{s', s'' \in V(t)} t \bullet_{s', s''} (t', t'') + t \bullet (t' \bullet t'').$$

So $(t \bullet t') \bullet t'' - t \bullet (t' \bullet t'')$ is clearly symmetric in $t$ and $t'$, and $\bullet$ is prelie.

Moreover, $(t \mathfrak{w} t') \mathfrak{w} t'' = t \mathfrak{w} (t' \mathfrak{w} t'')$ is the rooted forest $tt't''$, the partition being $\{ I_1 \sqcup J_1 \sqcup K_1, I_2, \ldots, I_k, J_2, \ldots, J_l, K_2, \ldots, K_m \}$, with immediate notations; $t \mathfrak{w} t' = t \mathfrak{w} t$ is the rooted forest $tt'$, the partition being $\{ I_1 \sqcup J_1, I_2, \ldots, I_k, J_2, \ldots, J_l \}$. So $\mathfrak{w}$ is an associative, commutative product.

Finally:

$$(t \mathfrak{w} t') \bullet t'' = \sum_{s \in V(t)} (t \mathfrak{w} t') \bullet_{s} t'' + \sum_{s' \in V(t')} (t \mathfrak{w} t') \bullet_{s'} t''$$

$$= \sum_{s \in V(t)} (t \bullet_{s} t'') \mathfrak{w} t' + \sum_{s' \in V(t')} t \mathfrak{w} (t' \bullet_{s'} t'')$$

$$= (t \bullet t') \mathfrak{w} t'' + t \mathfrak{w} (t' \bullet t'').$$

So $\mathfrak{g}_{\mathcal{PT}(\mathcal{D})}$ is Com-Prelie.

In particular, $\mathfrak{g}_{\mathcal{PT}(\mathcal{D})}$ is prelie. Let us use the extension of the prelie product $\bullet$ to $S(\mathfrak{g}_{\mathcal{PT}(\mathcal{D})})$ defined by Oudom and Guin [13, 14]:

1. If $t_1, \ldots, t_k \in \mathfrak{g}_{\mathcal{PT}(\mathcal{D})}$, $t_1 \ldots t_k \bullet 1 = t_1 \ldots t_k$.

2. If $t, t_1, \ldots, t_k \in \mathfrak{g}_{\mathcal{PT}(\mathcal{D})}$, $t \bullet t_1 \ldots t_k = (t \bullet t_1 \ldots t_{k-1}) \bullet t_k - t \bullet (t_1 \ldots t_{k-1} \bullet t_k)$. 

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3. If \( a, b, c \in S(\mathfrak{g}_{PT(D)}) \), \( ab \cdot c = (a \cdot c^{(1)})(b \cdot c^{(2)}) \), where \( \Delta(c) = c^{(1)} \otimes c^{(2)} \) is the usual coproduct of \( S(\mathfrak{g}_{PT(D)}) \). In particular, if \( t_1, \ldots, t_k, t \in PT(D) \):

\[
t_1 \ldots t_k \cdot t = \sum_{i=1}^{k} t_{1 \ldots i \cdot t} \ldots t_k.
\]

**Lemma 19** Let \( t = (t, I), t_1 = (t_1, I^{(1)}), \ldots, t_k = (t_k, I^{(k)}) \) be partitioned trees \((k \geq 1)\). Let \( s_1, \ldots, s_k \in V(t) \). The partitioned tree \( t \circ s_1, \ldots, s_k \) \((t_1, \ldots, t_k)\) is obtained by grafting the roots of \( t_i \) on \( s_i \) for all \( i \), the partition being \( I \sqcup I^{(1)} \sqcup \ldots \sqcup I^{(k)} \). Then:

\[
t \cdot t_1 \ldots t_k = \sum_{s_1, \ldots, s_k \in V(t)} t \cdot s_1, \ldots, s_k (t_1, \ldots, t_k).
\]

**Proof.** By induction on \( k \). This is obvious if \( k = 1 \). Let us assume the result at rank \( k \).

\[
t \cdot t_1 \ldots t_{k+1} = (t \cdot t_1 \ldots t_k) \cdot t_{k+1} - \sum_{i=1}^{k} t \cdot (t_1 \ldots (t_i \cdot t_{k+1}) \ldots t_k)
\]

\[
= \sum_{s_1, \ldots, s_k \in V(t)} (t \cdot s_1, \ldots, s_k (t_1, \ldots, t_k)) \cdot t_{k+1} - \sum_{i=1}^{k} \sum_{s \in V(t_i)} t \cdot (t_1 \ldots (t_i \cdot s) \cdot t_{k+1}) \ldots t_i
\]

\[
+ \sum_{i=1}^{k} \sum_{s \in V(t_i)} (t \cdot s_1, \ldots, s_k (t_1, \ldots, t_k)) \cdot s \cdot t_{k+1}
\]

\[
\quad - \sum_{i=1}^{k} \sum_{s \in V(t_i)} \sum_{s_1, \ldots, s_k \in V(t) \cdot s \in V(t)} t \cdot s_1, \ldots, s_k (t_1, \ldots, t_i \cdot s) \cdot t_{k+1}, \ldots, t_i
\]

\[
= \sum_{s_1, \ldots, s_k \in V(t)} t \cdot s_1, \ldots, s_k (t_1, \ldots, t_{k+1}).
\]

Hence, the result holds for all \( k \). \( \square \)

**Theorem 20** Let \( D \) be a set, let \( A \) be a Com-Prelie algebra, and let \( a_d \in A \) for all \( d \in D \). There exists a unique morphism \( \phi : \mathfrak{g}_{PT(D)} \to A \), such that \( \phi(a_d) = a_d \) for all \( d \in D \). In other words, \( \mathfrak{g}_{PT(D)} \) is the free Com-Prelie algebra generated by \( D \).

**Proof.** **Unicity.** Let \( t \in T^D \). We denote by \( r_1, \ldots, r_n \) its roots. For all \( 1 \leq i \leq n \), let \( t_{i,1}, \ldots, t_{i,k_i} \) be the partitioned trees born from \( r_i \) and let \( d_i \) be the decoration of \( r_i \). Then:

\[
t = (t_{i,1} \ldots t_{i,k_i}) \cdot \ldots \cdot (t_{n,1} \ldots t_{n,k_n}).
\]

So \( \phi \) is inductively defined by:

\[
\phi(t) = (a_{d_1} \cdot \phi(t_{1,1}) \ldots \phi(t_{1,k_1})) \cdot \ldots \cdot (a_{d_n} \cdot \phi(t_{n,1}) \ldots \phi(t_{n,k_n})).
\]

(1)

**Existence.** As the product \( \cdot \) of \( A \) is commutative and associative, (1) defines inductively a morphism \( \phi \) from \( \mathfrak{g}_{PT(D)} \) to \( A \). By definition, it is compatible with the product \( \cdot \). Let us prove the compatibility with the product \( \cdot \). Let \( t, t' \) be two partitioned trees, let us prove that \( \phi(t \cdot t') = \phi(t) \cdot \phi(t') \) by induction on the number \( N \) of vertices of \( t \). If \( N = 1 \), then \( t = \cdot a_d \) and:

\[
\phi(t \cdot t') = a_d \cdot \phi(t') = \phi(t) \cdot \phi(t'),
\]

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by definition of $t'$. If $N > 1$, two cases are possible. 

First case. If $t$ has only one root, then $t = \cdot \mathcal{d} \cdot t_1 \ldots t_k$, and:

$$t \cdot t' = \cdot \mathcal{d} \cdot t_1 \ldots t_k t' + \sum_{i=1}^{k} \cdot \mathcal{d} \cdot t_1 \ldots t_i t' \cdot t_k.$$ 

Using the induction hypothesis on $t_1, \ldots, t_k$:

$$\phi(t \cdot t') = \phi(t_1 \mathcal{U} t_2 \mathcal{U} \ldots \mathcal{U} t_k)$$

For all $i$, $\mathcal{d} \cdot t_i t' \mathcal{U} \ldots \mathcal{U} t_k = \mathcal{d} \cdot (t_1 \mathcal{U} t_2 \mathcal{U} \ldots \mathcal{U} t_k)$. Hence, $\phi$ is a morphism of Com-Prelie algebras. \hfill \Box

3.2 Presentation of $\mathbb{K}(x_0, x_1)$ as a Com-Prelie algebra

Proposition 21 As a Com-Prelie algebra, $\mathbb{K}(x_0, x_1)$ is generated by $\emptyset$ and $x_1$.

Proof. Let $A$ be the Com-Prelie subalgebra of $\mathbb{K}(x_0, x_1)$ generated by $\emptyset$ and $x_1$. For all $n \geq 1$, it contains $x_1^{\mathcal{U} n} = n! x_1^n$, so it contains $x_1^n$ for all $n \geq 0$. As $\mathbb{K}(x_0, x_1)$ is generated by these elements as a prelie algebra, $A = \mathbb{K}(x_0, x_1)$. \hfill \Box

We denote by $\phi_{CPL} : \mathbb{K}(\{1, 2\}) \rightarrow \mathbb{K}(x_0, x_1)$ the unique morphism of Com-Prelie algebras which sends $\cdot_1$ to $\emptyset$ and $\cdot_2$ to $x_1$. By proposition 21, it is surjective.

Lemma 22 Let $t_1, \ldots, t_k \in \mathcal{P}(\{1, 2\})$.

1. $\phi_{CPL}(\cdot_1 \cdot t_1 \ldots t_k) = 0$ if $k \geq 1$.
2. $\phi_{CPL}(\cdot_2 \cdot t_1 \ldots t_k) = 0$ if $k \geq 2$.
3. If $t \in \mathcal{P}(\{1, 2\})$, $\phi_{CPL}(\cdot_2 \cdot t) = x_0 \phi_{CPL}(t)$.

Proof. We prove 1.-3. by induction on $k$. If $k = 1$:

$$\phi_{CPL}(\cdot_1 \cdot t) = \emptyset \cdot \phi_{CPL}(t) = 0,$$
$$\phi_{CPL}(\cdot_2 \cdot t) = x_1 \cdot \phi_{CPL}(t) = x_0 \phi_{CPL}(t).$$
Let us assume the results at rank $k - 1 \geq 1$. Then:

$$
\phi_{CPL}(\cdot_1 \cdot t_1 \ldots t_k) = \emptyset \cdot \phi_{CPL}(t_1) \ldots \phi_{CPL}(t_k)
$$

$$
= (\emptyset \cdot \phi_{CPL}(t_1) \ldots \phi_{CPL}(t_{k-1})) \cdot \phi_{CPL}(t_k)
$$

$$
- \sum_{i=1}^{k} \emptyset \cdot \phi_{CPL}(t_1) \ldots \phi_{CPL}(t_i \cdot t_k) \ldots \phi_{CPL}(t_{k-1})
$$

$$
= 0,
$$

$$
\phi_{CPL}(\cdot_2 \cdot t_1 \ldots t_k) = x_1 \cdot \phi_{CPL}(t_1) \ldots \phi_{CPL}(t_k)
$$

$$
= (x_1 \cdot \phi_{CPL}(t_1) \ldots \phi_{CPL}(t_{k-1})) \cdot \phi_{CPL}(t_k)
$$

$$
- \sum_{i=1}^{k} x_1 \cdot \phi_{CPL}(t_1) \ldots \phi_{CPL}(t_i \cdot t_k) \ldots \phi_{CPL}(t_{k-1}).
$$

If $k \geq 3$, the induction hypothesis immediately allows to conclude that $\phi_{CPL}(\cdot_2 \cdot t_1 \ldots t_k) = 0 - 0 = 0$. If $k = 2$, this gives:

$$
\phi_{CPL}(\cdot_2 \cdot t_1 t_2) = (x_1 \cdot \phi_{CPL}(t_1)) \cdot \phi_{CPL}(t_2) - x_1 \cdot \phi_{CPL}(t_1 \cdot t_2)
$$

$$
= (x_0 \phi_{CPL}(t_1)) \cdot \phi_{CPL}(t_2) - x_0 \phi_{CPL}(t_1 \cdot t_2)
$$

$$
= x_0 (\phi_{CPL}(t_1) \cdot \phi_{CPL}(t_2)) \phi_{CPL}(t_1 \cdot t_2)
$$

$$
= 0.
$$

Hence, the result holds for all $k \geq 1$. \(\square\)

**Lemma 23** For all $t \in PT(\{1, 2\})$, $\phi_{CPL}(t)$ is a linear span of words of length the number of vertices of $t$ decorated by 2.

**Proof.** By induction on the number of vertices $N$ of $t$. If $N = 1$, then $t = \cdot_1$ or $\cdot_2$ and the result is obvious. Let us assume the result at all rank $< N$.

**First case.** If $t$ has only one root, we put $t = \cdot_1 \cdot t_1 \ldots t_k$. By the preceding lemma, we can assume that $i = 2$ and $k = 1$. Then $\phi_{CPL}(t) = x_0 \phi_{CPL}(t_1)$ and the result is obvious.

**Second case.** If $t$ has $k > 1$ roots, we put $t = t_1 \mathbf{wt} \ldots \mathbf{wt}_k$. Then $\phi_{CPL}(t_1)$ is equal to $\phi_{CPL}(t_1) \mathbf{w} \ldots \mathbf{w} \phi_{CPL}(t_k)$ and the result is immediate. \(\square\)

**Lemma 24** We define inductively a family $F$ of elements of $PT(\{1, 2\})$ by:

1. $F(1) = \{\cdot_1, \cdot_2\}$.
2. $F(n + 1) = (\cdot_2 \cdot F(n)) \cup \bigcup_{i=1}^{n}(F(i) \mathbf{w} F(n + 1 - i))$.
3. $F = \bigcup_{n \geq 1} F(n)$.

Let $t \in PT(\{1, 2\})$. Then $\phi_{CPL}(t) \neq 0$ if, and only if, $t \in F$.

**Proof.** $\implies$. We proceed by induction on the number $N$ of vertices of $t$. This is obvious if $N = 1$. Let us assume the result at all rank $< N$.

**First case.** If $N$ has only one root, we put $N = \cdot_1 \cdot t_1 \ldots t_k$. By lemma 22, $i = 2$ and $k = 1$. Then $\phi_{CPL}(t) = x_0 \phi_{CPL}(t_1)$. By the induction hypothesis, $t_1 \in F$, so $t \in F$. 

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Second case. If $N$ has $k > N$ roots, we put $t = t_1 \mathbb{W} \ldots \mathbb{W} t_k$. Then:

$$\phi_{CPL}(t) = \phi_{CPL}(t_1) \mathbb{W} \phi_{CPL}(t_2 \mathbb{W} \ldots \mathbb{W} t_k) \neq 0,$$

so by the induction hypothesis, $t_1$ and $t_2 \mathbb{W} \ldots \mathbb{W} t_k \in F$, and $t \in F$.

$$\iff$$ Let $t \in T(n)$. We proceed by induction on $n$. If $n = 1$, this is obvious. If $n > 1$ then $t = .2 \cdot t'$, with $t' \in F(n-1)$, or $t = t'' \mathbb{W} t''$, with $t' \in F(i)$, $t'' \in F(n-i)$. In the first case, by the induction hypothesis, $\phi_{CPL}(t') \neq 0$ and $\phi_{CPL}(t) = x_0 \phi_{CPL}(t') \neq 0$. In the second case, $\phi_{CPL}(t'), \phi_{CPL}(t'') \neq 0$ by the induction hypothesis, so $\phi_{CPL}(t) = \phi_{CPL}(t') \mathbb{W} \phi_{CPL}(t'') \neq 0$.  

**Examples.**

- $F(1) = \{1, 2\}$,
- $F(2) = \{1^2, 1^2, 1^2, 2, \ldots \}$,
- $F(3) = \left\{ 1^2, 1^2, \frac{1}{2}v_2, \frac{1}{2}v_2, \frac{1}{2}v_2, \frac{1}{2}v_2, \frac{1}{2}v_2, \frac{1}{2}v_2, \frac{1}{2}v_2, \frac{1}{2}v_2, \frac{1}{2}v_2, \frac{1}{2}v_2, \frac{1}{2}v_2 \right\}$.

We define a second family of elements of $\mathcal{P}T(\{1, 2\})$ in the following way:

1. $F'(1) = \{1, 2\}$.
2. $F'(2) = \{1^2, 1^2, 2, \ldots \}$.
3. $F'(n+1) = (n \cdot F'(n)) \cup \bigcup_{i=2}^{n-1} (F'(i) \mathbb{W} F'(n+1 - i)) \cup (n \mathbb{W} F'(n))$ if $n \geq 2$.
4. $F' = \bigcup_{n \geq 1} F'(n)$.

For example:

- $F'(3) = \left\{ 1^2, 1^2, 2, \ldots \right\}$,
- $F'(4) = \left\{ 1^2, 1^2, 2, \ldots \right\}$.

We define a map $\pi$ from $F$ to $\mathcal{P}T(\{1, 2\})$ in the following way:

1. $\pi(*_i) = *_i$ if $i = 1, 2$.
2. $\pi(*_i \mathbb{W} \ldots \mathbb{W} *_i) = *_i$.
3. If $t = *_i \mathbb{W} \ldots \mathbb{W} _i \mathbb{W} t_1 \mathbb{W} \ldots \mathbb{W} t_k$, $k \geq 1$, with $t_1, \ldots, t_k \neq *_1$, then $\pi(t) = \pi(t_1) \mathbb{W} \ldots \mathbb{W} \pi(t_k)$.
4. If $t = *_i \cdot t_1 \ldots t_k$, then $\pi(t) = *_i \cdot \pi(t_1) \ldots \pi(t_k)$.

**Lemma 25** $\pi$ is a projection on $F'$ and $\phi_{CPL} \circ \pi = \phi_{CPL}|F$.

**Proof.** Let $t \in F$. Let us prove by induction on the number $N$ of vertices of $t$ that:

1. $\pi(t) \in F'$.
2. If $t \in F'$, $\pi(t) = t$.
3. $\phi_{CPL} \circ \pi(t) = \phi_{CPL}(t)$.
4. If $\pi(t) = \cdot_1$, then $t = \cdot_1 w^N$.

All these points are immediate if $N = 1$. Let us assume the result at all ranks $< N$, $N \geq 2$. We put $t = \cdot_1 w \ldots w_{k} t_{1} w \ldots w_{k}$, $k \geq 0$, with $t_{1}, \ldots, t_{k} \neq \cdot_1$.

**First case.** If $k \geq 2$, then $\pi(t) = \pi(t_{1}) w \ldots w \pi(t_{k})$. Following the induction hypothesis, $\pi(t_{1}), \ldots, \pi(t_{k}) \in F'$ and are not equal to $\cdot_1$, so $\pi(t) \in F'$: moreover, $\pi(t_{1}) \neq \cdot_1$, so $\pi(t) \neq \cdot_1$.

$$\phi_{\text{CPL}}(t) = \phi_{\text{CPL}}(\cdot_1) w \ldots w \phi_{\text{CPL}}(\cdot_1) w \ldots w \phi_{\text{CPL}}(t_{k})$$

$$= \emptyset w \ldots w \emptyset w \phi_{\text{CPL}} \circ \pi(t_{1}) w \ldots w \phi_{\text{CPL}} \circ \pi(t_{k})$$

$$= \phi_{\text{CPL}}(\pi(t_{1}) w \ldots w \pi(t_{k}))$$

$$= \phi_{\text{CPL}} \circ \pi(t).$$

If $t \in F'$, necessarily $t = t_{1} w \ldots w t_{k}$, and $t_{1}, \ldots, t_{k} \in F'$. By the induction hypothesis, $\pi(t_{1}) = t_{1}, \ldots, \pi(t_{k}) = t_{k}$, so $\pi(t) = t$.

**Second case.** If $k = 1$, as $t_{1} \in F$, we put $t_{1} = \cdot_2 s$. Then $\pi(t) = \cdot_2 \cdot \pi(s)$. By the induction hypothesis, $\pi(s) \in F'$, so $\pi(t) \in F'$. Moreover:

$$\phi_{\text{CPL}}(t) = \phi_{\text{CPL}}(\cdot_1) w \ldots w \phi_{\text{CPL}}(\cdot_1) w \phi_{\text{CPL}}(\cdot_2) \cdot \phi_{\text{CPL}}(s)$$

$$= \emptyset w \ldots w \emptyset w (\phi_{\text{CPL}}(\cdot_1) w \phi_{\text{CPL}}(\cdot_2) \cdot \phi_{\text{CPL}}(s))$$

$$= \phi_{\text{CPL}} \circ \pi(\cdot_2) \cdot \phi_{\text{CPL}} \circ \pi(s)$$

$$= \phi_{\text{CPL}} \circ \pi(t).$$

If $t' \in F'$, then $s \in F'$, and $t = \cdot_2 s$. Then $\pi(t) = \cdot_2 \cdot \pi(s) = \cdot_2 \cdot s = t$.

**Last case.** If $k = 0$, all the results are obvious. □

**Lemma 26** Let $t, t' \in PT([1, 2])$. Then:

$$\phi_{\text{CPL}}\left((\cdot_2 \cdot t) w (\cdot_2 \cdot t')\right) = \phi_{\text{CPL}}\left(\cdot_2 \cdot ((\cdot_2 \cdot t) w t' + t w (\cdot_2 \cdot t'))\right).$$

**Proof.** Indeed, putting $w = \phi_{\text{CPL}}(t)$ and $w' = \phi_{\text{CPL}}(t')$:

$$\phi_{\text{CPL}}\left((\cdot_2 \cdot t) w (\cdot_2 \cdot t')\right) = x_{0} w w x_{0} w'$$

$$= x_{0} (w x_{0} w') + x_{0} (x_{0} w w')$$

$$= \phi_{\text{CPL}}\left(\cdot_2 \cdot ((\cdot_2 \cdot t) w t' + t w (\cdot_2 \cdot t'))\right).$$

We used lemma 22 for the first and third equalities. □

**Theorem 27** The kernel of $\phi_{\text{CPL}}$ is the Com-Prelie ideal generated by the elements:

1. $\cdot_1 \cdot t_{1} \ldots t_{k}$, where $k \geq 1$, $t_{1}, \ldots, t_{k} \in PT([1, 2])$.

2. $\cdot_2 \cdot t_{1} \ldots t_{k}$, where $k \geq 2$, $t_{1}, \ldots, t_{k} \in PT([1, 2])$.

3. $\cdot_1 w t - t$, where $t \in PT([1, 2])$.

4. $(\cdot_2 \cdot t) w (\cdot_2 \cdot t') - \cdot_2 \cdot ((\cdot_2 \cdot t) w t' - t w (\cdot_2 \cdot t'))$, where $t, t' \in PT([1, 2])$.

**Proof.** Let $I$ be the ideal generated by these elements. Lemmas 22 and 26 prove that the elements 1, 2 and 4 belong to $\text{Ker}(\phi_{\text{CPL}})$. Moreover, for all $t \in PT([1, 2])$, $\pi(\cdot_1 w t) = \pi(t)$. For all $t \in PT([1, 2])$:

$$\phi_{\text{CPL}}(\cdot_1 w t) = \emptyset w \phi_{\text{CPL}}(t) = \phi_{\text{CPL}}(t),$$

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so elements 3. also belong to \( \text{Ker}(\phi_{CPL}) \). Hence, \( I \subseteq \text{Ker}(\phi_{CPL}) \).

Let \( h = g_{\mathcal{P}T(\{1,2\})}/I \). As the elements 1 and 2 belong to \( I \), \( h \) is linearly spanned by the elements \( \overline{7}, t \in F' \). As the elements 3 belong to \( I \), for all \( t \in F \), \( \overline{\pi(t)} = \overline{7} \). As \( \pi \) is a projection on \( F' \), \( h \) is linearly spanned by the elements \( \overline{t}, t \in F' \).

We now define inductively two families of partitionned trees in the following way:

1. \( T''(1) = \{ \cdot_2 \} \) and \( F''(1) = \{ \cdot_1, \cdot_2 \} \).
2. \( T''(n + 1) = \cdot_2 \bullet F''(n) \).
3. \( F''(n + 1) = \bigcup_{i=1}^{n+1} T''(i) \bigcup_{i=1}^{n+1} \mathbb{W}(n+1-i) \).
4. \( F'' = \bigcup_{n \geq 1} F''(n) \).

For example:

\[
F''(3) = \left\{ I_2^1, I_2^2, I_2^3, 2I_2^2, 2I_2^1, 2I_2^3, 2I_2^2 \right\},
\]

\[
F''(4) = \left\{ I_2^1, I_2^2, I_2^3, 2I_2^2, 2I_2^1, 2I_2^3, 2I_2^2, 2I_2^1, 2I_2^3, 2I_2^2 \right\}.
\]

Let us prove that for all \( t \in F'' \), there exists \( t' \in \text{Vect}(F'') \) such that \( \overline{t} = \overline{t'} \). We proceed by induction on the number \( N \) of vertices of \( t \). If \( N = 1 \), then \( t = \cdot_1 \) or \( \cdot_2 \) and we take \( t' = t \). Let us assume the result at all rank \( \leq N - 1 \). We put \( t = t_1 \mathbb{W} \ldots \mathbb{W} t_k \mathbb{W} t_k \mathbb{W} \ldots \mathbb{W} t_2 \mathbb{W} \ldots \mathbb{W} t_1 \mathbb{W} \), with \( t_i = \cdot_2 \bullet s_i, s_i \neq 1 \), for all \( 1 \leq i \leq k \). We proceed by induction on \( k \). If \( k = 0 \), we take \( t' = t = \cdot_2 \mathbb{W} \ldots \mathbb{W} t_2 \mathbb{W} \).

If \( k = 1 \), then, by the induction hypothesis on \( N \) applied to \( s_1 \):

\[
\overline{t} = (\cdot_2 \bullet \overline{s_1}) \mathbb{W} \ldots \mathbb{W} \mathbb{W} = (\cdot_2 \bullet \overline{s_1'}) \mathbb{W} \ldots \mathbb{W} \mathbb{W} = (\cdot_2 \bullet \overline{s_1'}) \mathbb{W} \ldots \mathbb{W} \mathbb{W}.
\]

We take \( t' = (\cdot_2 \bullet s_1') \mathbb{W} \ldots \mathbb{W} t_2 \mathbb{W} \), which clearly belongs to \( \text{Vect}(F'') \), as \( s_1' \in \text{Vect}(F'') \). Let us assume the result at all rank \( \leq k - 1 \). Then, as the elements 4 belong to \( I \):

\[
t_1 \mathbb{W} t_2 = \cdot_2 \bullet (t_1 \mathbb{W} s_2) + \cdot_2 \bullet (s_1 \bullet t_2),
\]

so:

\[
\overline{t} = \overline{t_1} \mathbb{W} t_3 \mathbb{W} \ldots \mathbb{W} t_k \mathbb{W} t_2 \mathbb{W} \ldots \mathbb{W} t_2 \mathbb{W} + \overline{t_1} \mathbb{W} t_3 \mathbb{W} \ldots \mathbb{W} t_k \mathbb{W} \mathbb{W} \mathbb{W} t_2 \mathbb{W} \ldots \mathbb{W} t_2 \mathbb{W}.
\]

By the induction hypothesis on \( k \) applied to these two partitionned trees, there exists \( x_1' \) and \( x_2' \) in \( \text{Vect}(F'') \), such that \( \overline{t} = \overline{x_1'} + \overline{x_2'} \). We take \( t' = x_1' + x_2' \). Consequently, the elements \( \overline{t}, t \in F'' \), linearly span \( h \).

Let \( t \in F''(n) \). Then it has \( n \) vertices, and at most one of them is decorated by 1. We denote by \( F''(n) \) the set of elements of \( F''(n) \) with one vertex decorated by 1, and we put \( F_2''(n) = F''(n) \setminus F''(n) \). Let us prove that for all \( n \geq 1 \), \( |F''(n + 1)| \leq 2^{n-1} \) and \( |F_2''(n)| \leq 2^{n-1} \). For \( n = 0 \), as \( FF_1''(2) = \{ I_2^1 \} \) and \( F_2''(1) = \{ \cdot_2 \} \), this is immediate. Let us assume the result at all rank \( \leq n \). Then:

\[
F''(n + 1) = \bigcup_{i=1}^{n+1} \cdot_2 \mathbb{W}^{n+1-i} \bigcup_{i=1}^{n+1} \bigcup_{i=1}^{n} \cdot_2 \mathbb{W}^{n+1-i} \mathbb{W} t''(i) \mathbb{W} \bigcup_{i=1}^{n} \cdot_2 \mathbb{W}^{n+1-i} \mathbb{W} t''(i) \mathbb{W} t''(i).
\]
Hence, \(|F_2''(n + 1)| \leq 1 + 1 + 2 + \ldots + 2^{n-1} = 2^n\).

\[ F_1''(n + 2) = \bigcup_{i=1}^{n+2} \mathcal{U} \mathcal{T}''(i) \cap F_1''(i) = \bigcup_{i=2}^{n+2} \mathcal{U} \mathcal{T}''(i-1). \]

Hence, \(|F_1''(n + 2)| \leq +1 + 1 + \ldots + 2^{n-1} = 2^n\).

Let \(\overline{\phi}_{APL}\) be the linear map induced by \(\phi_{CPL}\) on \(h\). If \(t \in F_2''(n)\), by lemma 23, \(\overline{\phi}_{APL}(t)\) is a linear span of words of length \(n - 1\). If \(t \in F_1''(n)\), by lemma 23, \(\overline{\phi}_{APL}(t)\) is a linear span of words of length \(n\). Hence, for all \(n \geq 0\):

\[ \overline{\phi}_{APL}(\text{Vect}(F_2''(n)) + \text{Vect}(F_1''(n + 1))) \subseteq \text{Vect}(\text{words of length } n). \]

As \(\phi_{CPL}\) is surjective, we obtain:

\[ \overline{\phi}_{APL}(\text{Vect}(F_2''(n)) + \text{Vect}(F_1''(n + 1))) = \text{Vect}(\text{words of length } n). \]

Moreover, as \(\dim(\text{Vect}(\text{words of length } n)) = 2^n\) and \(\dim(\text{Vect}(F_2''(n)) + \text{Vect}(F_1''(n + 1))) \leq |F_2''(n)| + |F_1''(n + 1)| \leq 2^{n-1} + 2^{n-1} = 2^n\), the restriction of \(\overline{\phi}_{APL}\) to \(\text{Vect}(F_2''(n)) + \text{Vect}(F_1''(n + 1))\) is injective. Finally, \(\overline{\phi}_{APL}\) is injective, so \(\text{Ker}(\phi_{CPL}) = I\). \(\square\)

4 Presentation of \(K\langle x_0, x_1 \rangle\) as a prelie algebra

4.1 A surjective morphism

Let \(g_{T(N^*)}\) be the free prelie algebra generated by \(N^*\), as described in [1]. It can be seen as the subspace of \(g_{PT(N^*)}\) generated by rooted trees (which are seen as partitioned trees such that any part of the partition is a singleton), with the restriction of the prelie product \(\bullet\) defined by graftings. For example, in \(g_{T(N^*)}\), if \(a, b, c, d > 0\):

\[ 1^\bullet_a \cdot 1^d_b = 1^\bullet_{a+b} + 1^d_{b}. \]

This prelie algebra is graded, the degree of a tree being the sum of its decorations.

By theorem 12, there exists a unique surjective map of prelie algebras \(\Phi_{PL} : g_{T(N^*)} \rightarrow K\langle x_0, x_1 \rangle\), sending \(\cdot_n\) to \(x_1^{n-1}\) for all \(n \geq 1\). As \(x_1^{n-1}\) is homogeneous of degree \(n\) for all \(n\), this morphism is homogeneous of degree 0.

**Notation.** If \(t_1 \ldots t_k \in T(N^*)\) and \(n \in N^*\), we put:

\[ B_n(t_1 \ldots t_k) = \cdot_n \bullet t_1 \ldots t_k. \]

This is the tree obtained by grafting \(t_1, \ldots, t_k\) on a common root decorated by \(n\).

**Proposition 28** Let \(t = B_n(t_1 \ldots t_k) \in T(N^*)\). We put \(\phi_{PL}(t_i) = w_i\) for all \(1 \leq i \leq k\). Then:

\[ \phi_{PL}(t) = \begin{cases} x_0 w_1 \mathcal{U} \ldots x_0 w_k \mathcal{U} x_1^{n-1-k} & \text{if } k < n, \\ 0 & \text{otherwise}. \end{cases} \]

**Proof.** As \(g_{PT((1,2))}\) is prelie, there exists a unique morphism of prelie algebras:

\[ \psi : \begin{cases} g_{T(N^*)} & \rightarrow g_{PT((1,2))} \\ \cdot_n & \rightarrow (n-1)! \cdot 2 \mathcal{U} (n-1). \end{cases} \]

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Then $\phi_{APL} \circ \psi$ is a prelie algebra morphism sending $\star_n$ to $\frac{1}{(n-1)!}x_1^{w(n-1)} = x_1^{n-1}$ for all $n \geq 1$, so $\phi_{APL} \circ \psi = \phi_{PL}$. We obtain, by lemma 19:

$$\psi(\star_n \cdot t_1 \ldots t_k) = \frac{1}{(n-1)!} \sum_{I_1 \cup \ldots \cup I_{n-1} = \{1, \ldots, k\}} \sum_{i \in I_1} x_1 \cdot \left( \prod_{i \in I_1} w_i \right) \cdot \left( \prod_{i \in I_{n-1}} t_i \right)$$

Let us apply $\phi_{APL}$ to this expression. If $|I_j| \geq 2$, by theorem 27:

$$\phi_{APL} \left( \star_n \cdot \left( \prod_{i \in I_1} t_i \right) \right) = 0.$$ 

Consequently, if $k \geq n$, at least one of the $I_j$ contains two elements, so $\phi_{APL} \circ \psi(t) = \phi_{PL}(t) = 0$. Let us assume that $k < n$. Hence, using the commutativity of $\mathbf{w}$:

$$\phi_{PL}(\star_n \cdot t_1 \ldots t_k) = \frac{1}{(n-1)!} \sum_{I_1 \cup \ldots \cup I_{n-1} = \{1, \ldots, k\}} x_1 \cdot \left( \prod_{i \in I_1} w_i \right) \cdot \left( \prod_{i \in I_{n-1}} t_i \right)$$

which is the announced result. □

**Corollary 29** Let $s_1, \ldots, s_k, t_1, \ldots, t_l \in T(\{N^*\}, k, l \geq 0$. For all $i,j,n \geq 1$:

$$\phi_{PL}(B_{n+1}((B_i(s_1 \ldots s_k)B_j(t_1 \ldots t_l))) = \phi_{PL}(B_n(B_{i+1}(s_1 \ldots s_k)B_j(t_1 \ldots t_l))) + \phi_{PL}(B_n(B_{j+1}(s_1 \ldots s_k)t_1 \ldots t_l)).$$

**Proof.** We note:

- $T_1 = B_{n+1}((B_i(s_1 \ldots s_k)B_j(t_1 \ldots t_l))) = \star_{n+1} \cdot ((\star \cdot s_1 \ldots s_k)(\star \cdot t_1 \ldots t_l))$,
- $T_2 = B_n(B_{i+1}(s_1 \ldots s_k)B_j(t_1 \ldots t_l)) = \star_n \cdot ((\star \cdot s_1 \ldots s_k)(\star \cdot t_1 \ldots t_l))$,
- $T_3 = B_n(B_{j+1}(s_1 \ldots s_k)t_1 \ldots t_l) = \star_n \cdot ((\star \cdot s_1 \ldots s_k)(\star \cdot t_1 \ldots t_l))$.

If $k \geq i$, or $l \geq j$, or $n = 1$, all these elements are sent to zero by $\phi_{PL}$ by proposition 28. Let us assume now that $k < i, l < j, n < 1$. We put $v_i = \phi_{PL}(s_i)$ and $w_l = \phi_{PL}(t_l)$. Then:

$$\phi_{PL}(T_1) = x_0(x_0v_1 \mathbf{w} \ldots \mathbf{w} x_0v_k \mathbf{w} x_1^{i-1-k}) \mathbf{w} x_0(x_0w_1 \mathbf{w} \ldots \mathbf{w} x_0w_l \mathbf{w} x_1^{1-l}) \mathbf{w} x_1^{n-2}$$

$$= x_0X \mathbf{w} x_0Y \mathbf{w} x_1^{n-2},$$

$$\phi_{PL}(T_2) = x_0(x_0v_1 \mathbf{w} \ldots \mathbf{w} x_0w_1 \mathbf{w} \ldots \mathbf{w} x_0w_l \mathbf{w} x_1^{1-l} \mathbf{w} x_1^{1-k}) \mathbf{w} x_1^{n-2}$$

$$= x_0(X \mathbf{w} x_0Y) \mathbf{w} x_1^{n-2},$$

$$\phi_{PL}(T_3) = x_0(x_0x_1 \mathbf{w} \ldots \mathbf{w} x_0w_k \mathbf{w} x_1^{i-1-k} \mathbf{w} x_0w_1 \mathbf{w} x_0w_l \mathbf{w} x_1^{1-l} \mathbf{w} x_1^{1-k}) \mathbf{w} x_1^{n-2}$$

$$= x_0(x_0X \mathbf{w} x_0Y) \mathbf{w} x_1^{n-2}.$$

As $x_0X \mathbf{w} x_0Y = x_0(X \mathbf{w} x_0Y) + x_0(x_0X \mathbf{w} x_0Y)$, we obtain the result. □
Theorem 30 The kernel of $\phi_{PL}$ is the prelie ideal generated by:

1. $B_1(t_1 \ldots t_k)$, where $k \geq 1$, $t_1, \ldots, t_k \in \mathcal{T}(\mathbb{N}^*)$.

2. $B_{n+1}(B_l(s_1 \ldots s_k)B_j(t_1 \ldots t_l)) - B_n(B_{l+1}(s_1 \ldots s_k)B_j(t_1 \ldots t_l)) - B_{j+1}(B_l(s_1 \ldots s_k)t_1 \ldots t_l))$, where $k, l \geq 0$, $s_1, \ldots, s_k, t_1, \ldots, t_l \in \mathcal{T}(\mathbb{N}^*)$.

Proof. Let $I$ be the ideal generated by these elements. By proposition 28 and corollary 29, $I \subseteq Ker(\phi_{PL})$. We put $h = g_{\mathcal{T}(\mathbb{N}^*)}/I$. Applying repeatedly the relation given by elements of the second form, it is not difficult to prove that for any $t \in \mathcal{T}(\mathbb{N}^*)$, there exists a linear span of ladders $t'$ such that $\overrightarrow{t} = \overrightarrow{t'}$ in $h$. Moreover, by the relation given by elements 1., if one of the vertices of a ladder $t$ which is not the leaf is decorated by 1, then $\overrightarrow{t} = 0$. Let us denote by $L(n)$ the set of ladders decorated by $\mathbb{N}^*$, of weight $n$, such that all the vertices which are not the leaf are decorated by integers $> 1$. It turns out that $h$ is generated by the elements $\overrightarrow{t}$, $t \in L = \bigcup L(n)$.

Let $\phi_{PL}$ be the morphism form $h$ to $\mathbb{K}\langle x_0, x_1 \rangle$ induced by $\phi_{PL}$. By homogeneity, as $\phi_{PL}$ is surjective, for all $n \geq 1$:

$$\overrightarrow{\phi_{PL}}(\text{Vect}(L(n))) = \text{Vect}(\text{words of degree } n).$$

In order to prove that $I = Ker(\phi_{PL})$, it is enough to prove that $\overrightarrow{\phi_{PL}}$ is injective. By homogeneity, it is enough to prove that $\overrightarrow{\phi}_{|\text{Vect}(L(n))}$ is injective for all $n \geq 1$. Hence, it is enough to prove that for all $n \geq 1$,

$$|L(n)| = \text{dim}(\text{Vect}(\text{words of degree } n)) = p_n,$$

where the $p_n$ are the integers defined in proposition 8. Let $l_n = |L(n)|$ and $q_n$ be the number of $t \in L(n)$ with no vertex decorated by 1. Then for all $n \geq 2$, $l_n = q_n + q_{n-1}$, and $l_1 = 1$. We put:

$$L = \sum_{n=1}^\infty l_n x^n, \quad Q = \sum_{n=1}^\infty q_n x^n.$$

We obtain $P = X + Q + XQ$. Moreover:

$$Q = \frac{1}{1 - \sum_{i \geq 2} X^i} - 1 = \frac{1}{1 - \frac{X^2}{1 - X}} - 1 = \frac{X^2}{1 - X - X^2}.$$ 

Finally:

$$L = \frac{X}{1 - X - X^2} = F.$$ 

So, for all $n \geq 1$, $|L(n)| = p_n$. $\square$

As an immediate corollary, a basis of $h$ is given by the classes of the elements of $L$. Turning to $\mathbb{K}\langle x_0, x_1 \rangle$, we obtain:

Corollary 31 Let $w = a_1 \ldots a_k$ be a word with letters in $\mathbb{N}^*$.

1. We put:

$$m_w = x_1^{a_1-1} \cdot (x_1^{a_1-1} \cdot (\ldots (x_1^{a_k-1} \cdot x_1^{a_k}) \ldots)).$$

2. We shall say that $w$ is admissible if $a_1, \ldots, a_{k-1} > 1$. The set of admissible words is denoted by $\text{Adm}$. Then $(m_w)_{w \in \text{Adm}}$ is a basis of $\mathbb{K}\langle x_0, x_1 \rangle$.

Remark. If $w$ is not admissible, that is to say if there exists $1 \leq i < k$, such that $a_i = 1$, then $m_w = 0$ by proposition 28.

We extend the map $w \mapsto m_w$ by linearity.
4.2 Prelie product in the basis of admissible words

Notations.

1. For all $k,l$, we denote by $\text{Sh}(k,l)$ the set of $(k,l)$-shuffles, that is to say permutations $\zeta \in \mathfrak{S}_{k+l}$ such that $\zeta(1) < \ldots < \zeta(k)$, $\zeta(k+1) < \ldots < \zeta(k+l)$.

2. For all $k,l$ we denote by $\text{Sh}_{\prec}(k,l)$ the set of $(k,l)$-shuffles $\zeta$ such that $\zeta^{-1}(k+l) = k$.

3. For all $k,l$ we denote by $\text{Sh}_{\succ}(k,l)$ the set of $(k,l)$-shuffles $\zeta$ such that $\zeta^{-1}(k+l) = k+l$.

4. The symmetric group $\mathfrak{S}_n$ acts on the set of words with letters in $\mathbb{N}^*$ of length $n$ by permutation of the letters:
   $$\sigma.(a_1 \ldots a_n) = a_{\sigma^{-1}(1)} \ldots a_{\sigma^{-1}(n)}.$$

Proposition 32 Let $\mathbb{K} \langle \mathbb{N}^* \rangle$ be the space generated by words with letters in $\mathbb{N}^*$. We define a dendriform structure on this space by:

\[
(a_1 \ldots a_k) \prec (b_1 \ldots b_l) = \sum_{\zeta \in \text{Sh}_{\prec}(k,l)} \zeta.a_1 \ldots a_kb_1 \ldots b_{k-1}(b_k+1)
\]

\[
(a_1 \ldots a_k) \succ (b_1 \ldots b_l) = \sum_{\zeta \in \text{Sh}_{\succ}(k,l)} \zeta.a_1 \ldots a_{k-1}(a_k+1)b_1 \ldots b_l.
\]

The associative product $\prec + \succ$ is denoted by $\star$.

Proof. We denote by $\text{Sh}(k,l,m)$ the set of $k+l+m$-permutations such that $\zeta(1) < \ldots < \zeta(k)$, $\zeta(k+1) < \ldots < \zeta(k+l)$, $\zeta(k+l+1) < \ldots < \zeta(k+l+m)$. Then:

\[
(a_1 \ldots a_k \prec b_1 \ldots b_l) \prec c_1 \ldots c_m = a_1 \ldots a_k \prec (b_1 \ldots b_l \star c_1 \ldots c_m)
\]

\[
= \sum_{\zeta \in \text{Sh}(k,l,m); \zeta^{-1}(k+l+m) = k} \zeta.a_1 \ldots a_kb_1 \ldots (b_l+1)c_1 \ldots (c_m+1);
\]

\[
(a_1 \ldots a_k \succ b_1 \ldots b_l) \succ c_1 \ldots c_m = a_1 \ldots a_k \succ (b_1 \ldots b_l \succ c_1 \ldots c_m)
\]

\[
= \sum_{\zeta \in \text{Sh}(k,l,m); \zeta^{-1}(k+l+m) = k+l} \zeta.a_1 \ldots (a_k+1)b_1 \ldots b_lc_1 \ldots (c_m+1);
\]

\[
(a_1 \ldots a_k \star b_1 \ldots b_l) \succ c_1 \ldots c_m = a_1 \ldots a_k \succ (b_1 \ldots b_l \succ c_1 \ldots c_m)
\]

\[
= \sum_{\zeta \in \text{Sh}(k,l,m); \zeta^{-1}(k+l+m) = k+l+m} \zeta.a_1 \ldots (a_k+1)b_1 \ldots (b_l+1)c_1 \ldots c_m.
\]

So $\mathbb{K} \langle \langle \mathbb{N}^* \rangle \rangle$ is a dendriform algebra. $\square$

We postpone the study of this dendriform algebra to section 5.2.

Notations. For all $a_1, \ldots, a_k \in \mathbb{N}^*$, we denote by $l(a_1 \ldots a_k) = B_{a_1} \circ \ldots \circ B_{a_k}(1)$ the ladder decorated from the root to the leaf by $a_1, \ldots, a_k$. Note that $m_{a_1 \ldots a_k} = \phi_{PL}(l(a_1 \ldots a_k))$.

Lemma 33 Let $k,l \geq 1$ and let $a_1, \ldots, a_l, b_1, \ldots, b_l \in \mathbb{N}^*$. Then:

$$\phi_{PL}(B_{a_1+1}(l(a_2 \ldots a_k)l(b_1 \ldots b_l)) + B_{b_1+1}(l(a_1 \ldots a_k)l(b_2 \ldots b_l)) = m_{a_1 \ldots a_k \star b_1 \ldots b_l}$$

Proof. By induction on $k+l$. If $k = l = 1$, then:

$$\phi_{PL}(\mathbf{1}_{a_1}^{b_1} + \mathbf{1}_{b_1}^{a_1}) = m_{(a_1+1)b_1+(b_1+1)a_1} = m_{a_1 \star b_1}.$$
Let us assume the result at all ranks $< k + l$. If $k = 1$, then:

$$
\phi_{PL}(B_{a_1+1}(l(b_2 \ldots b_l)) + B_{b_1+1}(l(a_1)l(b_2 \ldots b_l))
= \phi_{PL}(\ast_{a_1+1} \cdot l(b_2 \ldots b_l) + \ast_{b_1+1} \cdot (l(a_1)l(b_2 \ldots b_l)))
= \phi_{PL}(l((a_1+1)b_2 \ldots b_l)) + \phi_{PL}(\ast_{b_1} \cdot (l((a_1+1)b_2 \ldots b_l) + \ast_{b_2+1} \cdot (l(a_1)l(b_3 \ldots b_l)))
= m_{(a_1+1)b_2 \ldots b_l} + m_{b_1(a_1) \ast b_2 \ldots b_l}
= m_{a_1b_1 \ldots b_l}.
$$

If $l = 1$, a similar computation, permuting the $a_i$'s and the $b_j$'s, proves the result. If $k, l > 1$, then:

$$
\phi_{PL}(B_{a_1+1}(l(a_2 \ldots a_k))(l(b_1 \ldots b_l)) + B_{b_1+1}(l(a_1 \ldots a_k)l(b_2 \ldots b_l))
= \phi_{PL}(\ast_{a_1} \cdot (l(a_2 \ldots a_k))(l(b_1 \ldots b_l)) + \ast_{b_1+1} \cdot (l(a_1 \ldots a_k)l(b_2 \ldots b_l)))
+ \phi_{PL}(\ast_{a_1} \cdot (l(a_2 \ldots a_k))(l(b_2 \ldots b_l)) + \ast_{b_2+1} \cdot (l(a_1 \ldots a_k)l(b_3 \ldots b_l)))
= m_{a_1(a_2 \ldots a_k \ast b_1 \ldots b_l)+b_1(a_1 \ldots a_k \ast b_2 \ldots b_l)}
= m_{a_1 \ast a_k \ast b_1 \ldots b_l}.
$$

Hence, the result holds for all $k, l \geq 1$. □

**Theorem 34** For all $a_1, \ldots, a_k, b_1, \ldots, b_l \in \mathbb{N}^*$:

$$
m_{a_1 \ldots a_k} \cdot m_{b_1 \ldots b_l} = \sum_{i=1}^{k-1} m_{a_1 \ldots a_{i-1}}(a_i a_{i+1} \ldots a_k \ast b_1 \ldots b_l) + m_{a_1 \ldots a_k b_1 \ldots b_l}.
$$

**Proof.** By definition of $m_{a_1 b_1 \ldots b_l}$, if $k = 1$, $m_{a_1} \cdot m_{b_1 \ldots b_l} = m_{a_1 b_1 \ldots b_l}$. So the result holds if $k = 1$. Let us assume that $k \geq 2$. In $\mathfrak{g}_T(\mathbb{N}^*)$, we have:

$$
l(a_1 \ldots a_k) \cdot l(b_1 \ldots b_l) = \ast_{a_1} \cdot (l(a_2 \ldots a_k) \cdot l(b_1 \ldots b_l)) + \ast_{a_2} \cdot l(a_2 \ldots a_k)l(b_1 \ldots b_l).
$$

Applying $\phi_{PL}$:

$$
m_{a_1 \ldots a_k} \cdot m_{b_1 \ldots b_l} = m_{a_1(a_2 \ldots a_k) \ast b_1 \ldots b_l}
+ \phi_{PL}(\ast_{a_1} \cdot (l(a_2 \ldots a_k)l(b_1 \ldots b_l)) + \ast_{a_2} \cdot (l(a_1 \ldots a_k)l(b_2 \ldots b_l)))
= m_{a_1(a_2 \ldots a_k) \ast b_1 \ldots b_l} + m_{(a_1-1)(a_2 \ldots a_k \ast b_1 \ldots b_l)},
$$

by the preceding lemma. The result follows from an easy induction. □

**Remark.** In particular, $m_1 \circ m_{b_1 \ldots b_l} = 0$.

**Corollary 35** Let $a_1 \ldots a_k, b_1 \ldots b_l$ be two words with letters in $\mathbb{N}^*$. Then $m_{a_1 \ldots a_k} \cdot m_{b_1 \ldots b_l}$ is a span of $m_w$, where $w$ is a word with $k + l$ letters and of weight $a_1 + \ldots + a_k + b_1 + \ldots + b_l$.

Hence, $\mathbb{K}(x_0, x_1)$ is a bigraded prelie algebra, with:

$$
\mathbb{K}(x_0, x_1)_{n,k} = Vect(m_{a_1 \ldots a_k} \mid a_1 + \ldots + a_k = n).
$$

We put:

$$
G = \sum_{k,n \geq 0} \text{dim}(\mathbb{K}(x_0, x_1)_{n,k})X^nY^k.
$$
Proposition 36 \[ G = \frac{XY}{1 - X - X^2Y} = \sum_{k=1}^{\infty} \sum_{l=2k-1}^{\infty} \binom{l-k}{k-1} X^l Y^k. \]

**Proof.** Note that \( \dim(\mathbb{K}(x_0, x_1)_{n,k}) \) is the number of words \( a_1 \ldots a_k \) of length \( k \), such that \( a_1, \ldots, a_{k-1} \geq 2 \), and \( a_1 + \ldots + a_k = n \). Hence:

\[ G = \sum_{k=1}^{\infty} \left( \frac{X^2Y}{1 - X} \right)^{k-1} \frac{XY}{1 - X} = \frac{XY}{1 - X - X^2Y}. \]

An easy development in formal series gives the second formula. \( \square \)

4.3 An associative product on \( g_T(\mathbb{N}^*) \)

We now define an associative product on \( g_T(\mathbb{N}^*) \), in such a way that \( \phi_{PL} \) becomes a morphism of Com-Prelie algebras.

**Proposition 37** We define a product \( \square \) on \( g_T(\mathbb{N}^*) \) by:

\[ B_p(s_1 \ldots s_k) \square B_q(t_1 \ldots t_l) = \binom{p+q-k-l-2}{p-k-1} B_{p+q-1}(s_1 \ldots s_k t_1 \ldots t_l). \]

Then \( g_T(\mathbb{N}^*) \) is a Com-Prelie algebra and \( \phi_{PL} \) is a morphism of Com-Prelie algebras.

**Proof.** As \( \binom{p+q-k-l-2}{p-k-1} = \binom{p+q-k-l-2}{q-l-1} \), \( \square \) is commutative. Let \( t = B_p(s_1 \ldots s_k) \), \( t' = B_q(\bullet t_1 \ldots t_l) \) and \( t'' = B_r(u_1 \ldots u_m) \). Then:

\[ t \square (t' \square t'') = \binom{p+q+r-k-l-m-2}{q-l-1} \binom{p+q+r-k-l-m-3}{q+r-l-m-2} B_{p+q+r-2}(s_1 \ldots s_k t_1 \ldots t_l u_1 \ldots u_m), \]

\[ (t \square t') \square t'' = \binom{p+q-k-l-2}{p-k-1} \binom{p+q+r-k-l-m-3}{p+q-k-l-2} B_{p+q+r-2}(s_1 \ldots s_k t_1 \ldots t_l u_1 \ldots u_m). \]

If \( p \leq k \) or \( q \leq l \) or \( r \leq m \), then \( A = B = 0 \). If \( p > k \) and \( q > l \) and \( r > m \), then:

\[ A = B = \frac{(p+q+r-k-l-m-3)!}{(p-k-1)!(q-l-1)!(r-m-1)!}. \]

So \( \square \) is associative.
Let \( t_1 = B_p(s_1 \ldots s_k) \), \( t_2 = B_q(t_1 \ldots t_l) \) and \( t \in \mathcal{T}(\mathbb{N}^*) \). Then:

\[
(t_1 \mathbf{w} t_2) \circ T = \left(\frac{p+q-k-l-2}{m-k-1}\right) B_{p+q-1}(s_1 \ldots s_k t_1 \ldots t_l) + \sum_{i=1}^{k} \left(\frac{p+q-k-l-2}{p-k-1}\right) B_{p+q-1}(s_1 \ldots (s_i \# t) \ldots s_k t_1 \ldots t_l) + \sum_{j=1}^{l} \left(\frac{p+q-k-l-2}{p-k-1}\right) B_{p+q-1}(s_1 \ldots s_k t_1 \ldots (t_j \# t) \ldots t_l),
\]

\[
(t_1 \# t) \mathbf{w} t_2 = \left(\sum_{i=1}^{k} B_p(s_1 \ldots (s_i \# t) \ldots s_k) + B_p(s_1 \ldots s_k t)\right) \mathbf{w} t_2
\]

\[
= \sum_{i=1}^{k} \left(\frac{p+q-k-l-2}{p-k-1}\right) B_{p+q-1}(s_1 \ldots (s_i \# t) \ldots s_k t_1 \ldots t_l) + \left(\frac{p+q-k-l-3}{p-k-2}\right) B_{p+q-1}(s_1 \ldots s_k t_1 \ldots t_l),
\]

\[
t_1 \mathbf{w} (t_2 \# t) = t_1 \mathbf{w} \left(\sum_{j=1}^{l} B_q(t_1 \ldots (t_j \# t) \ldots t_l) + B_q(t_1 \ldots t_j t)\right)
\]

\[
= \sum_{j=1}^{l} \left(\frac{p+q-k-l-2}{p-k-1}\right) B_{p+q-1}(s_1 \ldots s_k t_1 \ldots (t_j \# t) \ldots t_l) + \left(\frac{p+q-k-l-3}{p-k-2}\right) B_{p+q-1}(s_1 \ldots s_k t_1 \ldots t_l).
\]

As \( \left(\frac{p+q-k-l-3}{p-k-2}\right) + \left(\frac{p+q-k-l-3}{p-k-1}\right) = \left(\frac{p+q-k-l-2}{p-k-1}\right) \), we obtain \( (t_1 \mathbf{w} t_2) \# t = (t_1 \# t) \mathbf{w} t_2 + t_1 \mathbf{w} (t_2 \# t) \). So \( \mathcal{G}_{\mathcal{T}(\mathbb{N}^*)} \) is Com-Prelie.

Let \( t_1 = B_p(s_1 \ldots s_k) \) and \( t_2 = B_q(t_1 \ldots t_l) \). If \( k \geq p \), then \( \left(\frac{p+q-k-l-2}{p-k-1}\right) = 0 \), so \( t_1 \mathbf{w} t_2 = 0 \). By proposition 28, \( \phi_{PL}(t_1) = 0 \), so \( \phi_{PL}(t_1 \mathbf{w} t_2) = \phi_{PL}(t_1) \mathbf{w} \phi_{PL}(t_2) = 0 \). Similarly, if \( l \geq q \), \( \phi_{PL}(t_1 \mathbf{w} t_2) = \phi_{PL}(t_1) \mathbf{w} \phi_{PL}(t_2) = 0 \). If \( k < p \) and \( l < q \), we put \( w_i = \phi_{PL}(s_i) \) and \( w'_j = \phi_{PL}(t_j) \). Then:

\[
\phi_{PL}(t_1) \mathbf{w} \phi_{PL}(t_2) = x_0 w_1 \mathbf{w} \ldots \mathbf{w} x_0 w_k \mathbf{w} x_0 w'_1 \mathbf{w} \ldots \mathbf{w} x_0 w'_l \mathbf{w} x_0^{p+q-k-l-1}.
\]

So \( \phi_{PL} \) is a Com-Prelie algebra morphism. \( \square \)

**Remark.** By the proof of proposition 28, we have a commutative diagram of prelie algebra morphisms:

\[
\mathcal{G}_{\mathcal{T}(\mathbb{N}^*)} \xrightarrow{\psi} \mathcal{G}_{\mathcal{PT}(\{1,2\})} \xrightarrow{\phi_{CPL}} \mathbb{K}\langle x_0, x_1 \rangle
\]

\[
\mathcal{G}_{\mathcal{T}(\mathbb{N}^*)} \xrightarrow{\phi_{PL}}
\]

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Moreover, $\phi_{CPL}$ is a morphism of Com-Prelie algebra. With the commutative, associative product previously defined on $\mathcal{O}(\mathbb{N}^*)$, $\phi_{PL}$ is now a morphism of Com-Prelie algebra. However, $\psi$ is not compatible with $\mathcal{W}$. Indeed, $\psi(1^2_1) = \psi(\cdot) \bullet \psi(\cdot) = 1^2_1$, so:

$$\psi(1^2_1) \mathcal{W} \psi(1^2_1) = 1^2_1 \mathcal{W} 1^2_1 = 1^2_1 \mathcal{W}.$$ 

Moreover, $1^2_1 \mathcal{W} 1^2_1 = 1^4_1$, so:

$$\psi(1^2_1 \mathcal{W} 1^2_1) = \psi(\cdot) \bullet \psi(\cdot) = 1^4_1 \mathcal{W}.$$ 

5 Appendix

5.1 Enumeration of partitioned trees

Let $d \geq 1$. For all $n \geq 1$, let $f_n$ be the number of partitioned trees decorated by $\{1, \ldots, d\}$ with $n$ vertices and let $t_n$ be the number of partitioned trees decorated by $\{1, \ldots, d\}$ with $n$ vertices and one root. By convention, $f_0 = 1$. We put:

$$T = \sum_{n=1}^{\infty} t_n X^n, \quad F = \sum_{n=0}^{\infty} f_n X^n.$$ 

Let $V_T$ be the vector space generated by the set of partitioned trees decorated by $\{1, \ldots, d\}$ and $V_F$ be the vector space generated by the set of partitioned trees decorated by $\{1, \ldots, d\}$ with only one root. There is a bijection:

$$\begin{cases}
S(V_T) & \longrightarrow V_F \\
t_1 \ldots t_k & \longrightarrow t_1 \mathcal{W} \ldots \mathcal{W} t_k.
\end{cases}$$ 

Hence:

$$F = \prod_{i=1}^{\infty} \frac{1}{(1 - X^k)^k}. \quad (2)$$

There is a bijection:

$$\begin{cases}
\bigoplus_{i=1}^{d} S(V_F) & \longrightarrow V_T \\
(F_{i_1,1} \ldots, F_{i_1,k_1}, \ldots, F_{d,1} \ldots F_{d,k_d}) & \longrightarrow \prod_{i=1}^{d} \bullet (F_{i_1} \ldots F_{i_1}).
\end{cases}$$

This gives:

$$T = dX \prod_{i=1}^{\infty} \frac{1}{(1 - X^k)^{k-1}}. \quad (3)$$

Formulas (2) and (3) allow to compute inductively $f_k$ and $t_k$ for all $k \geq 1$. This gives:

$$\begin{align*}
f_1 &= d \\
f_2 &= \frac{d(3d + 1)}{2} \\
f_3 &= \frac{d(19d^2 + 9d + 2)}{6} \\
f_4 &= \frac{d(63d^2 + 34d^2 + 13d + 2)}{8} \\
f_5 &= \frac{d(644d^4 + 400d^3 + 175d^2 + 35d + 6)}{30}
\end{align*}$$
Here are examples of $f_n$ for $d = 1$ or 2:

<table>
<thead>
<tr>
<th>$n$</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
<th>5</th>
<th>6</th>
<th>7</th>
<th>8</th>
<th>9</th>
<th>10</th>
</tr>
</thead>
<tbody>
<tr>
<td>$d = 1$</td>
<td>1</td>
<td>2</td>
<td>5</td>
<td>14</td>
<td>42</td>
<td>134</td>
<td>444</td>
<td>1518</td>
<td>5318</td>
<td>18989</td>
</tr>
<tr>
<td>$d = 2$</td>
<td>2</td>
<td>7</td>
<td>32</td>
<td>167</td>
<td>952</td>
<td>5759</td>
<td>36340</td>
<td>236498</td>
<td>1576156</td>
<td>10702333</td>
</tr>
</tbody>
</table>

The row $d = 1$ is sequence A035052 of [16].

### 5.2 Study of the dendriform structure on admissible words

We here study the dendriform algebra $\mathbb{K}\langle \mathbb{N}^* \rangle$ of proposition 32. It is clearly commutative, via the bijection from $Sh_{\prec}(k,l)$ to $Sh_{\prec}(l,k)$ given by the composition (on the left) by the permutation $(l + 1 \ldots l + k 1 \ldots l)$: in other terms, it is a Zinbiel algebra [8].

Let $V$ be a vector space. The shuffle dendriform algebra $Sh(V)$ is $T_+(V)$, with the products given by:

$$
(a_1 \ldots a_k) \prec (b_1 \ldots b_l) = \sum_{\zeta \in Sh_{\prec}(k,l)} \zeta, a_1 \ldots a_k b_1 \ldots b_{l-1} b_k
$$

$$
(a_1 \ldots a_k) \succ (b_1 \ldots b_l) = \sum_{\zeta \in Sh_{\prec}(k,l)} \zeta, a_1 \ldots a_{k-1} a_k b_1 \ldots b_k.
$$

Moreover, this is the free commutative dendriform algebra generated by $V$, that is to say if $A$ is a commutative dendriform algebra and $f : V \rightarrow A$ is any linear map, there exists a morphism of dendriform algebras $\phi : Sh(V) \rightarrow A$ such that $\phi(V) = f$. As $a_1 \ldots a_k \succ b = a_1 \ldots a_k b$ in $Sh(V)$ for all $a_1, \ldots, a_k, b \in V$, this morphism $\phi$ is defined by:

$$
\phi(a_1 \ldots a_k) = (\ldots (a_1 \succ a_2) \succ a_3) \ldots) \succ a_k.
$$

**Proposition 38**  
1. Let $V$ be the space generated by the words $1^k i$, $k \in \mathbb{N}$, $i \geq 1$. Then $K\langle \mathbb{N}^* \rangle$ is isomorphic, as a dendriform algebra, to $Sh(V)$.

2. Let $A$ be the subspace of $K\langle \mathbb{N}^* \rangle$ generated by admissible words. Then it is a dendriform subalgebra of $K\langle \mathbb{N}^* \rangle$. Moreover, if $W$ is the space generated by the letters $i$, $i \geq 1$, then $A$ is isomorphic, as a dendriform algebra, to $Sh(W)$.

**Proof.** Let $w = a_1 \ldots a_k$ be a word with letters in $\mathbb{N}^*$. We denote by $o(w)$ the sequence of indices $j \in \{1, \ldots, k-1\}$ such that $a_j \neq 1$. These sequences are totally ordered in this way: $(j_1, \ldots, j_k) < (j'_1, \ldots, j'_l)$ if there exists a $p$ such that $j_k = j'_1$, $j_{k-1} = j'_1$, \ldots, $j_{k-p+1} = j'_{l-p+1}$, $j_{k-p} < j'_{l-p}$, with the convention $j_0 = j_{-1} = \ldots = j'_0 = j'_{-1} = \ldots = 0$.

Let $\phi : Sh(V) \rightarrow K\langle \mathbb{N}^* \rangle$ be the unique morphism of dendriform algebras which extends the identity of $V$. Then:

$$
\phi((1^{k_1-1} a_1) \ldots (1^{k_n-1} a_n)) = 1^{k_1-1}(a_1 + 1) \ldots 1^{k_{n-1}-1}(a_{n-1} + 1) 1^{k_{n-1}-1} a_n
$$

+ words $w'$ such that $o(w') > (k_1, \ldots, k_{n-1})$.

By triangularity, $\phi$ is an isomorphism. Moreover, for all $a_1, \ldots, a_n \geq 1$:

$$
\phi(a_1 \ldots a_n) = (a_1 + 1) \ldots (a_{n-1} + 1) a_n.
$$

Consequently, $\phi(Sh(W)) = A$, so $A$ is a dendriform subalgebra of $K\langle \mathbb{N}^* \rangle$ and is isomorphic to $Sh(W)$.

\[ \square \]
5.3 Freeness of the pre-Lie algebra \( g_{\mathcal{PT}(D)} \)

**Notations.** Let \( k \geq 1, d_1, \ldots, d_k \in D \) and let \( F_1, \ldots, F_k \) be decorated partitioned forests. We put:

\[
B_{d_1, \ldots, d_k}(F_1, \ldots, F_k) = (\cdot_{d_1} \cdot F_1) \mathcal{W} \cdots \mathcal{W}(\cdot_{d_k} \cdot F_k).
\]

Note that any partitioned tree can be written under the form \( B_{d_1, \ldots, d_k}(F_1, \ldots, F_k) \). This writing is unique up to a common permutation of the \( d_i \)'s and the \( F_i \)'s.

**Proposition 39** We define a coproduct \( \delta \) on \( g_{\mathcal{PT}(D)} \) in the following way: for any decorated partitioned tree \( t = B_{d_1, \ldots, d_k}(t_1, \ldots, t_{n_1}, \ldots, t_{k,1} \cdots t_{k,n_k}) \),

\[
\delta(t) = \frac{1}{k} \sum_{i=1}^{k} \sum_{j=1}^{n_i} B_{d_1, \ldots, d_k}(t_1, \ldots, t_{n_1}, \ldots, t_{i,1} \cdots t_{i,j-1} t_{i,j+1} \cdots t_{i,n_i}, \ldots, t_{k,1} \cdots t_{k,n_k}) \otimes t_{i,j}.
\]

1. For all \( x \in g_{\mathcal{PT}(D)} \), \( (\delta \otimes \text{Id}) \circ \delta(x) = (23)(\delta \otimes \text{Id}) \circ \delta(x) \).
2. For all \( x, y \in g_{\mathcal{PT}(D)} \), \( \delta(x \bullet y) = x \otimes y + \delta(x) \bullet y \).

**Proof.** 1. Let \( t = B_{d_1, \ldots, d_k}(t_1, \ldots, t_{n_1}, \ldots, t_{k,1} \cdots t_{k,n_k}) \). For all \( i, j \), we put:

\[
t_{i,j} = B_{d_1, \ldots, d_k}(t_1, \ldots, t_{n_1}, \ldots, t_{i,1} \cdots t_{i,j-1} t_{i,j+1} \cdots t_{i,n_i}, \ldots, t_{k,1} \cdots t_{k,n_k}).
\]

Then:

\[
\delta(t) = \frac{1}{k} \sum_{i,j} t_{i,j} \otimes t_{i,j}.
\]

Hence:

\[
(\delta \otimes \text{Id}) \circ \delta(t) = \sum_{(i,j) \neq (i',j')} (t_{i,j})/t_{i',j'} \otimes t_{i,j}
\]

As \( (t_{i,j})/t_{i',j'} \) and \( (t_{i',j'})/t_{i,j} \) are both the partitioned tree obtained by cutting \( t_{i,j} \) and \( t_{i',j'} \) in \( t \), they are equal, so \( (\delta \otimes \text{Id}) \circ \delta(t) \) is invariant under the action of (23).

2. Let \( t' \) be a decorated partitioned tree.

\[
\delta(t \bullet t') = \sum_{i=1}^{k} \delta(B_{d_1, \ldots, d_k}(t_1, \ldots, t_{n_1}, \ldots, t_{i,1} \cdots t_{i,n_i} t', \ldots, t_{k,1} \cdots t_{k,n_k}))
+ \sum_{i,j} \delta(B_{d_1, \ldots, d_k}(t_1, \ldots, t_{n_1}, \ldots, t_{i,1} \cdots t_{i,j} \bullet t' \cdots t_{i,n_i}, \ldots, t_{k,1} \cdots t_{k,n_k}))
\]

\[
= \frac{1}{k} t \otimes t' + \frac{1}{k} \sum_{i} B_{d_1, \ldots, d_k}(t_1, \ldots, t_{n_1}, \ldots, t_{i,1} \cdots t_{i,n_i} t', \ldots, t_{k,1} \cdots t_{k,n_k})/t_{i',j'} \otimes t_{i',j'}
+ \frac{1}{k} \sum_{i,j} B_{d_1, \ldots, d_k}(t_1, \ldots, t_{n_1}, \ldots, t_{i,1} \cdots t_{i,j} \bullet t' \cdots t_{i,n_i}, \ldots, t_{k,1} \cdots t_{k,n_k})/t_{i',j'} \otimes t_{i',j'}
+ \frac{1}{k} \sum_{i,j} t_{i,j} \otimes t'
\]

\[
= t \otimes t' + \sum t^{(1)} \otimes t^{(2)} \bullet t' + \sum t^{(1)} \otimes t^{(2)} \bullet t'.
\]

So \( \delta(t \bullet t') = t \otimes t' + \delta(t) \bullet t' \) by Livernet's pre-Lie rigidity theorem [7].

**Corollary 40** The pre-Lie algebra \( g_{\mathcal{PT}(D)} \) is freely generated by \( \text{Ker}(\delta) \).
Remarks.

1. It is not difficult to prove that for any \( x, y \in \mathcal{P}(\mathcal{T}) \):
\[
\delta(x \mathcal{U} y) = \sum x^{(1)} \otimes x^{(2)} \mathcal{U} y + \sum y^{(1)} \otimes x \mathcal{U} y^{(2)}.\]
Hence, \( \text{Ker}(\delta) \) is an algebra for the product \( \mathcal{U} \).

2. Here are elements of \( \text{Ker}(\delta) \) in the non decorated case. Let \( t_1, t_2, t_3, t_4 \) be partitioned trees.
\[
X = B(t_1t_2, 1) - B(t_1, t_2),
Y = B(t_1t_2t_3, 1, 1) - B(t_1t_2, t_3, 1) - B(t_1t_3, t_1, 1) - B(t_2t_3, t_1, 1) + 2B(t_1, t_2, t_3),
Z = B(t_1t_2t_3t_4, 1) - B(t_1t_2t_3, t_4) - B(t_1t_2t_4, t_3) - B(t_1t_3t_4, t_2) - B(t_2t_3t_4, t_1)
\quad + B(t_1t_2, t_3t_4) + B(t_1t_3, t_2t_4) + b(t_1t_4, t_2t_3),
T = B(t_1t_2, t_3t_4, 1, 1) + B(t_1t_3, t_2t_4, 1, 1) + B(t_1t_4, t_2t_3, 1, 1) - B(t_1t_2, t_3, 4, 1)
\quad - B(t_1t_3, t_2, t_4, 1) - B(t_1t_4, t_2, t_3, 1) - B(t_2t_3, t_1, t_4, 1) - B(t_2t_4, t_1, t_3, 1)
\quad - B(t_3t_4, t_1, t_2, 1) + 3B(t_1, t_2, t_3, t_4).
\]

References


