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## NEW PERMUTATION CODING AND EQUIDISTRIBUTION OF SET-VALUED STATISTICS

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ABSTRACT. A new coding for permutations is explicitly constructed and its association with the classical Lehmer coding provides a bijection of the symmetric group onto itself serving to show that six bivariable set-valued statistics are equidistributed on that group. This extends a recent result due to Cori valid for integer-valued statistics.

### 1. Introduction

In a recent paper Cori [Cor08] updates a classical algorithm constructed by Ossona de Mendez and Rosenstiehl [OR04] that provides a one-to-one correspondence between rooted hypermaps and indecomposable permutations. He further constructs a bijection of the symmetric group  $\mathfrak{S}_n$  onto itself that maps each permutation having  $p$  cycles and  $q$  left-to-right maxima onto another one having  $q$  cycles and  $p$  left-to-right maxima. Moreover, by using an encoding of permutations by Dyck paths due to Roblet and Viennot [RV96] he also shows that three bivariable *integer-valued* statistics, introduced in the next paragraph, are equidistributed on  $\mathfrak{S}_n$ . The purpose of this paper is to show that all those results can be extended to *set-valued* statistics and that the construction of the underlying bijection is very simple; it involves two permutation codings called the *A-code* and the *B-code*.

The first one, classically referred to as the *Lehmer code* [Le60] or the *inversion table*, goes back, in fact, to more ancient authors (Rothe, Rodrigues, Netto), as knowledgeably stated by Knuth ([Kn98], Ex. 5.1.1-7, p. 14). The second one is a *new* coding that takes the cycle decomposition of permutations into account. Although the motivation of the paper was to prove the equidistribution of several set-valued statistics, its novelty is to fully describe that B-code and exploit its basic properties.

The set-valued statistics in question are introduced as follows. Let  $w = x_1x_2 \cdots x_n$  be a word of length  $n$ , whose letters are positive integers. The **L**eft to right **m**aximum **p**lace set,  $\text{Lmap } w$ , of  $w$  is defined to be the set of all *places*  $i$  such that  $x_j < x_i$  for all  $j < i$ , while the **R**ight to left **m**inimum **l**etter set,  $\text{Rmil } w$ , of  $w$  is the set of all *letters*  $x_i$  such that  $x_j > x_i$  for all  $j > i$ .

When the word  $w$  is a permutation of  $12 \cdots n$  that we shall preferably denote by  $\sigma = \sigma(1)\sigma(2) \cdots \sigma(n)$  and the bijection  $i \mapsto \sigma(i)$  ( $1 \leq i \leq n$ ) has

$r$  disjoint cycles, whose *minimum* elements are  $c_1, c_2, \dots, c_r$ , respectively, define  $\text{Cyc } \sigma$  to be the set

$$\text{Cyc } \sigma := \{c_1, c_2, \dots, c_r\}.$$

When  $\sigma$  is a permutation, the *cardinalities* of  $\text{Lmap } \sigma$ ,  $\text{Rmil } \sigma$  and  $\text{Cyc } \sigma$  are denoted by  $\text{lmap } \sigma$ ,  $\text{rmil } \sigma$  and  $\text{cyc } \sigma$ , respectively, and classically referred to as the *number of left-to-right maxima*, *number of right-to-left minima*, *number of cycles*.

In Fig. 1 the graphs of the permutation  $\sigma = 5, 7, 1, 4, 9, 2, 6, 3, 8$  and its inverse  $\sigma^{-1} = 3, 6, 8, 4, 1, 7, 2, 9, 5$  have been drawn. The set  $\text{Lmap } \sigma$  (resp.  $\text{Lmap } \sigma^{-1}$ ) is the set of the *abscissas* of the “bullets,” while  $\text{Rmil } \sigma$  (resp.  $\text{Rmil } \sigma^{-1}$ ) is the set of the *ordinates* of the “crosses.” The set-valued statistics “Leh,” “RmilLeh” and “MaxLeh” will be further introduced. Notice that  $\text{lmap } \sigma = \text{rmil } \sigma^{-1} = 3$ ,  $\text{rmil } \sigma = \text{lmap } \sigma^{-1} = 4$ . As  $\sigma$  is the product of the disjoint cycles  $(15983)(4)(276)$ , we have  $\text{Cyc } \sigma = \text{Cyc } \sigma^{-1} = \{1, 2, 4\}$  and  $\text{cyc } \sigma = \text{cyc } \sigma^{-1} = 3$ .

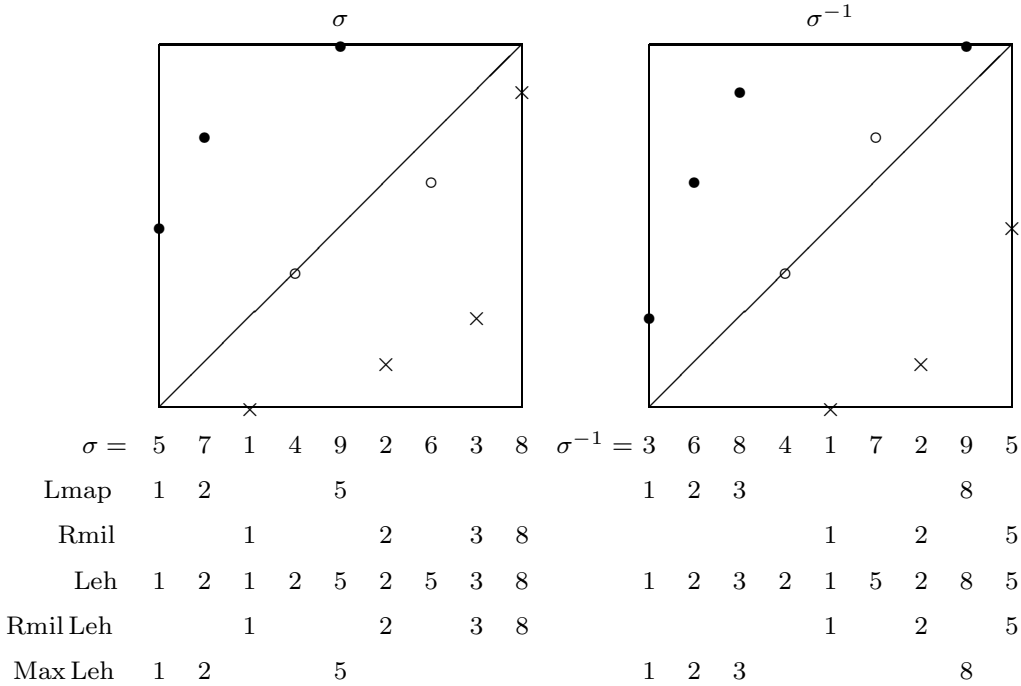


Fig. 1. Graphs of  $\sigma$  and of its inverse  $\sigma^{-1}$

First, recall Cori’s result [Cor08].

*The three pairs of integer-valued statistics  $(\text{rmil}, \text{cyc})$ ,  $(\text{cyc}, \text{rmil})$  and  $(\text{lmap}, \text{rmil})$  are equidistributed on  $\mathfrak{S}_n$ .*

The equidistribution of the first two pairs (resp. of the last two ones) is proved by updating the Ossona-de-Mendez-Rosenstiehl algorithm [OR04]

on hypermaps (resp. by using the Roblet-Viennot Dyck path encoding [RV96]). Second, the set-valued statistics “Cyc” and “Rmil” (or “Lmap”) are known to be equidistributed on  $\mathfrak{S}_n$ . This is one of the properties of the first fundamental transformation [Lo83, chap. 10]. Our main result is the following theorem.

**Theorem 1.** *The six bivariable set-valued statistics (Cyc, Rmil), (Cyc, Lmap), (Rmil, Lmap), (Rmil, Cyc), (Lmap, Rmil), (Lmap, Cyc) are all equidistributed on  $\mathfrak{S}_n$ .*

Based on two permutation codings, the A-code and B-code, introduced in Sections 2 and 3, respectively, we construct a bijection  $\phi$  of  $\mathfrak{S}_n$  onto itself (see (4.1)) having the following property:

$$(1.1) \quad (\text{Lmap, Rmil}) \sigma = (\text{Lmap, Cyc}) \phi(\sigma) \quad (\sigma \in \mathfrak{S}_n).$$

Let  $\mathbf{i} : \sigma \mapsto \sigma^{-1}$ . As

$$(1.2) \quad \text{Cyc } \mathbf{i} \sigma = \text{Cyc } \sigma;$$

$$(1.3) \quad \text{Rmil } \mathbf{i} \sigma = \text{Lmap } \sigma;$$

(see Fig. 1 for the second relation), it follows from (1.1) that the chain

$$(1.4) \quad \begin{array}{ccccccccc} \mathfrak{S}_n & \xrightarrow{\mathbf{i}} & \mathfrak{S}_n & \xrightarrow{\phi^{-1}} & \mathfrak{S}_n & \xrightarrow{\mathbf{i}} & \mathfrak{S}_n & \xrightarrow{\phi} & \mathfrak{S}_n & \xrightarrow{\mathbf{i}} & \mathfrak{S}_n \\ (\text{Cyc}) & & (\text{Cyc}) & & (\text{Rmil}) & & (\text{Lmap}) & & (\text{Lmap}) & & (\text{Rmil}) \\ (\text{Rmil}) & & (\text{Lmap}) & & (\text{Lmap}) & & (\text{Rmil}) & & (\text{Cyc}) & & (\text{Cyc}) \end{array}$$

provides all the bijections needed to prove Theorem 1. Note that (1.1), on the one hand, and (1.2)–(1.3), on the other hand, are reproduced as

$$\begin{array}{ccc} \mathfrak{S}_n & \xrightarrow{\phi} & \mathfrak{S}_n & \text{and} & \mathfrak{S}_n & \xrightarrow{\mathbf{i}} & \mathfrak{S}_n \\ (\text{Lmap}) & & (\text{Lmap}) & & (\text{Cyc}) & & (\text{Cyc}) \\ (\text{Rmil}) & & (\text{Cyc}) & & (\text{Rmil}) & & (\text{Lmap}) \end{array}$$

Let  $A = (I_1, I_2, \dots, I_h)$  be an ordered partition of the set  $[n] := \{1, 2, \dots, n\}$  into disjoint non-empty *intervals*, such that  $\max I_j + 1 = \min I_{j+1}$  for  $j = 1, 2, \dots, h - 1$ . A permutation  $\sigma$  from  $\mathfrak{S}_n$  is said to be *A-decomposable*, if each  $I_j$  is the *smallest interval* such that  $\sigma(I_j) = I_j$  (see [Com74], p. 261, exercise 14). For instance,  $\sigma = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 2 & 1 & 5 & 4 & 3 \end{pmatrix}$  is *A-decomposable*, with  $A = (\{1, 2\}, \{3, 4, 5\})$ . It is convenient to write  $\text{Decomp } \sigma = A$ , if  $\sigma$  is *A-decomposable*. A permutation is said to be *indecomposable*, if it is *A-decomposable*, with  $A = ([n])$ . The bijection  $\phi$  defined in (4.1) has the further property

$$(1.5) \quad \text{Decomp } \phi(\sigma) = \text{Decomp } \sigma \quad (\sigma \in \mathfrak{S}_n).$$

As we evidently have

$$(1.6) \quad \text{Decomp } \mathbf{i} \sigma = \text{Decomp } \sigma,$$

the following result holds.

**Theorem 2.** *Let  $A$  be an ordered partition of the set  $[n]$  into disjoint consecutive non-empty intervals. Then,  $(\text{Cyc}, \text{Rmil})$ ,  $(\text{Cyc}, \text{Lmap})$ ,  $(\text{Rmil}, \text{Lmap})$ ,  $(\text{Rmil}, \text{Cyc})$ ,  $(\text{Lmap}, \text{Rmil})$ ,  $(\text{Lmap}, \text{Cyc})$  are equidistributed on the set of all  $A$ -decomposable permutations from  $\mathfrak{S}_n$ .*

The next corollary is relevant to the study of hypermaps, as the set of rooted hypermaps with darts  $1, 2, \dots, n$  is in one-to-one correspondence with the subset of indecomposable permutations from  $\mathfrak{S}_{n+1}$  (see [Cor08, CM92]).

**Corollary 3.** *The statistics  $(\text{Cyc}, \text{Rmil})$ ,  $(\text{Cyc}, \text{Lmap})$ ,  $(\text{Rmil}, \text{Lmap})$ ,  $(\text{Rmil}, \text{Cyc})$ ,  $(\text{Lmap}, \text{Rmil})$ ,  $(\text{Lmap}, \text{Cyc})$  are equidistributed on the set of all indecomposable permutations from  $\mathfrak{S}_n$ .*

The construction of the bijection  $\phi$  together with the proofs of Theorem 2, and Corollary 3 are given in Section 4. It is followed by the algorithmic definitions of both A-code and B-code in Section 5. Tables and concluding remarks are reproduced in Section 6.

## 2. The A-code

The *Lehmer code* [Le60] of a permutation  $\sigma = \sigma(1)\sigma(2)\cdots\sigma(n)$  of  $12\cdots n$  is defined to be the sequence  $\text{Leh } w = (a_1, a_2, \dots, a_n)$ , where for each  $i = 1, 2, \dots, n$

$$a_i := \#\{j : 1 \leq j \leq i, \sigma(j) \leq \sigma(i)\}.$$

The sequence  $\text{Leh } w$  belongs to  $\text{SE}_n$  of all sequences  $a = (a_1, a_2, \dots, a_n)$ , called *subexcedant*, such that  $1 \leq a_i \leq i$  for each  $i = 1, 2, \dots, n$ . For such a sequence it makes sense to define the set, denoted by  $\text{Max } a$ , of all letters (or places!)  $a_i$  such that  $a_i = i$ .

Under the graphs drawn in Fig. 1 the Lehmer codes  $\text{Leh } \sigma$  and  $\text{Leh } \sigma^{-1}$  have been calculated, as well as the four sets  $\text{Rmil } \text{Leh } \sigma$ ,  $\text{Rmil } \text{Leh } \sigma^{-1}$ ,  $\text{Max } \text{Leh } \sigma$  and  $\text{Max } \text{Leh } \sigma^{-1}$ . The next Proposition is geometrically evident and given without proof. It shows that the set-valued statistics “Lmap” and “Rmip” can be directly read from the Lehmer code.

**Proposition 4.** *For each permutation  $\sigma$  we have:*

$$(2.1) \quad \text{Rmil } \text{Leh } \sigma = \text{Rmil } \sigma;$$

$$(2.2) \quad \text{Max } \text{Leh } \sigma = \text{Lmap } \sigma.$$

We then define the *A-code* of a permutation  $\sigma$  to be

$$(2.3) \quad \text{A-code } \sigma := \text{Leh } \mathbf{i} \sigma.$$

Hence,  $\text{Max } \text{A-code } \sigma = \text{Max } \text{Leh } \mathbf{i} \sigma = \text{Lmap } \mathbf{i} \sigma = \text{Rmil } \sigma$ . Furthermore,  $\text{Rmil } \text{A-code } \sigma = \text{Rmil } \text{Leh } \mathbf{i} \sigma = \text{Rmil } \mathbf{i} \sigma = \text{Lmap } \sigma$ . As  $\text{Leh}$  is a bijection of the symmetric group  $\mathfrak{S}_n$  onto  $\text{SE}_n$ , we obtain the following result.

**Theorem 5.** *The A-code is a bijection of  $\mathfrak{S}_n$  onto  $SE_n$  having the property:*

$$(2.4) \quad (\text{Rmil}, \text{Lmap}) \sigma = (\text{Max}, \text{Rmil}) \text{ A-code } \sigma \quad (\sigma \in \mathfrak{S}_n).$$

An algorithmic definition of the A-code will be given in Section 5.

### 3. The B-code

The B-code is based on the decomposition of each permutation as product of disjoint cycles. For a permutation  $\sigma = \sigma(1)\sigma(2)\cdots\sigma(n)$  and each  $i = 1, 2, \dots, n$  let  $k := k(i)$  be the *smallest* integer  $k \geq 1$  such that  $\sigma^{-k}(i) \leq i$ . Then, define

$$\text{B-code } \sigma = (b_1, b_2, \dots, b_n) \quad \text{with} \quad b_i := \sigma^{-k(i)}(i) \quad (1 \leq i \leq n).$$

For example, with the permutation  $\sigma = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ 3 & 4 & 5 & 2 & 6 & 1 \end{pmatrix}$  we have:  
 $\sigma^{-1}(1) = 6, \sigma^{-2}(1) = 5, \sigma^{-3}(1) = 3, \sigma^{-4}(1) = 1$ , so that  $b_1 = 1$ ;  
 $\sigma^{-1}(2) = 4, \sigma^{-2}(2) = 2$ , so that  $b_2 = 2$ ;  $\sigma^{-1}(3) = 1$ , so that  $b_3 = 1$ ;  
 $\sigma^{-1}(4) = 2$ , so that  $b_4 = 2$ ;  $\sigma^{-1}(5) = 3$ , so that  $b_5 = 3$ ;  
 $\sigma^{-1}(6) = 5$ , so that  $b_6 = 5$ . Thus, B-code  $\sigma = (1, 2, 1, 2, 3, 5)$ .

An alternate definition is the following. First, the B-code of the unique permutation from  $\mathfrak{S}_1$  is defined to be the sequence  $(1) \in SE_1$ . Let  $n \geq 2$ . When writing each permutation  $\sigma \in \mathfrak{S}_n$  of order  $n \geq 2$  as a product of its disjoint cycles, the removal of  $n$  yields a permutation  $\sigma'$  of order  $(n - 1)$ . Let  $b' = (b'_1, b'_2, \dots, b'_{n-1})$  be the B-code of  $\sigma'$ . We define the B-code of  $\sigma$  to be  $b := (b'_1, b'_2, \dots, b'_{n-1}, \sigma^{-1}(n))$ . By induction on  $n$ , we immediately see that the B-code is a bijection of  $\mathfrak{S}_n$  onto  $SE_n$ .

The following Theorem shows that the set-valued statistics “Lmap” and “Cyc” can be directly read from the B-code.

**Theorem 6.** *The B-code is a bijection of  $\mathfrak{S}_n$  onto  $SE_n$  having the property:*

$$(3.1) \quad (\text{Cyc}, \text{Lmap}) \sigma = (\text{Max}, \text{Rmil}) \text{ B-code } \sigma \quad (\sigma \in \mathfrak{S}_n).$$

*Proof.* By induction, suppose that  $\text{Lmap } \sigma' = \text{Rmil } b'$  and  $\text{Cyc } \sigma' = \text{Max } b'$ . If  $n$  is a fixed point of  $\sigma$ , so that  $\sigma^{-1}(n) = n$  and  $b = (b'_1, \dots, b'_{n-1}, n)$ , then  $\text{Lmap } \sigma = \text{Lmap } \sigma' \cup \{n\} = \text{Rmil } b' \cup \{n\} = \text{Rmil } \sigma$ . Also,  $\text{Cyc } \sigma = \text{Cyc } \sigma' \cup \{n\} = \text{Max } b' \cup \{n\} = \text{Max } b$ .

When  $n$  is not a fixed point of  $\sigma$ , then  $\sigma$  is a product of the form

$$\sigma = \cdots (\cdots \sigma^{-1}(n)n\sigma(n)\cdots) \cdots$$

while  $\sigma'$  may be expressed as

$$\sigma' = \cdots (\cdots \sigma^{-1}(n)\sigma(n)\cdots) \cdots$$

In particular,  $\sigma^{-1}(n) < n, \sigma(n) < n$  and  $\sigma'(\sigma^{-1}(n)) = \sigma(n)$ . We have  $\text{Cyc } \sigma = \text{Cyc } \sigma' = \text{Max } b' = \text{Max } b$  since  $\sigma^{-1}(n) < n$ .

To prove  $\text{Lmap } \sigma = \text{Rmil } b$ , three cases are to be considered, (i)  $\sigma(n) = n-1$ ; (ii)  $\sigma(n) \neq n-1$  and  $\sigma^{-1}(n-1) < \sigma^{-1}(n)$ ; (iii)  $\sigma(n) \neq n-1$  and  $\sigma^{-1}(n-1) > \sigma^{-1}(n)$ , each of them materialized by the following three tableaux:

$$(i) \quad \begin{array}{ccccccc} \text{Id} = & 1 & \cdots & \sigma^{-1}(n) & \cdots & n-1 & n \\ \sigma = & \sigma(1) & \cdots & n & \cdots & \sigma(n-1) & \sigma(n) = n-1 \\ \sigma' = & \sigma(1) & \cdots & \sigma(n) = n-1 & \cdots & \sigma(n-1) & * \end{array}$$

$$(ii) \quad \begin{array}{ccccccc} \text{Id} = & 1 & \cdots & \sigma^{-1}(n-1) & \cdots & \sigma^{-1}(n) & \cdots & n-1 & n \\ \sigma = & \sigma(1) & \cdots & n-1 & \cdots & n & \cdots & \sigma(n-1) & \sigma(n) \\ \sigma' = & \sigma(1) & \cdots & n-1 & \cdots & \sigma(n) & \cdots & \sigma(n-1) & * \end{array}$$

$$(iii) \quad \begin{array}{ccccccc} \text{Id} = & 1 & \cdots & \sigma^{-1}(n) & \cdots & \sigma^{-1}(n-1) & \cdots & n-1 & n \\ \sigma = & \sigma(1) & \cdots & n & \cdots & n-1 & \cdots & \sigma(n-1) & \sigma(n) \\ \sigma' = & \sigma(1) & \cdots & \sigma(n) & \cdots & n-1 & \cdots & \sigma(n-1) & * \end{array}$$

In case (i) we get  $\text{Lmap } \sigma = \text{Lmap } \sigma'$ ,  $b' = (\dots, \sigma^{-1}(n))$  and  $b = (\dots, \sigma^{-1}(n), \sigma^{-1}(n))$ , then  $\text{Rmil } b = \text{Rmil } b'$ .

In case (ii) we clearly have:  $\text{Lmap } \sigma = \text{Lmap } \sigma' \cup \{\sigma^{-1}(n)\}$ . Also,  $b' = (\dots, \sigma^{-1}(n-1))$  and  $b = (\dots, \sigma^{-1}(n-1), \sigma^{-1}(n))$ . Hence,  $\text{Lmap } \sigma = \text{Lmap } \sigma' \cup \{\sigma^{-1}(n)\} = \text{Rmil } b' \cup \{\sigma^{-1}(n)\} = \text{Rmil } b$ .

Finally, comes case (iii), which is the hardest one. We have  $\text{Lmap } \sigma = (\text{Lmap } \sigma' \cap [1, \sigma^{-1}(n) - 1]) \cup \{\sigma^{-1}(n)\}$ , also  $b' = (\dots, b'_{n-2}, \sigma^{-1}(n-1))$ ,  $b = (\dots, b'_{n-2}, \sigma^{-1}(n-1), \sigma^{-1}(n))$ . But as  $\sigma^{-1}(n) < \sigma^{-1}(n-1)$ , we have  $\text{Rmil } b = (\text{Rmil } b' \cap [1, \sigma^{-1}(n) - 1]) \cup \{\sigma^{-1}(n)\} = (\text{Lmap } \sigma' \cap [1, \sigma^{-1}(n) - 1]) \cup \{\sigma^{-1}(n)\} = \text{Lmap } \sigma$ .  $\square$

#### 4. The bijection $\phi$

The bijection  $\phi$ , which is the main ingredient in the chain displayed in (1.4), is simply defined as

$$(4.1) \quad \phi := (\text{B-code})^{-1} \circ \text{A-code}.$$

It follows from Theorems 6 and 5 that

$$\begin{aligned} (\text{Cyc}, \text{Lmap}) \phi(\sigma) &= (\text{Max}, \text{Rmil}) \text{B-code } \phi(\sigma) \\ &= (\text{Max}, \text{Rmil}) \text{A-code } \sigma = (\text{Rmil}, \text{Lmap}) \sigma. \end{aligned}$$

This proves relation (1.1) and consequently Theorem 1. It also follows from Theorem 5 and/or 6 that the distribution of each pair of statistics stated in Theorem 1 is also equal to the distribution of  $(\text{Max}, \text{Rmil})$  on  $\text{SE}_n$ .

It remains to prove identity (1.5) to achieve the proofs of Theorem 2 and its Corollary. Let  $A = ([p_1, q_1], [p_2, q_2], \dots, [p_h, q_h])$  be an ordered partition of  $[n]$  into disjoint non-empty intervals, such that  $p_j + 1 = q_{j+1}$  for  $j = 1, 2, \dots, h-1$  and  $p_1 = 1, q_h = n$ . Let  $G(\sigma) = \{(i, \sigma(i)) : 1 \leq i \leq n\}$  be the graph of a permutation  $\sigma$  from  $\mathfrak{S}_n$ . Referring to Fig. 2, where the square  $[p_j, q_j] \times [p_j, q_j]$  has been materialized by the four points  $B, B'', D'', D$ , we see that  $\sigma$  is  $A$ -indecomposable, if for every  $j = 1, 2, \dots, h$

- (i) the square  $[BB''D''D]$  contains the subgraph  $\{(i, \sigma(i)) : p_j \leq i \leq q_j\}$ ;
- (ii) for every  $l$  such that  $p_j + 1 \leq l \leq q_j$  the rectangle  $[B'B''C''C']$  contains at least one element from  $G(\sigma)$ .

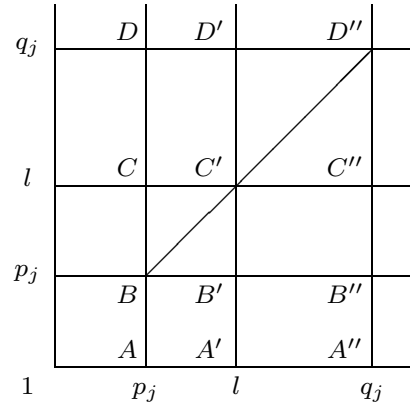


Fig. 2. Graphs of  $\sigma$  and  $c$

We are then led to the following definition.

*Definition.* Each subexcedant sequence  $c = (c_1, c_2, \dots, c_n)$  from  $\text{SE}_n$  is said to be  $A$ -decomposable, if for every  $j = 1, 2, \dots, h$

- (i) the triangle  $[BB''D'']$  contains the subgraph  $\{(i, c_i) : p_j \leq i \leq q_j\}$ ;
- (ii) for every  $l$  such that  $p_j + 1 \leq l \leq q_j$  the rectangle  $[B'B''C''C']$  contains at least one element  $(i, c_i)$  ( $l \leq i \leq q_j$ ).

**Proposition 6.** A permutation  $\sigma$  from  $\mathfrak{S}_n$  is  $A$ -decomposable, if and only if its  $A$ -code (resp.  $B$ -code) is  $A$ -decomposable.

*Proof.* Let  $a = (a_1, a_2, \dots, a_n)$  be the  $A$ -code of a permutation  $\sigma$ . If  $\sigma$  is  $A$ -decomposable, then for every  $j = 1, 2, \dots, h$  and  $l = p_j, p_j + 1, \dots, q_j$  the point  $(\sigma^{-1}(l), l)$  belongs to the square  $[BB''D''D]$ . As  $a_l$  is equal to 1 plus the number of points  $(i, \sigma(i))$  such that  $1 \leq i < \sigma^{-1}(l)$  and  $\sigma(i) < l$ , we have  $a_l \geq p_j$ , so that the point  $(l, a_l)$  belongs to the triangle  $[BB''D'']$ . Conversely, if  $(l, a_l) \in [BB''D'']$ , then  $(\sigma^{-1}(l), l) \in [BB''D''D]$ .

Now, the rectangle  $[B'B''C''C']$  contains no element from  $G(\sigma)$  if and only if all the points  $(\sigma^{-1}(l), l), \dots, (\sigma^{-1}(q_j), q_j)$  are in the square  $[C'C''D''D']$ . This is equivalent to saying that all the quantities  $\sigma^{-1}(l), l, \dots, \sigma^{-1}(q_j), q_j$  lie between  $l$  and  $q_j$ , which is also equivalent to the fact

that  $a_l, \dots, a_{q_j}$  lie between  $l$  and  $q_j$ , that is, the rectangle  $[B'B''C''C']$  has no element  $(i, a_i)$  ( $l \leq i \leq q_j$ ).

Next, let  $b = (b_1, b_2, \dots, b_n)$  be the B-code of  $\sigma$ . If  $\sigma$  is  $A$ -decomposable, the restriction of  $\sigma$  to the interval  $[p_j, q_j]$  is a product of cycles all elements of which lie between  $p_j$  and  $q_j$ . By definition of the B-code all the terms  $b_{p_j}, \dots, b_{q_j}$  also lie between  $p_j$  and  $q_j$  and conversely, if it is the case, all the points  $(p_j, \sigma(p_j)), \dots, (q_j, \sigma(q_j))$  belong to the square  $[BB''D''D]$ . The same argument can be applied when all the points  $(l, \sigma(l)), \dots, (q_j, \sigma(q_j))$  belong to the square  $[C'C''D''D']$ . All terms  $b_l, \dots, b_{q_j}$  are greater than or equal to  $l$  and the rectangle  $[B'B''C''C']$  contains no element of the form  $(i, b_i)$  with  $l \leq i \leq q_j$ .  $\square$

Thus, if  $\sigma$  is  $A$ -decomposable, so are A-code  $\sigma$  and the composition product  $(\text{B-code})^{-1} \text{A-code}(\sigma) = \phi(\sigma)$ . This proves identity (1.5) and then Theorem 2 and its corollary.

### 5. Algorithmic definitions and examples

Although the A-code has been greatly described in various forms (see, e.g., [Kn98], p. 14), we give a full algorithmic definition, which is to be compared with the analogous definition for the B-code.

*Algorithmic definition of A-code.* Let  $\sigma = \sigma(1)\sigma(2)\cdots\sigma(n)$  be a permutation of  $12\cdots n$ . By definition the A-code of  $\sigma$  is the sequence  $a = (a_1, a_2, \dots, a_n)$  where for each  $i = 1, 2, \dots, n$

$$a_i := \#\{j : 1 \leq j \leq i, \sigma^{-1}(j) \leq \sigma^{-1}(i)\},$$

or still

$$(5.1) \quad a_i := \#\{\sigma(k) : 1 \leq \sigma(k) \leq i, k \leq \sigma^{-1}(i)\}.$$

Thus,  $a_i$  is equal to 1 plus the number of letters less than  $i$ , to the left of  $i$ , in the word  $\sigma = \sigma(1)\sigma(2)\cdots\sigma(n)$ .

For instance, with  $\sigma = 4, 6, 1, 2, 3, 5$  the A-code of  $\sigma$  is equal to  $a = (1, 2, 3, 1, 5, 2)$ :  $a_1 = 1$ ,  $a_2 = 2$  as 1 is to the left of 2,  $a_3 = 3$  as 1 and 2 are to the left of 3,  $a_4 = 1$ , as 4 is the leftmost letter of  $\sigma$ , etc. Thus,

$$(5.2) \quad \text{A-code}(4, 6, 1, 2, 3, 5) = (1, 2, 3, 1, 5, 2).$$

*Algorithmic definition of A-code<sup>-1</sup>.* Given  $a = (a_1, a_2, \dots, a_n) \in \text{SE}_n$  write a word with  $n$  empty places numbered 1 to  $n$  from left to right. First, move the letter  $n$  to the  $a_n$ -th leftmost place; let  $\sigma_n$  be the resulting word (having one non-empty letter!). Next, move  $(n-1)$  to the place having  $a_{n-1} - 1$  empty letters to its left. Let  $\sigma_{n-2}$  be the resulting word (having

two non-empty letters). Move  $(n - 2)$  to the place having  $a_{n-2} - 1$  empty letters to its left, etc. Thus,  $\text{A-code}^{-1}(a)$  is the final permutation  $\sigma_1$ .

For instance, start with  $a = (1, 2, 1, 2, 3, 5)$ . We successively get:

$$\begin{array}{rcl}
 & * * * * * & \\
 \sigma_6 = & * * * * 6 * & a_6 = 5 \\
 \sigma_5 = & * * 5 * 6 * & a_5 = 3 \\
 \sigma_4 = & * 4 5 * 6 * & a_4 = 2 \\
 \sigma_3 = & 3 4 5 * 6 * & a_3 = 1 \\
 \sigma_2 = & 3 4 5 * 6 2 & a_2 = 2 \\
 \sigma_1 = & 3 4 5 1 6 2 & a_1 = 1
 \end{array}$$

Thus

$$(5.3) \quad \text{A-code}^{-1}(1, 2, 1, 2, 3, 5) = 3, 4, 5, 1, 6, 2.$$

*Algorithmic definition of B-code.* Let  $\sigma = \sigma(1)\sigma(2)\cdots\sigma(n) \in \mathfrak{S}_n$ . Its B-code  $b = (b_1, b_2, \dots, b_n)$  is calculated as follows. First,  $b_n$  is the place occupied by  $n$  in  $\sigma_n := \sigma$ . Permute the two letters  $n$  and  $\sigma(n)$  in the word  $\sigma$ . Let  $\sigma_{n-1}$  be the resulting word. Then,  $b_{n-1}$  is the place occupied by  $(n - 1)$  in  $\sigma_{n-1}$ . Next, permute the two letters  $(n - 2)$  and  $\sigma(n - 2)$  in  $\sigma_{n-1}$  and let  $\sigma_{n-2}$  be the resulting word. Let  $b_{n-2}$  is the place occupied by  $(n - 2)$  in  $\sigma_{n-2}$ . Permute  $(n - 3)$  and  $\sigma(n - 3)$  in  $\sigma_{n-2}$ , etc. The B-code of  $\sigma$  is  $(b_1, b_2, \dots, b_n)$ .

Start with  $\sigma = 3, 4, 5, 2, 6, 1$ . We successively get:

$$\begin{array}{rcl}
 \text{Id} = & 1 2 3 4 5 6 & \\
 \sigma_6 = & 3 4 5 2 6 1 & b_6 = 5 \\
 \sigma_5 = & 3 4 5 2 1 \mathbf{6} & b_5 = 3 \\
 \sigma_4 = & 3 4 1 2 \mathbf{5 6} & b_4 = 2 \\
 \sigma_3 = & 3 2 1 \mathbf{4 5 6} & b_3 = 1 \\
 \sigma_2 = & 1 2 \mathbf{3 4 5 6} & b_2 = 2 \\
 \sigma_1 = & 1 \mathbf{2 3 4 5 6} & b_1 = 1
 \end{array}$$

Thus

$$(5.4) \quad \text{B-code}(3, 4, 5, 2, 6, 1) = (1, 2, 1, 2, 3, 5).$$

*Algorithmic definition of B-code<sup>-1</sup>.* Let  $b = (b_1, b_2, \dots, b_n) \in \text{SE}_n$ . Start with the identity permutation  $\sigma_1 = 1, 2, \dots, n$ . In  $\sigma_1$  exchange 2 and the letter at the  $b_2$ -th place. Let  $\sigma_2$  be the resulting word. In  $\sigma_2$  permute 3 and the letter at the  $b_3$ -th place. Let  $\sigma_3$  be the resulting word. In  $\sigma_3$  permute 4 and the letter at the  $b_4$ -th place, etc. The permutation  $\sigma = \text{B-code}^{-1} b$  is the permutation  $\sigma_n$ .

For example, starting with  $b = (1, 2, 3, 1, 5, 2)$ . We successively form:

$$\begin{aligned} \sigma_1 &= 1 \mathbf{2} \mathbf{3} \mathbf{4} \mathbf{5} \mathbf{6} & b_1 &= 1 \\ \sigma_2 &= 1 \mathbf{2} \mathbf{3} \mathbf{4} \mathbf{5} \mathbf{6} & b_2 &= 2 \\ \sigma_3 &= 1 \mathbf{2} \mathbf{3} \mathbf{4} \mathbf{5} \mathbf{6} & b_3 &= 3 \\ \sigma_4 &= 4 \mathbf{2} \mathbf{3} \mathbf{1} \mathbf{5} \mathbf{6} & b_4 &= 1 \\ \sigma_5 &= 4 \mathbf{2} \mathbf{3} \mathbf{1} \mathbf{5} \mathbf{6} & b_5 &= 5 \\ \sigma_6 &= 4 \mathbf{6} \mathbf{3} \mathbf{1} \mathbf{5} \mathbf{2} & b_6 &= 2 \end{aligned}$$

Thus,

$$(5.5) \quad \text{B-code}^{-1}(1, 2, 3, 1, 5, 2) = 4, 6, 3, 1, 5, 2.$$

Let  $\Phi := \mathbf{i}\phi\mathbf{i}\phi^{-1}\mathbf{i}$  be the product of the bijections occurring in (1.4). With  $\sigma = 6, 4, 1, 2, 3, 5$  the computation of  $\Phi(\sigma)$  can be made as follows.

$$\begin{aligned} \text{Id} &= 1 \ 2 \ 3 \ 4 \ 5 \ 6 \\ \sigma &= 6 \ 4 \ 1 \ 2 \ 3 \ 5 \\ \mathbf{i}\sigma &= 3 \ 4 \ 5 \ 2 \ 6 \ 1 \\ \text{B-code } \mathbf{i}\sigma &= 1 \ 2 \ 1 \ 2 \ 3 \ 5 && \text{(by (5.4))} \\ \text{A-code}^{-1} \text{B-code } \mathbf{i}\sigma &= \phi^{-1} \mathbf{i}\sigma = 3 \ 4 \ 5 \ 1 \ 6 \ 2 && \text{(by (5.3))} \\ \mathbf{i}\phi^{-1} \mathbf{i}\sigma &= 4 \ 6 \ 1 \ 2 \ 3 \ 5 \\ \text{A-code } \mathbf{i}\phi^{-1} \mathbf{i}\sigma &= 1 \ 2 \ 3 \ 1 \ 5 \ 2 && \text{(by (5.2))} \\ \text{B-code}^{-1} \text{A-code } \mathbf{i}\phi^{-1} \mathbf{i}\sigma &= \phi \mathbf{i}\phi^{-1} \mathbf{i}\sigma = 4 \ 6 \ 3 \ 1 \ 5 \ 2 && \text{(by (5.5))} \\ \Phi(\sigma) &= \mathbf{i}\phi\mathbf{i}\phi^{-1}\mathbf{i}\sigma = 4 \ 6 \ 3 \ 1 \ 5 \ 2. \end{aligned}$$

We verify that

$$(\text{Cyc}, \text{Rmil}) \sigma = (\text{Rmil}, \text{Cyc}) \Phi(\sigma) = (\{1, 2\}, \{1, 2, 3, 5\}).$$

## 6. Concluding remarks and Tables

The bijection constructed by Cori [Cor08] only preserves the *cardinalities* “cyc” and “lmap”, but not the sets “Cyc” and “Lmap.” With the example used in his paper, the permutation

$$\theta = 6, 5, 7, 4, 2, 10, 3, 8, 9 = (1, 6, 10)(2, 5)(3, 7)(4)(8)(9)$$

is mapped onto

$$\theta' = 4, 6, 5, 7, 3, 8, 1, 9, 10, 2 = (1, 4, 7)(2, 6, 8, 9, 10)(3, 5),$$

so that  $(\text{Lmap}, \text{Cyc}) \theta' = (\{1, 2, 4, 6, 8, 9\}, \{1, 2, 3\}) \neq (\{1, 2, 3, 4, 8, 9\}, \{1, 3, 6\}) = (\text{Cyc}, \text{Lmap}) \theta$ . However,  $(\text{cyc}, \text{lmap}) \theta = (\text{lmap}, \text{cyc}) \theta' = (6, 3)$ .



B=	1, 2, 3, 4	1, 2, 4	1, 3, 4	2, 3, 4	1, 4	2, 4	3, 4	4	$\Sigma$
A=1,2,3,4	1								1
1,2,4		1							6
1,3,4			1		1				
2,3,4				1			1	1	
1,4			1		1				11
2,4						1	1	1	
3,4				1		1	2	2	
4				1		1	2	2	6
$\Sigma$	1	6			11			6	

$$n = 4$$

Fig. 3. Distribution of (Cyc, Rmil) over  $\mathfrak{S}_n$ .

There exist other bijections  $\sigma \mapsto a$  such that the sum  $\sum_i (a_i - 1)$  is equal to a statistic different from the inversion number “inv,” but having interesting properties. Let us quote the *Tompkins-Paige method* ([To56, Le60, We61]) for generating permutations on a computer. That method was further used in [Ha92, Ha94] to show that the corresponding sum  $\sum_i (a_i - 1)$  is equal to the *major index* “maj”. Let us also mention the *Denert coding* [FZ90, Ha94], whose sum  $\sum_i (a_i - 1)$  is equal to the *Denert statistic* “den”. Those codings serve to prove that the statistics “inv,” “maj” and “den” are equidistributed on  $\mathfrak{S}_n$ , their common distribution being called *Mahonian*.

Let  $b = (b_1, b_2, \dots, b_n)$  be the B-code of a permutation  $\sigma \in \mathfrak{S}_n$ . In its turn the sum  $\text{env } \sigma := \sum_i (b_i - 1)$  becomes a new Mahonian statistic. Moreover, it follows from the properties of the bijection  $\phi$  defined in (4.1) that the two three-variable statistics (env, Cyc, Lmap) and (inv, Rmil, Lmap) are equidistributed on  $\mathfrak{S}_n$ .

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