

On some sets of dictionaries whose ω -powers have a given complexity

Olivier Finkel *

Equipe de Logique Mathématique
 CNRS et Université Paris Diderot Paris 7
 UFR de Mathématiques case 7012, site Chevaleret,
 75205 Paris Cedex 13, France

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A dictionary is a set of finite words over some finite alphabet X . The ω -power of a dictionary V is the set of infinite words obtained by infinite concatenation of words in V . Lecomte studied in [1] the complexity of the set of dictionaries whose associated ω -powers have a given complexity. In particular, he considered the sets $\mathcal{W}(\Sigma_k^0)$ (respectively, $\mathcal{W}(\Pi_k^0)$, $\mathcal{W}(\Delta_1^1)$) of dictionaries $V \subseteq 2^*$ whose ω -powers are Σ_k^0 -sets (respectively, Π_k^0 -sets, Borel sets). In this paper we first establish a new relation between the sets $\mathcal{W}(\Sigma_2^0)$ and $\mathcal{W}(\Delta_1^1)$, showing that the set $\mathcal{W}(\Delta_1^1)$ is “more complex” than the set $\mathcal{W}(\Sigma_2^0)$. As an application we improve the lower bound on the complexity of $\mathcal{W}(\Delta_1^1)$ given by Lecomte, showing that $\mathcal{W}(\Delta_1^1)$ is in $\Sigma_2^1(2^{2^*}) \setminus \Pi_2^0$. Then we prove that, for every integer $k \geq 2$, (respectively, $k \geq 3$) the set of dictionaries $\mathcal{W}(\Pi_{k+1}^0)$ (respectively, $\mathcal{W}(\Sigma_{k+1}^0)$) is “more complex” than the set of dictionaries $\mathcal{W}(\Pi_k^0)$ (respectively, $\mathcal{W}(\Sigma_k^0)$).

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1 Introduction

A finitary language, called here also a dictionary as in [1], is a set of finite words over some finite alphabet X . The ω -power of a dictionary V is the set of infinite words obtained by infinite concatenation of words in V . The ω -powers appear very naturally in Theoretical Computer Science and in Formal Language Theory, in the characterization of the classes of languages of infinite words accepted by finite automata or by pushdown automata, [2].

Since the set of infinite words over a finite alphabet X is usually equipped with the Cantor topology, the question of the topological complexity of the ω -powers of finitary languages naturally arises. It has been posed by Niwinski [3], Simonnet [4] and Staiger [5].

Firstly it is easy to see that the ω -power of a finitary language V is always an analytic set because it is the continuous image of either a compact set $\{1, \dots, n\}^\omega$ for $n \geq 0$, or the Baire space ω^ω .

The first example of a finitary language L such that the ω -power L^ω is analytic but not Borel, and even Σ_1^1 -complete, was obtained in [6]. Amazingly the language L has a very simple description and was obtained via a coding of the infinite labelled binary trees. The construction will be recalled below. For the Borel ω -powers, after some partial results obtained in [7–9], the question of the Borel hierarchy of ω -powers of finitary languages has been solved recently by Finkel and Lecomte in [10], where a very surprising result is proved, showing that actually ω -powers exhibit a great topological complexity. For every non-null countable ordinal α there exist some Σ_α^0 -complete ω -powers and also some Π_α^0 -complete ω -powers.

* Corresponding author E-mail: finkel@logique.jussieu.fr

Another question naturally arises about ω -powers and descriptive set theory. It has been firstly studied by Lecomte in [1]. He asked about the complexity of the set of dictionaries whose associated ω -powers have a given complexity. The set $\mathcal{W}(\Sigma_\xi^0)$ (respectively, $\mathcal{W}(\Pi_\xi^0)$, $\mathcal{W}(\Delta_1^1)$) is the set of dictionaries over the alphabet $2 = \{0, 1\}$ whose ω -powers are Σ_ξ^0 -sets (respectively, Π_ξ^0 -sets, Borel sets). The set of dictionaries over the alphabet $2 = \{0, 1\}$ can be naturally equipped with the Cantor topology. Then Lecomte proved that $\mathcal{W}(\Sigma_2^0)$ is in $\Sigma_2^1(2^{2^*}) \setminus \Pi_2^0$ and that all the other sets $\mathcal{W}(\Sigma_\xi^0)$, $\mathcal{W}(\Pi_\xi^0)$, and $\mathcal{W}(\Delta_1^1)$ are in $\Sigma_2^1(2^{2^*}) \setminus D_2(\Sigma_1^0)$, where $D_2(\Sigma_1^0)$ is the class of 2-differences of open sets, that is, the class of sets which are intersections of an open set and of a closed set. It is proved in [11] that for each countable ordinal $\xi \geq 3$ the sets $\mathcal{W}(\Sigma_\xi^0)$ and $\mathcal{W}(\Pi_\xi^0)$ are actually Π_1^1 -hard. In this paper we obtain first a new relation between the sets $\mathcal{W}(\Sigma_2^0)$ and $\mathcal{W}(\Delta_1^1)$, showing that $\mathcal{W}(\Sigma_2^0)$ is continuously reducible to $\mathcal{W}(\Delta_1^1)$, which means that the set $\mathcal{W}(\Delta_1^1)$ is “more complex” than the set $\mathcal{W}(\Sigma_2^0)$. As an application we improve the lower bound on the complexity of $\mathcal{W}(\Delta_1^1)$ given by Lecomte, showing that $\mathcal{W}(\Delta_1^1)$ is in $\Sigma_2^1(2^{2^*}) \setminus \Pi_2^0$. Then we prove that, for every integer $k \geq 2$, (respectively, $k \geq 3$) the set of dictionaries $\mathcal{W}(\Pi_{k+1}^0)$ (respectively, $\mathcal{W}(\Sigma_{k+1}^0)$) is “more complex” than the set of dictionaries $\mathcal{W}(\Pi_k^0)$ (respectively, $\mathcal{W}(\Sigma_k^0)$).

The paper is organized as follows. In Section 2 we recall some notations of formal language theory and some notions of topology. We prove our results in Section 3. Some concluding remarks are given in Section 4.

2 Borel and projective hierarchies

We use usual notations of formal language theory which may be found for instance in [2, 12].

When X is a finite alphabet, a *non-empty finite word* over X is any sequence $x = a_1 \dots a_k$, where $a_i \in X$ for $i = 1, \dots, k$, and k is an integer ≥ 1 . The *length* of x is k , denoted by $|x|$. The *empty word* has no letter and is denoted by λ ; its length is 0. X^* is the *set of finite words* (including the empty word) over X , and $X^+ = X^* \setminus \{\lambda\}$ is the set of *non-empty finite words*. A *finitary language*, called here also a *dictionary*, over the alphabet X is a subset of X^* .

An ω -word over X is an ω -sequence $a_1 \dots a_n \dots$, where for all integers $i \geq 1$, $a_i \in X$. When σ is an ω -word over X , we write $\sigma = \sigma(1)\sigma(2) \dots \sigma(n) \dots$, where for all i , $\sigma(i) \in X$, and $\sigma[n] = \sigma(1)\sigma(2) \dots \sigma(n)$ for all $n \geq 1$ and $\sigma[0] = \lambda$.

The usual concatenation product of two finite words u and v is denoted $u \cdot v$ (and sometimes just uv). This product is extended to the product of a finite word u and an ω -word v : the infinite word $u \cdot v$ is then the ω -word such that:

$$(u \cdot v)(k) = u(k) \text{ if } k \leq |u|, \text{ and } (u \cdot v)(k) = v(k - |u|) \text{ if } k > |u|.$$

The *prefix relation* is denoted \sqsubseteq : a finite word u is a *prefix* of a finite word v (respectively, an infinite word v), denoted $u \sqsubseteq v$, if and only if there exists a finite word w (respectively, an infinite word w), such that $v = u \cdot w$. The *set of ω -words* over the alphabet X is denoted by X^ω . An ω -language over an alphabet X is a subset of X^ω .

We shall denote $X^{\leq \omega} = X^* \cup X^\omega$ the set of *finite or infinite words* over the alphabet X .

We assume the reader to be familiar with basic notions of topology which may be found in [2, 13–16]. There is a natural metric on the set X^ω of infinite words over a finite alphabet X containing at least two letters. It is called the *prefix metric* and is defined as follows. For $u, v \in X^\omega$ and $u \neq v$ let $\delta(u, v) = 2^{-l_{\text{pref}(u,v)}}$ where $l_{\text{pref}(u,v)}$ is the first integer n such that the $(n+1)^{\text{st}}$ letter of u is different from the $(n+1)^{\text{st}}$ letter of v . This metric induces on X^ω the usual Cantor topology for which the *open subsets* of X^ω are of the form $W \cdot X^\omega$, where $W \subseteq X^*$. A set $L \subseteq X^\omega$ is a *closed set* iff its complement $X^\omega \setminus L$ is an open set. Define now the *Borel Hierarchy* of subsets of X^ω :

Definition 2.1 For a non-null countable ordinal α , the classes Σ_α^0 and Π_α^0 of the Borel Hierarchy on the topological space X^ω are defined as follows:

Σ_1^0 is the class of open subsets of X^ω , Π_1^0 is the class of closed subsets of X^ω ,

and for any countable ordinal $\alpha \geq 2$:

Σ_α^0 is the class of countable unions of subsets of X^ω in $\bigcup_{\gamma < \alpha} \Pi_\gamma^0$.

Π_α^0 is the class of countable intersections of subsets of X^ω in $\bigcup_{\gamma < \alpha} \Sigma_\gamma^0$.

For a countable ordinal α , a subset of Σ^ω is a Borel set of *rank* α iff it is in $\Sigma_\alpha^0 \cup \Pi_\alpha^0$ but not in $\bigcup_{\gamma < \alpha} (\Sigma_\gamma^0 \cup \Pi_\gamma^0)$.

There exists another hierarchy beyond the Borel hierarchy, which is called the projective hierarchy. The classes Σ_n^1 and Π_n^1 , for integers $n \geq 1$, of the projective hierarchy are obtained from the Borel hierarchy by successive applications of operations of projection and complementation. The first level of the projective hierarchy consists of the class of *analytic sets*, and the class of *co-analytic sets* which are complements of analytic sets. In particular, the class of Borel subsets of X^ω is strictly included in the class Σ_1^1 of *analytic sets*. The class of analytic sets is also the class of the continuous images of Borel sets.

We now recall the notion of Wadge reducibility, which will be fundamental in the sequel.

Definition 2.2 (Wadge [17]) Let X, Y be two finite alphabets. For $L \subseteq X^\omega$ and $L' \subseteq Y^\omega$, L is said to be Wadge reducible to L' ($L \leq_W L'$) iff there exists a continuous function $f : X^\omega \rightarrow Y^\omega$, such that $L = f^{-1}(L')$. L and L' are Wadge equivalent iff $L \leq_W L'$ and $L' \leq_W L$. This is denoted by $L \equiv_W L'$.

The relation \leq_W is reflexive and transitive, and \equiv_W is an equivalence relation.

The *equivalence classes* of \equiv_W are called *Wadge degrees*.

For $L \subseteq X^\omega$ and $L' \subseteq Y^\omega$, if $L \leq_W L'$ and f is a continuous function from X^ω into Y^ω with $L = f^{-1}(L')$, then f is called a continuous reduction of L to L' . Intuitively it means that L is less complicated than L' because to check whether $x \in L$ it suffices to check whether $f(x) \in L'$ where f is a continuous function.

Recall that each Borel class Σ_α^0 and Π_α^0 is closed under inverse images by continuous functions and that a set $L \subseteq X^\omega$ is a Σ_α^0 (respectively Π_α^0)-*complete set* iff for any set $L' \subseteq Y^\omega$, L' is in Σ_α^0 (respectively Π_α^0) iff $L' \leq_W L$.

There is a close relationship between Wadge reducibility and games that we now introduce.

Definition 2.3 Let $L \subseteq X^\omega$ and $L' \subseteq Y^\omega$. The Wadge game $W(L, L')$ is a game with perfect information between two players. Player 1 is in charge of L and Player 2 is in charge of L' .

Player 1 first writes a letter $a_1 \in X$, then Player 2 writes a letter $b_1 \in Y$, then Player 1 writes a letter $a_2 \in X$, and so on.

The two players alternatively write letters a_n of X for Player 1 and b_n of Y for Player 2.

After ω steps, Player 1 has written an ω -word $a \in X^\omega$ and Player 2 has written an ω -word $b \in Y^\omega$. Player 2 is allowed to skip, even infinitely often, provided he really writes an ω -word in ω steps.

Player 2 wins the play iff $[a \in L \leftrightarrow b \in L']$, i.e. iff :

$$[(a \in L \text{ and } b \in L') \text{ or } (a \notin L \text{ and } b \notin L' \text{ and } b \text{ is infinite})].$$

Recall that a strategy for Player 1 is a function $\sigma : (Y \cup \{s\})^* \rightarrow X$. And a strategy for Player 2 is a function $f : X^+ \rightarrow Y \cup \{s\}$.

A strategy σ is a winning strategy for Player 1 iff he always wins the play when he uses the strategy σ , i.e. when the n^{th} letter he writes is given by $a_n = \sigma(b_1 \cdots b_{n-1})$, where b_i is the letter written by Player 2 at the step i and $b_i = s$ if Player 2 skips at the step i . A winning strategy for Player 2 is defined in a similar manner.

Martin's Theorem states that every Gale-Stewart game $G(B)$, where B is a Borel set, is determined, see [15]. This implies the following determinacy result:

Theorem 2.4 (Wadge) *Let $L \subseteq X^\omega$ and $L' \subseteq Y^\omega$ be two Borel sets, where X and Y are finite alphabets. Then the Wadge game $W(L, L')$ is determined: one of the two players has a winning strategy. And $L \leq_W L'$ iff Player 2 has a winning strategy in the game $W(L, L')$.*

3 ω -powers and sets of dictionaries

Recall that, for $V \subseteq X^*$, the ω -language

$$V^\omega = \{u_1 \cdot u_2 \cdots u_n \cdots \mid \forall i \geq 1 u_i \in V \setminus \{\lambda\}\}$$

is the ω -power of the language, or dictionary, V .

A dictionary over the alphabet X may be seen as an element of the space 2^{X^*} , i.e. the set of functions from X^* into 2 , where $2 = \{0, 1\}$ is a two letter alphabet. The space 2^{X^*} is naturally equipped with the product topology of the discrete topology on $2 = \{0, 1\}$. The set X^* of finite words over the alphabet X is countable so there is a bijection between X^* and ω and the topological space 2^{X^*} is in fact homeomorphic to the Cantor space 2^ω . The notions of Borel and projective hierarchies on the space 2^{X^*} are obtained in the same way as above in the case of the Cantor space X^ω .

Lecomte introduced in [1] the following sets of dictionaries. For a non null countable ordinal ξ , we set

$$\mathcal{W}(\Sigma_\xi^0) := \{A \subseteq 2^* \mid A^\omega \text{ is a } \Sigma_\xi^0\text{-set}\},$$

$$\mathcal{W}(\Pi_\xi^0) := \{A \subseteq 2^* \mid A^\omega \text{ is a } \Pi_\xi^0\text{-set}\},$$

$$\mathcal{W}(\Delta_1^1) := \{A \subseteq 2^* \mid A^\omega \text{ is a Borel set}\}.$$

Lecomte proved in [1] that $\mathcal{W}(\Sigma_2^0)$ is in $\Sigma_2^1(2^{2^*}) \setminus \Pi_2^0$ and that all the other sets $\mathcal{W}(\Sigma_\xi^0)$, $\mathcal{W}(\Pi_\xi^0)$, and $\mathcal{W}(\Delta_1^1)$ are in $\Sigma_2^1(2^{2^*}) \setminus D_2(\Sigma_1^0)$, where $D_2(\Sigma_1^0)$ is the class of 2-differences of open sets, that is, the class of sets which are intersections of an open set and of a closed set. Finkel and Lecomte showed in [11] that for each countable ordinal $\xi \geq 3$ the sets $\mathcal{W}(\Sigma_\xi^0)$ and $\mathcal{W}(\Pi_\xi^0)$ are actually Π_1^1 -hard. This gives a much better lower bound on the complexity of these sets, but their complexity is not completely determined.

Staiger gave in [5] a characterization of the set $\mathcal{W}(\Sigma_1^0)$ (respectively, $\mathcal{W}(\Pi_1^0)$). He gave in [5] an example of a dictionary $V \in \mathcal{W}(\Sigma_1^0 \setminus \Pi_1^0)$, and also an example of a $W \in \mathcal{W}(\Delta_2^0) \setminus \mathcal{W}(\Sigma_1^0 \cup \Pi_1^0)$. We refer the reader to [10, 11] for an example of a $W \in \mathcal{W}(\Sigma_2^0 \setminus \Pi_2^0)$.

In this paper we show that the set $\mathcal{W}(\Delta_1^1)$ is more complex than the set $\mathcal{W}(\Sigma_2^0)$. As an application we improve the lower bound on the complexity of the set $\mathcal{W}(\Delta_1^1)$.

We have already mentioned in the introduction the existence of a dictionary L such that L^ω is Σ_1^1 -complete, and hence non Borel. We now give a simple construction of such a language L using the notion of substitution that we now recall, (see [6] for more details).

A *substitution* is defined by a mapping $f : X \longrightarrow \mathcal{P}(Y^*)$, where $X = \{a_1, \dots, a_n\}$ and Y are two finite alphabets. For each integer $i \in [1; n]$, $f(a_i) = L_i$ is a finitary language over the alphabet Y .

Now this mapping is extended in the usual manner to finite words: $f(a_{i_1} \cdots a_{i_n}) = L_{i_1} \cdots L_{i_n}$, and to finitary languages $L \subseteq X^*$: $f(L) = \cup_{x \in L} f(x)$.

If for each integer $i \in [1; n]$ the language L_i does not contain the empty word, then the mapping f may be extended to ω -words:

$$f(x(1) \cdots x(n) \cdots) = \{u_1 \cdots u_n \cdots \mid \forall i \geq 1 \quad u_i \in f(x(i))\}$$

and to ω -languages $L \subseteq X^\omega$ by setting $f(L) = \cup_{x \in L} f(x)$.

Now let $X = \{0, 1\}$, d be a new letter not in X , and

$$D = \{u \cdot d \cdot v \mid u, v \in X^* \text{ and } [(|v| = 2|u|) \text{ or } (|v| = 2|u| + 1)]\}$$

Let $g : X \longrightarrow \mathcal{P}((X \cup \{d\})^*)$ be the substitution defined by $g(a) = a \cdot D$.

Notice that if V^ω is an ω -power then $g(V^\omega) = (g(V))^\omega$ is also an ω -power.

If $W = 0^* \cdot 1$ then $W^\omega = (0^* \cdot 1)^\omega$ is the set of ω -words over the alphabet X containing infinitely many occurrences of the letter 1. It is a well known example of an ω -language which is a Π_2^0 -complete subset of X^ω . One can prove that $(g(W))^\omega$ is Σ_1^1 -complete, and hence a non Borel set. This is done by reducing to this ω -language a well-known example of a Σ_1^1 -complete set: the set of infinite binary trees labelled in the alphabet $\{0, 1\}$ having an infinite branch in the Π_2^0 -complete set $(0^* \cdot 1)^\omega$.

More generally it is proved in [6, proof of Theorem 4.5 and Section 5] that if $W^\omega \subseteq X^\omega$ is an ω -power which is Π_2^0 -hard, then the ω -power $(g(W))^\omega \subseteq (X \cup \{d\})^\omega$ is Σ_1^1 -complete, and hence non Borel.

We use this result to prove our first proposition. In the sequel, for two sets $A, B \subseteq 2^{X^*}$ we denote $A \leq B$ iff there is a continuous function $H : 2^{X^*} \rightarrow 2^{X^*}$ such that $A = H^{-1}(B)$. So the relation \leq is in fact the Wadge reducibility relation \leq_W .

Proposition 3.1 *The following relation holds : $\mathcal{W}(\Sigma_2^0) \leq \mathcal{W}(\Delta_1^1)$.*

Proof. We shall use the substitution g defined above. Then let $g' : X \cup \{d\} \rightarrow \mathcal{P}(X^*)$ be the substitution simply defined by $g'(0) = \{0 \cdot 1\}$, $g'(1) = \{0 \cdot 1^2\}$, and $g'(d) = \{0 \cdot 1^3\}$. And let $G = g' \circ g$ be the substitution obtained by the composition of g followed by g' . Then, for every dictionary $V \subseteq X^*$, the language $G(V)$ is also a dictionary over the alphabet X and $G(V^\omega) = (G(V))^\omega$. The substitution G will provide the reduction $G : 2^{X^*} \rightarrow 2^{X^*}$.

Firstly, it is easy to see that the mapping $G : 2^{X^*} \rightarrow 2^{X^*}$ is continuous, [13].

Secondly, we claim that for every dictionary $V \subseteq X^*$, it holds that:

$$V \in \mathcal{W}(\Sigma_2^0) \text{ if and only if } G(V) \in \mathcal{W}(\Delta_1^1).$$

Assume first that $V \notin \mathcal{W}(\Sigma_2^0)$. By definition of $\mathcal{W}(\Sigma_2^0)$ this means that V^ω is not a Σ_2^0 -subset of 2^ω . Then we can infer from Hurewicz's Theorem, see [15, page 160], that the ω -power V^ω is Π_2^0 -hard because it is an analytic subset of 2^ω which is not a Σ_2^0 -set. Then it follows from [6, proof of Theorem 4.5 and Section 5] that the ω -power $(g(V))^\omega \subseteq (X \cup \{d\})^\omega$ is Σ_1^1 -complete, and hence non Borel. It is now very easy to check, applying the second substitution g' , that the ω -power $(G(V))^\omega \subseteq X^\omega$ is also non Borel. This means that $G(V)$ does not belong to the set $\mathcal{W}(\Delta_1^1)$.

Conversely assume now that $V \in \mathcal{W}(\Sigma_2^0)$. By definition of $\mathcal{W}(\Sigma_2^0)$ this means that V^ω is a Σ_2^0 -subset of X^ω , i.e. is a countable union of closed sets $F_n \subseteq X^\omega$, $n \geq 1$. Thus $V^\omega = \bigcup_{n \geq 1} F_n$ and $G(V^\omega) = G(\bigcup_{n \geq 1} F_n) = \bigcup_{n \geq 1} G(F_n)$.

We are going to show that for every closed set $F \subseteq X^\omega$, it holds that $G(F)$ is a Borel subset of X^ω .

Let then $F \subseteq X^\omega$ be a closed set. Then there is a tree $T \subseteq X^*$ such that $F = [T]$, i.e. F is the set of the infinite branches of T . We first prove that $g(F)$ is Borel. For any ω -word $y \in (X \cup \{d\})^\omega$, it holds that $y \in g(F)$ if and only if there exist $x \in F$ and sequences $u_i, v_i \in X^*$, $i \geq 1$, such that :

$$y = x(1) \cdot (u_1 \cdot d \cdot v_1) \cdot x(2) \cdot (u_2 \cdot d \cdot v_2) \cdot x(3) \cdots$$

where for each integer $i \geq 1$, $(|v_i| = 2|u_i|)$ or $(|v_i| = 2|u_i| + 1)$.

Let then T_1 be the set of finite prefixes of such ω -words in the set $g(F)$. The set $T_1 \subseteq (X \cup \{d\})^*$ is a tree. We claim that $g(F) = [T_1] \cap (\{0, 1\}^* \cdot d)^\omega$.

The inclusion $g(F) \subseteq [T_1] \cap (\{0, 1\}^* \cdot d)^\omega$ is straightforward.

To prove the inverse inclusion, let us consider an ω -word $x \in [T_1] \cap (\{0, 1\}^* \cdot d)^\omega$.

Then for each integer $n \geq 1$ there exists (at least) one finite sequence $(\varepsilon_i)_{1 \leq i \leq n} \in \{0, 1\}^n$ and one finite word $a_1 \cdot a_2 \cdots a_n \in X^*$ and finite words u_i and v_i in X^* , for $1 \leq i \leq n - 1$, and $u \in X^*$, such that :

$$a_1 \cdot (u_1 \cdot d \cdot v_1) \cdot a_2 \cdot (u_2 \cdot d \cdot v_2) \cdots a_{n-1} (u_{n-1} \cdot d \cdot v_{n-1}) \cdot a_n \cdot u \cdot d \sqsubseteq x$$

where for each integer $i \in [1, n]$, $(|v_i| = 2|u_i|)$ iff $\varepsilon_i = 0$ and $(|v_i| = 2|u_i| + 1)$ iff $\varepsilon_i = 1$, and $a_1 \cdot a_2 \cdots a_{n-1} \in T$.

Consider now all the “suitable” sequences $(\varepsilon_i)_{1 \leq i \leq n} \in \{0, 1\}^n$ defined as above. The set of these suitable sequences is closed under prefix. Therefore this set form a subtree of $(\{0, 1\}^*, \sqsubseteq)$, which is finitely branching. This tree is infinite so by König’s Lemma it has an infinite branch. therefore there exists an *infinite* sequence $(\varepsilon_i)_{1 \leq i < \omega} \in \{0, 1\}^\omega$ and one infinite word $a_1 \cdot a_2 \cdots a_n \cdots \in \Sigma^\omega$ and finite words u_i and v_i in X^* , for $1 \leq i < \omega$, such that :

$$x = a_1 \cdot (u_1 \cdot d \cdot v_1) \cdot a_2 \cdot (u_2 \cdot d \cdot v_2) \cdots a_n (u_n \cdot d \cdot v_n) \cdots$$

where for each integer $i \geq 1$, $(|v_i| = 2|u_i|)$ iff $\varepsilon_i = 0$ and $(|v_i| = 2|u_i| + 1)$ iff $\varepsilon_i = 1$, and $a_1 \cdot a_2 \cdots a_n \cdots \in [T] = F$.

This shows that $x \in g(F)$.

Thus $g(F) = [T_1] \cap (\{0, 1\}^* \cdot d)^\omega$ is the intersection of the closed set $[T_1]$ and of the Π_2^0 -set $(\{0, 1\}^* \cdot d)^\omega$. Then $g(F)$ is a Borel Π_2^0 -set, and it is easy to see that $G(F)$ is also Borel.

Assume now that $V \in \mathcal{W}(\Sigma_2^0)$, then $V^\omega = \bigcup_{n \geq 1} F_n$, where $F_n \subseteq X^\omega$ are closed sets. Then $G(V)^\omega = G(V^\omega) = G(\bigcup_{n \geq 1} F_n) = \bigcup_{n \geq 1} G(F_n)$ is a Borel subset of X^ω , so $G(V)$ belongs to the set $\mathcal{W}(\Delta_1^1)$. \square

We can now improve the result : $\mathcal{W}(\Delta_1^1) \in \Sigma_2^1(2^{2^*}) \setminus D_2(\Sigma_1^0)$ proved in [1].

Corollary 3.2 *The following relation holds : $\mathcal{W}(\Delta_1^1) \in \Sigma_2^1(2^{2^*}) \setminus \Pi_2^0$*

Proof. It follows directly from the relations $\mathcal{W}(\Sigma_2^0) \in \Sigma_2^1(2^{2^*}) \setminus \Pi_2^0$ and $\mathcal{W}(\Delta_1^1) \in \Sigma_2^1(2^{2^*})$, proved by Lecomte in [1], and from Proposition 3.1 stating that $\mathcal{W}(\Sigma_2^0) \leq \mathcal{W}(\Delta_1^1)$. \square

Remark 3.3 We have obtained only a slight improvement of Lecomte’s result that $\mathcal{W}(\Delta_1^1) \in \Sigma_2^1(2^{2^*}) \setminus D_2(\Sigma_1^0)$. The question is still open of the exact complexity of the two sets $\mathcal{W}(\Delta_1^1)$ and $\mathcal{W}(\Sigma_2^0)$ (and also of the other sets $\mathcal{W}(\Sigma_k^0)$ and $\mathcal{W}(\Pi_k^0)$).

However, Proposition 3.1 could provide a better improvement of the lower bound on the complexity of $\mathcal{W}(\Delta_1^1)$ as soon as a better improvement of the lower bound on the complexity of $\mathcal{W}(\Sigma_2^0)$ would be obtained. On the other hand, if one could obtain a better upper bound on the complexity of the set $\mathcal{W}(\Delta_1^1)$, then this would provide, by Proposition 3.1, a better upper bound on the complexity of the set $\mathcal{W}(\Sigma_2^0)$.

We consider now Borel ω -powers. It has been proved in [7] that for each integer $n \geq 1$, there exist some ω -powers of (context-free) languages which are Π_n^0 -complete Borel sets. (We refer the reader for instance to [18] for definitions and properties of context-free languages). These results were obtained by the use of an operation $A \longrightarrow A^\approx$ over ω -languages which is a variant of Duparc’s operation of exponentiation $A \longrightarrow A^\sim$, [19].

We first recall the definition of the operation $A \longrightarrow A^\sim$. Notice that this operation is defined over sets of finite or infinite words, called *conciliating sets* in [19].

Definition 3.4 (Duparc [19]) Let X be a finite alphabet, $\leftarrow \notin X$, and x be a finite or infinite word over the alphabet $Y = X \cup \{\leftarrow\}$.

Then x^{\leftarrow} is inductively defined by:

$$\lambda^{\leftarrow} = \lambda,$$

and for a finite word $u \in (X \cup \{\leftarrow\})^*$:

$$(u \cdot a)^{\leftarrow} = u^{\leftarrow} \cdot a, \text{ if } a \in X,$$

$$(u \cdot \leftarrow)^{\leftarrow} = u^{\leftarrow} \text{ with its last letter removed if } |u^{\leftarrow}| > 0,$$

$$\text{i.e. } (u \cdot \leftarrow)^{\leftarrow} = u^{\leftarrow}(1) \cdot u^{\leftarrow}(2) \cdots u^{\leftarrow}(|u^{\leftarrow}| - 1) \text{ if } |u^{\leftarrow}| > 0,$$

$$(u \cdot \leftarrow)^{\leftarrow} = \lambda \text{ if } |u^{\leftarrow}| = 0,$$

and for u infinite:

$(u)^{\leftarrow} = \lim_{n \in \omega} (u[n])^{\leftarrow}$, where, given β_n and v in X^* ,

$v \sqsubseteq \lim_{n \in \omega} \beta_n \leftrightarrow \exists n \forall p \geq n \quad \beta_p \upharpoonright [v] = v$.

(The finite or infinite word $\lim_{n \in \omega} \beta_n$ is determined by the set of its (finite) prefixes).

Remark 3.5 For $x \in Y^{\leq \omega}$, x^{\leftarrow} denotes the string x , once every \leftarrow occurring in x has been “evaluated” to the back space operation, proceeding from left to right inside x . In other words $x^{\leftarrow} = x$ from which every interval of the form “ $a \leftarrow$ ” ($a \in X$) is removed. The letter \leftarrow may be called an “eraser”.

For example if $u = (a \leftarrow)^n$, for n an integer ≥ 1 , or $u = (a \leftarrow)^\omega$, or $u = (a \leftarrow \leftarrow)^\omega$, then $(u)^{\leftarrow} = \lambda$. If $u = (ab \leftarrow)^\omega$ then $(u)^{\leftarrow} = a^\omega$ and if $u = bb(\leftarrow a)^\omega$ then $(u)^{\leftarrow} = b$.

Let us notice that in Definition 3.4 the limit is not defined in the usual way:

for example if $u = bb(\leftarrow a)^\omega$ the finite word $u[n]^{\leftarrow}$ is alternatively equal to b or to ba : more precisely $u[2n+1]^{\leftarrow} = b$ and $u[2n+2]^{\leftarrow} = ba$ for every integer $n \geq 1$ (it holds also that $u[1]^{\leftarrow} = b$ and $u[2]^{\leftarrow} = bb$). Thus Definition 3.4 implies that $\lim_{n \in \omega} (u[n])^{\leftarrow} = b$ so $u^{\leftarrow} = b$.

We can now define the operation $A \longrightarrow A^\sim$ of *exponentiation of conciliating sets*:

Definition 3.6 (Duparc [19]) For $A \subseteq X^{\leq \omega}$ and $\leftarrow \notin X$, let

$$A^\sim =_{df} \{x \in (X \cup \{\leftarrow\})^{\leq \omega} \mid x^{\leftarrow} \in A\}.$$

We now define the variant $A \longrightarrow A^\approx$ of the operation $A \longrightarrow A^\sim$.

Definition 3.7 ([7]) Let X be a finite alphabet, $\leftarrow \notin X$, and x be a finite or infinite word over the alphabet $Y = X \cup \{\leftarrow\}$.

Then x^{\leftrightarrow} is inductively defined by:

$\lambda^{\leftrightarrow} = \lambda$,

and for a finite word $u \in (X \cup \{\leftarrow\})^*$:

$(u \cdot a)^{\leftrightarrow} = u^{\leftrightarrow} \cdot a$, if $a \in X$,

$(u \leftarrow)^{\leftrightarrow} = u^{\leftrightarrow}$ with its last letter removed if $|u^{\leftrightarrow}| > 0$,

$(u \leftarrow)^{\leftrightarrow}$ is undefined if $|u^{\leftrightarrow}| = 0$,

and for u infinite:

$(u)^{\leftrightarrow} = \lim_{n \in \omega} (u[n])^{\leftrightarrow}$, where, given β_n and v in X^* ,

$v \sqsubseteq \lim_{n \in \omega} \beta_n \leftrightarrow \exists n \forall p \geq n \quad \beta_p \upharpoonright [v] = v$.

The difference between the definitions of x^{\leftarrow} and x^{\leftrightarrow} is that here we have added the convention that $(u \leftarrow)^{\leftrightarrow}$ is undefined if $|u^{\leftrightarrow}| = 0$, i.e. when the last letter \leftarrow can not be used as an eraser (because every letter of X in u has already been erased by some erasers \leftarrow placed in u). For example if $u = \leftarrow (a \leftarrow)^\omega$ or $u = a \leftarrow \leftarrow a^\omega$ or $u = (a \leftarrow \leftarrow)^\omega$, then $(u)^{\leftrightarrow}$ is undefined.

Definition 3.8 For $A \subseteq X^{\leq \omega}$, $A^\approx = \{x \in (X \cup \{\leftarrow\})^{\leq \omega} \mid x^{\leftrightarrow} \in A\}$.

The operation $A \longrightarrow A^\approx$ was used by Duparc in his study of the Wadge hierarchy, [19]. The result stated in the following lemma will be important in the sequel.

Lemma 3.9 Let X be a finite alphabet and $L \subseteq X^\omega$. Then the two ω -languages L^\sim and L^\approx are Wadge equivalent, i.e. $L^\sim \equiv_W L^\approx$.

Proof. Let X be a finite alphabet and $L \subseteq X^\omega$. We are going to prove that $L^\sim \equiv_W L^\approx$, using Wadge games.

- a) In the Wadge game $W(L^\sim, L^\approx)$ the player in charge of L^\approx has clearly a winning strategy which consists in copying the play of the other player except if player 1 writes the eraser \leftarrow but he has nothing to erase. In this case player 2 writes for example a letter $a \in X$ and the eraser \leftarrow at the next step of the play. Now if, in ω steps, player 1 has written the ω -word α and player 2 has written the ω -word β , then it is easy to see that $[\alpha^{\leftarrow} = \beta^{\leftarrow}]$ and then $\alpha \in L^\sim$ iff $\beta \in L^\approx$. Thus player 2 has a winning strategy in the Wadge game $W(L^\sim, L^\approx)$.

- b) Consider now the Wadge game $W(L^\approx, L^\sim)$. The only extra possibility for player 1 in charge of L^\approx is to get out of the set L^\approx by writing the eraser \leftarrow when in fact there is no letter of his previous play to erase. But then his final play is surely outside L^\approx . If this happens at some point of the play, then player 2 may write the eraser \leftarrow forever. Then, after ω steps, player 2 has written an infinite word β such that $\beta^{\leftarrow} = \lambda$. In particular, $\beta^{\leftarrow} \notin L$ because β^{\leftarrow} is not an infinite word, and $\beta \notin L^\sim$. On the other hand player 1 has written an infinite word α such that α^{\leftarrow} is undefined, hence $\alpha \notin L^\approx$. Therefore player 2 wins the play in this case too, and player 2 has a winning strategy in the Wadge game $W(L^\approx, L^\sim)$. □

The operation $A \rightarrow A^\approx$ is very useful in the study of ω -powers because it can be defined with the notion of substitution and preserves the ω -powers of finitary languages. Let $L_1 = \{w \in (X \cup \{\leftarrow\})^* \mid w^{\leftarrow} = \lambda\}$. L_1 is a context free (finitary) language generated by the context free grammar with the following productions: $(S, aS \leftarrow S)$ with $a \in X$; and (S, λ) .

Then, for each ω -language $A \subseteq X^\omega$, the ω -language $A^\approx \subseteq (X \cup \{\leftarrow\})^\omega$ is obtained by substituting in A the language $L_1 \cdot a$ for each letter $a \in X$. This implies that the operation $A \rightarrow A^\approx$ preserves the ω -powers of finitary languages. This is stated in the following lemma.

Lemma 3.10 ([7]) *Let X be a finite alphabet and let h be the substitution defined by $h(a) = L_1 \cdot a$ for every letter $a \in X$.*

If $A = V^\omega$ for some language $V \subseteq X^$, then $A^\approx = h(V^\omega) = (h(V))^\omega$. Thus, if A is an ω -power, then A^\approx is also an ω -power.*

We now recall the operation $A \rightarrow A^b$ used by Duparc in his study of the Wadge hierarchy, [19]. For $A \subseteq X^{\leq \omega}$ and b a letter not in X , A^b is the ω -language over $X \cup \{b\}$ which is defined by :

$$A^b = \{x \in (X \cup \{b\})^\omega \mid x(/b) \in A\}$$

where $x(/b)$ is the sequence obtained from x when removing every occurrence of the letter b .

We can now state the following lemma.

Lemma 3.11 *Let X be a finite alphabet having at least two elements and $A \subseteq X^\omega$.*

1. *For each integer $k \geq 2$, A is a Π_k^0 -subset of X^ω iff A^b is a Π_k^0 -subset of $(X \cup \{b\})^\omega$.*
2. *For each integer $k \geq 3$, A is a Σ_k^0 -subset of X^ω iff A^b is a Σ_k^0 -subset of $(X \cup \{b\})^\omega$.*

Proof. We denote by Z^∞ the set of infinite words in $(X \cup \{b\})^\omega$ having infinitely many letters in X . The set $Z^\infty = \{x \in (X \cup \{b\})^\omega \mid x(/b) \in X^\omega\}$ is a well known example of Π_2^0 -subset of $(X \cup \{b\})^\omega$, [15, 16]. Notice that Z^∞ , equipped with the induced topology, is a topological subspace of the Cantor space $(X \cup \{b\})^\omega$. One can define the Borel hierarchy on the topological space Z^∞ as in the case of the Cantor space, see [15, page 68]. Then one can prove by induction that, for each non-null countable ordinal α , the Σ_α^0 (respectively, Π_α^0)-subsets of Z^∞ are the intersections of Σ_α^0 (respectively, Π_α^0)-subsets of $(X \cup \{b\})^\omega$ with the set Z^∞ , see [15, page 167].

Let now ϕ be the function from Z^∞ into X^ω defined by $\phi(x) = x(/b)$. It is easy to see that, for each $A \subseteq X^\omega$, it holds that $\phi^{-1}(A) = A^b$. On the other hand, the function ϕ is continuous. Thus the inverse image of an open (respectively, closed) subset of X^ω is an open (respectively, closed) subset of Z^∞ . And one can prove by induction that, for each non-null countable ordinal α , the inverse image of a Σ_α^0 (respectively, Π_α^0)-subset of X^ω is a Σ_α^0 (respectively, Π_α^0)-subset of Z^∞ , i.e. the intersection of a Σ_α^0 (respectively, Π_α^0)-subset of $(X \cup \{b\})^\omega$ with the set Z^∞ .

Let now $k \geq 2$ and $A \subseteq X^\omega$ be a Π_k^0 -subset of X^ω . Then $\phi^{-1}(A) = A^b$ is a Π_k^0 -subset of Z^∞ , i.e. the intersection of a Π_k^0 -subset of $(X \cup \{b\})^\omega$ with the set Z^∞ . But Z^∞ is a Π_2^0 -subset of $(X \cup \{b\})^\omega$ thus $\phi^{-1}(A) = A^b$ is the intersection of two Π_k^0 -subsets of $(X \cup \{b\})^\omega$, hence also a Π_k^0 -subset of $(X \cup \{b\})^\omega$.

In a similar way we prove that if $k \geq 3$ and $A \subseteq X^\omega$ is a Σ_k^0 -subset of X^ω , then $\phi^{-1}(A) = A^b$ is a Σ_k^0 -subset of $(X \cup \{b\})^\omega$.

Conversely assume that for some integer $k \geq 2$ and $A \subseteq X^\omega$ the set A^b is a Π_k^0 -subset of $(X \cup \{b\})^\omega$. Notice that X^ω is a closed subset of $(X \cup \{b\})^\omega$. Thus $A = A^b \cap X^\omega$ is the intersection of two Π_k^0 -subsets of $(X \cup \{b\})^\omega$, hence also a Π_k^0 -subset of $(X \cup \{b\})^\omega$. And $A = A \cap X^\omega$ so A is also a Π_k^0 -subset of X^ω .

In a similar way we prove that if for some integer $k \geq 3$ and $A \subseteq X^\omega$ the set A^b is a Σ_k^0 -subset of $(X \cup \{b\})^\omega$, then A is also a Σ_k^0 -subset of X^ω . \square

Lemma 3.12 *Let X be a finite alphabet having at least two elements and $A \subseteq X^\omega$.*

1. *For each integer $k \geq 3$, A is a Σ_k^0 -subset of X^ω iff A^\approx is a Σ_{k+1}^0 -subset of $(X \cup \{\leftarrow\})^\omega$.*
2. *For each integer $k \geq 2$, A is a Π_k^0 -subset of X^ω iff A^\approx is a Π_{k+1}^0 -subset of $(X \cup \{\leftarrow\})^\omega$.*

Proof. Let X be a finite alphabet having at least two elements, $A \subseteq X^\omega$, and $k \geq 3$ be an integer. Then the following equivalences hold:

- $A \in \Sigma_k^0$
- $\longleftrightarrow A^b \in \Sigma_k^0$ by Lemma 3.11.
- $\longleftrightarrow A^b \leq_W B^b$ for some $B \subseteq X^{\leq\omega}$ such that B^b is Σ_k^0 -complete.
- $\longleftrightarrow (A^\sim)^b \leq_W (B^\sim)^b$ by [19, Proposition 23].
- $\longleftrightarrow (A^\sim)^b \in \Sigma_{k+1}^0$, because $(B^\sim)^b$ is Σ_{k+1}^0 -complete by [19, Lemma 31].
- $\longleftrightarrow A^\sim \in \Sigma_{k+1}^0$ by Lemma 3.11.
- $\longleftrightarrow A^\approx \in \Sigma_{k+1}^0$ by Lemma 3.9.

In a very similar way we prove that if $k \geq 2$ is an integer, then $A \in \Pi_k^0$ iff $A^\approx \in \Pi_{k+1}^0$. \square

We now state the following result about the classes $\mathcal{W}(\Pi_k^0)$.

Proposition 3.13 *For each integer $k \geq 2$ it holds that: $\mathcal{W}(\Pi_k^0) \leq \mathcal{W}(\Pi_{k+1}^0)$.*

Proof. We shall use the substitution h defined above.

Let then $h' : \{0, 1, \leftarrow\} \rightarrow \mathcal{P}(\{0, 1\}^*)$ be the substitution simply defined by $h'(0) = \{0 \cdot 1\}$, $h'(1) = \{0 \cdot 1^2\}$, and $h'(\leftarrow) = \{0 \cdot 1^3\}$. And let $H = h' \circ h$ be the substitution obtained by the composition of h followed by h' . Then, for every dictionary $V \subseteq X^* = \{0, 1\}^*$, the language $H(V)$ is also a dictionary over the alphabet X and $H(V^\omega) = (H(V))^\omega$. The substitution H will provide the reduction $H : 2^{X^*} \rightarrow 2^{X^*}$.

It is easy to see that the mapping $H : 2^{X^*} \rightarrow 2^{X^*}$ is continuous, [13].

We claim that for every dictionary $V \subseteq X^*$, it holds that $V \in \mathcal{W}(\Pi_k^0)$ if and only if $H(V) \in \mathcal{W}(\Pi_{k+1}^0)$.

Firstly by definition of the class $\mathcal{W}(\Pi_k^0)$ it holds that for every dictionary $V \subseteq X^*$, V is in the class $\mathcal{W}(\Pi_k^0)$ iff V^ω is a Π_k^0 -set. By Lemma 3.12, V^ω is a Π_k^0 -set iff $(V^\omega)^\approx$ is a Π_{k+1}^0 -set. But $(V^\omega)^\approx = h(V^\omega) = h(V)^\omega$. Thus V is in the class $\mathcal{W}(\Pi_k^0)$ iff $h(V)^\omega$ is in the class Π_{k+1}^0 . It is now easy to see, using the coding h' that this is equivalent to the assertion “ $(H(V))^\omega$ is in the class Π_{k+1}^0 ”, i.e. $H(V)$ is in the class $\mathcal{W}(\Pi_{k+1}^0)$. \square

In a very similar manner, we can prove the following result about the sets $\mathcal{W}(\Sigma_k^0)$ for integers $k \geq 3$.

Proposition 3.14 *For each integer $k \geq 3$ it holds that: $\mathcal{W}(\Sigma_k^0) \leq \mathcal{W}(\Sigma_{k+1}^0)$.*

Remark 3.15 Notice that here $k \geq 3$ because for $L \subseteq X^\omega$ then L may be in the class Σ_2^0 while $L^b \subseteq (X \cup \{b\})^\omega$ is not in the class Σ_2^0 . For instance $L = \{0, 1\}^\omega \subseteq \{0, 1\}^\omega$ is open and closed hence also in the class Σ_2^0 . But the ω -language L^b is simply the set of ω -words over the alphabet $\{0, 1, b\}$ which contain infinitely many letters 0 or 1 and it is a Π_2^0 -complete, hence non Σ_2^0 , subset of $\{0, 1, b\}^\omega$.

4 Concluding remarks

Lecomte proved that for every countable ordinal $\xi \geq 2$ (respectively, $\xi \geq 3$), $\mathcal{W}(\Pi_\xi^0) \in \Sigma_2^1(2^{2^*}) \setminus D_2(\Sigma_1^0)$ (respectively, $\mathcal{W}(\Sigma_\xi^0) \in \Sigma_2^1(2^{2^*}) \setminus D_2(\Sigma_1^0)$). Finkel and Lecomte proved that for every countable ordinal $\xi \geq 3$, the sets $\mathcal{W}(\Pi_\xi^0)$ and $\mathcal{W}(\Sigma_\xi^0)$ are actually Π_1^1 -hard. The exact complexity of the sets $\mathcal{W}(\Pi_\xi^0)$ and $\mathcal{W}(\Sigma_\xi^0)$ is still unknown, but our new results could help to determine it.

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