Parallel Programming with SystemC for Loosely Timed Models: A Non-Intrusive Approach

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Abstract—The SystemC/TLM technologies are widely accepted in the industry for fast system-level simulation. An important limitation of SystemC regarding performance is that the reference implementation is sequential, and the official semantics makes parallel executions difficult. As the number of cores in computers increase quickly, the ability to take advantage of the host parallelism during a simulation is becoming a major concern. Most existing work on parallelization of SystemC targets cycle-accurate simulation, and would be inefficient on loosely timed systems since they cannot run in parallel processes that do not execute simultaneously.

We propose an approach that explicitly targets loosely timed systems, and offers the user a set of primitives to express tasks with duration, as opposed to the notion of time in SystemC which allows only instantaneous computations and time elapses without computation. Our tool exploits this notion of duration to run the simulation in parallel. It runs on top of any (unmodified) SystemC implementation, which lets legacy SystemC code continue running as-it-is. This allows the user to focus on the performance-critical parts of the program that need to be parallelized.

I. INTRODUCTION

Systems-on-a-Chip are complex systems involving dedicated hardware components and one or several processors. Because of time to market pressure, it is not possible to wait for the physical chip to develop the software, hence the developers need simulators, or virtual prototypes, earlier in the design flow. These prototypes need to be fast enough to make it possible to run complex applications, and must expose enough details of the hardware to allow non-portable software to run. The RTL level of abstraction is detailed enough, but far too slow, and available too late in the design cycle, hence a new level of abstraction called Transaction Level Modeling (TLM) has been introduced. We are interested in the SystemC [1] implementation of TLM.

There are several sub-levels of abstraction within TLM depending on the target usage: if the virtual prototype is meant for performance evaluation, it needs to be precise with respect to timing (either cycle-accurate or approximately timed). On the other hand, if only the functional correctness is important, timing can be relaxed considerably. TLM allows a purely untimed coding style, as well as a loosely timed (LT) coding style where the behavior is timed only if the timing is needed for the functionality. For example, a loosely timed model of a timer will usually be timed properly, but most computations can be considered instantaneous.

Although SystemC represents concurrent systems, it should be noted that timing and concurrency in SystemC are simulated, and the actual execution is usually sequential (The SystemC standard imposes co-routine semantics). Running the simulation in parallel is needed to take advantage of today’s machines, and SystemC’s sequentiality could become a major performance bottleneck in the future. Many approaches have been proposed for SystemC parallelization, but most of them are only able to run processes that run at the same simulation cycle, and hence apply only on cycle-accurate systems.

As opposed to this, the parallelization approach we propose targets loosely timed systems. We present a programming model in which processes can not only execute instantaneous tasks or wait statements, but also express the notion of duration, i.e. specify the start and end date of the computation and let the scheduler freely spread it across time. These tasks with duration are executed in parallel with the SystemC scheduler.

This approach is implemented in a library, that runs on top of an unmodified SystemC kernel (including proprietary ones that cannot be modified). Also, this means that legacy code will continue running with our approach, although they will not automatically benefit from parallelism. The user can identify performance-critical parts of the system, and focus on them to parallelize the program (adding synchronization as required).

The paper is structured as follows: Section [II] discusses related works. Section [III] presents the main contribution: the during tasks principle and implementation, for which Section [IV] give experimental results. Section [V] is the conclusion.

II. RELATED WORK

As the official semantics for SystemC is to use co-routine semantics, the natural way to execute a SystemC program is to execute it sequentially. To take advantage of the parallelism of the host machine, it is tempting to run SystemC threads in parallel. The SystemC standard allows this, “provided that the behavior appears identical to the co-routine semantics” [1]. This implies two constraints on a parallel implementation:

1) It should not change the order in which processes are allowed to be executed. In particular, the simulated time imposes an order on the execution of processes.

2) It should not introduce new race conditions. For example, two SystemC processes may safely execute $x++$ on a shared variable, but running two such processes in parallel cannot be allowed.

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Approaches for parallel execution of SystemC programs can be classified depending on the way they deal with these constraints: conservative ensure them by construction, while optimistic approaches relax them, ideally with a verification and rollback mechanism to correct them after the fact.

\subsection{A. Semantics-Preserving Parallelization}

An ideal conservative approach would preserve the official semantics completely. To solve problem 1 above, one can introduce parallelism within a \( \delta \)-cycle (i.e. within one iteration inside a simulation instant), but keep a global synchronization barrier between each \( \delta \)-cycle. Since the order of execution within an evaluation phase is unspecified in SystemC, it is correct to run processes in any order.

To solve problem 2 one needs to ensure that two processes accessing the same variable will never be run in parallel. This can be done structurally, under the assumption that a process will only access variables declared within its own module. It is then sufficient to ensure that at most one process is running within each module. SystemCASS \[2\] is an example tool using this assumption to perform a dependency analysis (for static scheduling). Unfortunately, communication using TLM breaks this assumption: TLM communication channels transport transaction through function calls, hence the code processing a transaction in a target module is executed in the context of the initiator module. In particular, the natural implementation of a RAM makes it a shared array between all components accessing it.

Another option is to statically analyze the code of the processes to detect accesses to shared variables. \[3\] proposes an approach to parallelization of SystemC based on a dependency analysis performed on a SystemC-transition (i.e. piece of code executed between two \textit{wait} statements) basis. \[4\] applies similar static analysis for SpecC. Scoot \[5\] proposes a finer analysis using model-checking (applied to scheduling and model-checking, not to parallelization). These approaches are heavyweight since they require a static analysis of arbitrary C++ code, and require the complete source code of the model to be available. The efficiency of these approaches is disappointing on TLM models for which a few target components create a synchronization bottleneck (back to our RAM example, two processes that may access the RAM never be executed in parallel).

\subsection{B. Parallelism within \( \delta \)-cycles}

A second category of approaches is to be conservative with respect to problem 1 but to leave it up to the user to deal with problem 2 by avoiding shared variables, and using proper synchronization and locking when needed. \[6\] actually distributes the simulation over several machines, running one SystemC scheduler on each node. The first version enforced a global synchronization at each \( \delta \)-cycle, and further optimizations \[7\] allowed relaxing this synchronization by computing the next simulated time locally on each node.

Another approach is to modify the SystemC kernel to allow a parallel simulation \[8\]. The scheduler maintains a set of runnable processes, and several threads run an evaluation phase that pick processes from this set and run them, finishing the evaluation phase with a global synchronization barrier. This work is currently focused on performance, but leaves the responsibility of protecting the atomicity of accesses to shared data-structure to the SystemC processes. Similar work propose some performance optimizations using various load-balancing techniques like work-stealing, work-sharing or manual grouping \[9, 10\] to improve the cache usage.

\subsection{C. Parallelism Across Simulated Instants}

All the works presented above have in common the fact that they parallelize the content of \( \delta \)-cycle, but cannot run in parallel processes that should be executed at different simulated times. In other words, they require a simultaneity between processes to take advantage of parallelism. This is appropriate for cycle-accurate simulation, where clock ticks usually wakes up many processes at the same time. But higher levels of abstraction tend to eliminate clock, and deal with quantitative time instead. Figure 1 shows an example modeling a system performing two computations in parallel. Computations that take some time are usually modeled by instantaneous computations followed by a SystemC wait. For example, TLM-2’s temporal decoupling \[11\] produces this effect (the call to \textit{wait} is done when synchronizing a process’s local clock with the simulation time). In simulation, this program will run without ever having two processes runnable at the same time. In the presence of loose timing annotations (e.g. \textit{wait}(X\pm 50\%)), the situation is even worse since the argument to \textit{wait} becomes a random value.

\begin{verbatim}
void P::compute() {
  P1(); wait(25, SC_MS);
  P2(); wait(12, SC_MS);
}

void Q::compute() {
  Q1(); wait(10, SC_MS);
  Q2(); wait(11, SC_MS);
}
\end{verbatim}

\textbf{fig. 1. Parallelism Without Simultaneity}

For such systems, semantics-preserving approaches are even harder. The scheme proposed in \[4\] performs a static analysis for processes that may be executed at different instants, but in SystemC/TLM programs with a shared bus, target components will still create conflicts in many cases, and this approach requires the availability of the whole source code. The alternative is to use optimistic approaches, i.e. one processes should be allowed to advance its local time without waiting for other computations to terminate. Maintaining the same execution order as the reference semantics becomes intractable, hence optimistic simulation either (a) requires a violation-detection together with a rollback mechanism or (b) deliberately accepts some order violations, even though it can control them with some logical clock mechanism in addition to the quantitative simulated time. Approaches in category...
(a) have been successfully applied to VHDL simulations, but
the case of SystemC is much harder to deal with: since the
user can call arbitrary C++ code and system calls, a rollback
mechanism is not implementable in general.

One solution in category (b) is proposed in [12], introducing
the TLM with Distributed Time (TLM-DT) abstraction level
and the associated parallel simulator. In TLM-DT, processes
do not rely on the SystemC simulated time, but manage a local
clock, that is synchronized with other components regularly
(on every communication, and using null messages when a pre-
defined time quantum elapsed without actual communication).
This way, ordering violations can happen only with processes
running with a difference of at most the time quantum. This
approach does not need simultaneity for parallelization, and
relaxes the synchronization barrier, but still enforces it every
time quantum (through null messages). It requires a particular
modeling style and forbids some SystemC primitives, hence
is not applicable on legacy code. It is designed to be efficient
on approximately timed model, but not on loosely timed ones.

Another approach relying on the notion of time quantum is
presented in [13]. It performs a distributed simulation similarly
to [6], but instead of performing a synchronization barrier
for every δ-cycle, the authors propose to force processes to
execute only on a multiple of the quantum.

Tasks and Duration: An experiment of transaction-level
modeling outside the SystemC world is jTLM [14]. The
management of simulated time is different in that jTLM
provides a notion of “task with duration” [15]. Instead of
modeling such task as an instantaneous computation followed
by a wait, the user can specify the duration of the task, and let
the scheduler spread the computation over the specified period
of time. The example of Figure 1 becomes the one in Figure 2.

```
void P_compute() {
  consumesTime(25) {P1();
  consumesTime(12) {P2();
}
```

Fig. 2. Parallelism With Duration in jTLM

This way, tasks P1() and Q1(), although not simultaneous,
have an overlap, hence can run in parallel.

This notion of duration and overlap is interesting for parallel
execution of models with a coarse timing granularity. The
coarsest the task are, the less likely it is for their start and end
points to be simultaneous, but the most likely it is for them
to overlap. Unlike cycle-based approaches, the efficiency will
therefore increase with the size of tasks. An obvious limitation
of this approach is that jTLM is not compatible with SystemC,
therefore unable to use legacy SystemC code.

III. TASKS WITH DURATION IN SYSTEMC

None of the approaches mentioned above can parallelize
the execution of SystemC programs across instants safely
and without changing the code. It does not seem realistic
to target fully automatic parallelization for real-life TLM
programs. Instead, our contribution is a new programming
model that complements the one of SystemC and allows the
programmer to solve the parallelization problem by using the
new primitives on performance-critical parts of the program.
The semantics is inspired from the one of jTLM, but the
implementation is made as a library called sc-during
(while jTLM implements tasks with a dedicated scheduler).
It runs on top of SystemC, without modifying it. The library
can be freely downloaded from http://sc-during.forge.imag.fr/

A. Principle of Tasks with Duration

An example code using sc-during is shown in Figure 3.

```
1 extern void P1(), P2(), Q1(), Q2();
2 struct P : sc_module, sc_during {
3   void compute() {
4     during(sc_time(25, SC_MS), P1);
5     during(sc_time(12, SC_MS), P2);
6   };
7   SCCTOR(P) (SC_THREAD(compute));
8 };
9 struct Q : sc_module, sc_during {
10   void compute() {
11     wait(10, SC_MS);
12     during(sc_time(13, SC_MS), Q1);
13     during(sc_time(11, SC_MS), Q2);
14   };
15   SCCTOR(Q) (SC_THREAD(compute));
16 };
17 int sc_main(int, char **) {
18   P p("p"); Q q("q"); sc_start(); return 1; }
```

Fig. 3. Parallelism With Duration in SystemC

The execution of this program is illustrated in Figure 4.
It starts the normal way, using the SystemC scheduler and
global simulated time. Once a call to during is encountered,
the execution of the function passed as argument (e.g. P1, at
t = 0) can be delegated to another operating system thread.
The simulation will continue with other processes on the SystemC side, in parallel with the execution of P1. SystemC and P1 will join at t = 25 ms, i.e. either the SystemC simulation reaches 25 and waits for P1 to complete, or P1 terminates at t < 25.
While executing $P_1$ as a separate task, the SystemC scheduler continues its execution, executes $\text{wait}(10, \text{SC_MS})$; and reaches $t = 10$. Then, it encounters another call to during and launches $Q_1$ in a separate task, just like it did for $P_1$. At this point, the simulation can contain 3 OS threads: one for each task, and one running the SystemC kernel. As $P_1$ and $Q_1$ have an overlap in time, they can be executed in any order, or with any interleaving. On the other hand, the end of task $Q_1$ predates the start of $P_2$, so the scheduler ensures that $Q_1$ is completed before starting $Q_2$.

It is of course possible to have instantaneous computations too. Any code within a SystemC process and outside a call to during will run instantaneously, following the SystemC semantics and independently from the sc-during library (another way to create instantaneous tasks is during(SC_ZERO_TIME, f) which creates a separate task that must complete before the end of the next $\delta$-cycle).

The sc-during library contains 4 different implementations of the notion of task with duration. We first illustrate the idea with the one called THREAD, which is straightforward and unoptimized. Figure 5 shows a simplified version of this implementation (more implementation details are needed to introduce other primitives in section III-B).

```
1 void during(sc_core::sc_time duration,
2     boost::function<void()> routine) {
3     boost::thread t(routine); // create thread
4     sc_core::wait(t); // let SystemC execute
5     t.join(); // wait for thread completion
6 }
```

Fig. 5. Simplified Implementation of during()

B. Synchronization and Time Modeling

The ideas presented up to now allow running a piece of code in parallel with a SystemC simulation, but up to now, the task can hardly communicate with the SystemC thread. In particular, it cannot call SystemC primitives like wait() and notify(), and shared variables between the SystemC thread would create race-conditions. In order to provide a complete programming model to the user, we need to introduce additional synchronization primitives.

1) Full Synchronization with SystemC: sc_call(): A during task may want to execute a piece of code within the context of the SystemC thread. This way, this piece of code will be allowed to execute SystemC primitives (e.g. wait, notify ...) the normal way, and to access variables shared with SystemC without specific locking or race condition (like two SystemC processes access shared variables).

The sc_call() API function does this: it is callable from a during task and takes a function as argument. The current task is blocked, and the function is scheduled to be executed next time SystemC allows it. The during task is unlocked once the function has completed its execution.

We provide two implementations for the scheduling of the function to be executed: In SystemC 2.2, we simply record the function to be called, and call it when the call to sc_core::wait associated to the task terminates. In SystemC 2.3, we use async_request_update() to schedule the execution during the next $\delta$-cycle.

2) Controlling Time Elapse and Duration: extra_time() and catch_up(): Up to now, the only way to specify the duration of a task is to provide it as argument to the during() method. This is not always possible, since it would require knowing how long the computation done within the task would take, before executing it.

To solve this issue, the sc-during library provides the function extra_time(sc_time t), that can be called from a during task. This function increases the duration of the current task by t. In the implementation of the library, this means the wait statement line 4 in Fig 5 is actually inside a loop along the lines of while (remaining_time) wait(remaining_time), and extra_time increments remaining_time with proper synchronization. An example is shown in Figure 6.

```
1 void trace(string s) {
2     cout << sc_time_stamp() << ": " << s << endl; }
3 SC_MODULE(A), sc_core {
4     void P() {
5         wait(5, SC_MS);
6         during(5, SC_MS, boost::bind(&A::f, this));
7         trace("done");
8     }
9     void f(void) {
10        trace("before extra_time");
11        extra_time(7, SC_MS);
12        trace("after extra_time");
13    };
14    SCCTOR(A) { SC_THREAD(P); };
15 }
```

Fig. 6. Example Code Using extra_time(), and Execution Trace

This program will produce an output of the form (output is generated by the trace function defined line 1 which displays the current simulated time and a string):

```
X ms: before extra_time
Y ms: after extra_time
17 ms: done
```

The semantics ensures that the duration of the task is the initial duration plus the sum of times accumulated with extra_time(), hence the last line 17 ms: done. Before the call to extra_time(), the during task spans only over $[5,10]$, hence $X \in [5,10]$ is ensured. After calling extra_time(), the task spans over $[5,17]$, hence $Y \in [X,17]$. In general, the semantics of during tasks with extra_time is that a piece of code within the task can be executed at a simulated time $t \in [s,s+d+e]$ where $s$ is the starting time of the task, $d$ is the initial duration, and $e$ is the sum of the extra-time accumulated up to this point.

By default, time may flow freely during the execution of
a during task. This may be problematic if the program relies on some fairness between a during task and the SystemC simulation. Take the example of a polling loop like `while (t) {}`. In SystemC, such a loop which doesn’t let time elapse will create an infinite loop, and will have to be rewritten `while ((t) {wait(t);})` (for some time constant t). Within a during task, it is possible to achieve a similar effect with `while ((t) {extra_time(t);})`. Unfortunately, while this piece of code creates a during task with an unbounded duration, it does not ensure that the SystemC simulated time will actually increase. The `catch_up` function may be used as follows to ensure fairness:

```c
void catch_up();
void wait();
```

If this code is executed starting at t = 0, in a during task whose duration is initially 0, it will produce an output of the form:

```
10 ms: after
20 ms: after
```

with X1 ∈ [0, 10] and X2 ∈ [10, 20]. The simulated time after returning from `catch_up` is fixed to the duration of the task at this point.

Alternatively, the user can use the relaxed form `catch_up(t)`, providing as argument a time t and that only ensures that the simulated time will advance until the difference between the current time and the end of the task is at most t.

### C. Temporal Decoupling and Duration

**Temporal decoupling** is a technique standardized by TLM-2 to reduce the number of `wait` statements executed. The idea is to allow processes to maintain a local lock t_local to model the fact that the process is “in advance” with respect to the SystemC time. Instead of using `wait(d)`, processes increase t_local by d (function `annotate` in Figure 7), and periodically call a function to actually perform the `wait()` statement (function `synchronize` in Figure 8).

```c
void annotate(sc_time d) {
  t_local = t_local + d;
  t_local = 0;
  t_local = t_local + d;
}

void synchronize() {
  t_local = t_local + d;
  t_local = 0;
  t_local = t_local + d;
}
```

*Fig. 7. Simple Temporal Decoupling API*

```c
void annotate(sc_time d) {
  extra_time(d);
  catch_up();
}

void synchronize() {
  catch_up();
}
```

*Fig. 8. Temporal Decoupling API within during Tasks*

This `synchronize` function can be called whenever t_local becomes greater than a predefined time quantum (a call can be inserted in `annotate` for instance), or before synchronization points [10]. The SC-during library can benefit from temporal decoupling. The primitives of Figure 7 can be re-written as Figure 8 when used inside (possibly unbounded) during tasks. This preserves the property that important actions (i.e. the ones after calling `synchronize`) are executed at the right SystemC time, and the computations calling `annotate` are modeled with the right duration.

### D. Implementation

We gave in section III-C a sketch of the THREAD implementation for during tasks. The actual implementation is more complex since it needs to manage `sc_call`, `extra_time` and `catch_up`. We also provide several implementation of parallelism: POOL pre-allocates a set of threads and works in a producer-consumer scheme, ONDEMAND creates new threads as needed and re-uses them instead of deleting and re-creating them, and SEQ is a reference sequential implementation.

### IV. EXPERIMENTAL RESULTS

#### A. Experimental Setup and Platform Synchronization

We experimented our approach on a SystemC model of an existing platform containing a CPU, a timer, a GPIO, a RAM and an interrupt controller (ITC), running on an FPGA, and executing an embedded software running Conway’s Game of Life (the same C code is running on the FPGA and on the TLM platform). To illustrate parallelization of simulation, we added multi-processing capabilities to the system, adding N additional CPUs and ITCs. CPU0 delegates parts of the computation (slices of the image to compute) to other CPUs. It does so by setting a flag in RAM, and triggering an interrupt using the ITC associated to the CPU. Bus-communication is abstracted with a TLM-2 channel. Transactions (either read or write) are initiated by a component, routed through a Bus module towards a target module.

The most performance-consuming components are the CPUs, on which we are now going to focus. We use a simple ISS, but our approach would work equally well with any other SystemC module provided the approach has some notion of timing (possibly approximate). One ISS step takes a constant amount of time (PERIOD), hence we can basically count the number N of steps, and wait for N*PERIOD when synchronizing with SystemC time. We tried several strategies to choose when and how to synchronize:

The first one is called during_quant and is directly inspired from the notion of quantum. “Packs” of 5000 instructions are executed in a during task of duration d = 5000*PERIOD. The second one, called during_sync, requires the user to explicitly set the synchronization points of the model (writing to some magic values in the addressmap), and is an application of Section III-C: we start a during task initially empty, call `extra_time(5000*PERIOD)` every 5000 ISS steps, and call `extra_time(N*PERIOD)` N = 0; `catch_up()` at every synchronization point (i.e. before sending or waiting for an interrupt, within polling loops, ...).

The ISS triggers reads and writes on the bus, which triggers execution of handlers in the target modules. When these
handlers do not call SystemC primitives, and if the data-access they perform are safe, we can let them run in the context of the during task. But for example, writing on ITC component may trigger an e.notify(), which cannot be run from a during task. When triggering a write transaction to such modules, the ISS wraps the bus access in an sc_call, so that the transaction be executed in the context of the SC_THREAD, like in the sequential version. We also wrap wait for interrupts in sc_calls. Most accesses to the bus are performed without synchronization; the bus component itself is thread-safe (no state variable except the addressmap, which is read-only during simulation), and memory access that do not require particular synchronization in the actual system are performed without synchronization in simulation.

B. Speedup and Scalability

We tested the scalability of the approach on two machines. The first one (Figure 9(a)) has 4 Intel Xeon cores. The second (Figure 9(b)) has 4 AMD Opteron-6176 processors, each containing 12 cores (total of 48 cores).

![Speedup](image)

Fig. 9. Speedup Compared to SEQ for Different Implementation (average, min and max over 10 runs)

Both graphs show the benefit of defining explicit synchronization points, as the strategy during_sync always gives a better speedup than the quantum-based one. This also illustrate the fact that sc-during targets loosely timed systems, since coarser granularities models give better performance (unlike most alternative approaches). For low numbers of CPU in the model, the speedup is almost optimal, it does not decrease significantly when parallelizing the model too much (i.e. instantiation more CPU in the model than physical CPUs). With a large number of processors, even on the 48 cores machines, we can notice that the speedup is smaller than the number of processors. While disappointing, this result is not so surprising since the during task need some synchronization with SystemC (some bus accesses being protected with sc_call). We can further optimize the platform by making the ITC component itself thread-safe (by using async_request_update()), and get 13% additional speedup in ONDEMAND implementation and 23% in THREAD implementation, on machine (b) with 60 threads.

It should be noted that the suboptimal speedups are not due to extra computation, but essentially to latency of synchronization. Our experiments show that the CPU time spent in the computation increase by less than 35% even for very loaded simulation (60 CPUs). Unlike approaches using busy-waiting, this means that our parallel simulation will remain efficient on multi-users machines.

The during_quant strategy has a slightly lower speedup when the parallelism is low, but stagnates (and even decreases for the ONDEMAND implementation) with a large number of threads, because OS threads need frequent synchronization with SystemC, which becomes the synchronization bottleneck. This problem can be solved by increasing the quantum. For example, by setting the quantum to 50,000 instructions (i.e. 1 ms) instead of 5000, we speed up the simulation in THREAD implementation by a factor of 3.7. With a quantum of 10 ms, we get back to the same performance as the during_sync version. This illustrates the potential difficulty of choosing the right quantum, and the advantage of explicit synchronizations.

V. Conclusion

We presented an approach for parallel execution of SystemC models. A deliberate choice of the approach is that we require the user to use new primitives, as the alternatives would either be inefficient or would break the SystemC semantics for existing code. We showed that our approach allows a significant speedup with a reasonable modeling effort. By providing primitives to express the notion of duration, we allow the user to create more parallelism in the model than cycle-based approach, hence can efficiently parallelize loosely timed systems.

Our approach works with any SystemC scheduler, it is a simple library whose primitives are called directly from the SC_THREAD, hence intrusion in the user code is minimal. We could in theory run our library on top of a parallel implementation of SystemC, although such setup has not been experimented. Running code in a separate OS thread limits the interactions with SystemC, hence the approach applies best to parallelize computation that do not interact with the rest of the simulation, but we also provide primitives to allow synchronizing with SystemC, hence communication with other processes (sc_call, extra_time and catch_up).

Further work include providing more tools for synchronization to the user. For example, we manually wrapped in sc_call some transactions, but the Bus component could do it automatically for user-defined address ranges. This approach can be compared with the SpecC approach, where the language
reference allows parallelism within a component, but the communication channels act as monitors [17].

REFERENCES


